

# PATENT ABSTRACTS OF JAPAN

(11)Publication number : 07-123413  
(43)Date of publication of application : 12.05.1995

(51)Int.Cl. H04N 7/30  
G06T 9/00  
H03M 7/30  
H03M 7/40  
H04N 1/41

(21)Application number : **05-265812** (71)Applicant : **SEIKO EPSON CORP  
HUDSON SOFT CO LTD**  
(22)Date of filing : **28.09.1993** (72)Inventor : **MIYANE TOSHIKI  
SEKIMOTO UICHI**

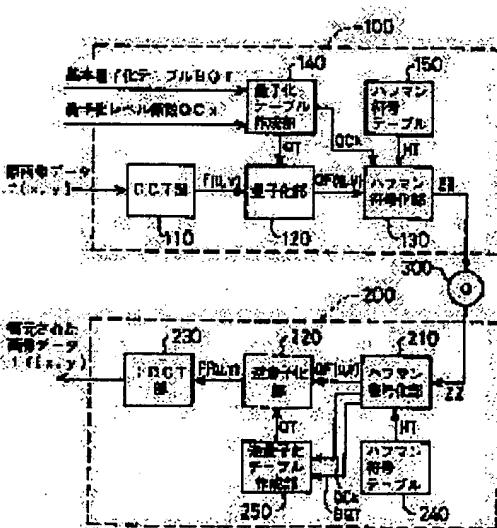
(30)Priority  
Priority number : 04258398 Priority date : 28.09.1992 Priority country : JP  
05 66831 25.03.1993 JP

**(54) METHOD AND DEVICE FOR COMPRESSING/EXTENDING IMAGE DATA**

(57) Abstract:

**PURPOSE:** To reduce the data amount of compressed image data in comparison with the conventional data amount.

**CONSTITUTION:** The compressed image data contains code data showing a quantizing level coefficient QCx among block data. An inverse quantizing table preparation part 250 prepares a quantizing table QT by multiplying a DCT coefficient QF (u,v) provided by decoding compressed image data ZZ and the quantizing level coefficient QCx. An inverse quantizing part 220 performs inverse quantizing by using the quantizing table QT. Since the quantizing level coefficient QCx is inserted between the block data of the compressed image data, the quantizing table QT is updated each time the quantizing level coefficient QCx is decoded. On the other hand, data expressing plural continuous picture element blocks in the same pattern color are inserted to the compressed image data in any special data format (null run data).



## LEGAL STATUS

[Date of request for examination] 28.04.2000

[Date of sending the examiner's decision of rejection] 13.01.2004

[Kind of final disposal of application other than the examiner's decision of rejection or application

converted registration]

[Date of final disposal for application]

[Patent number] 3673529

[Date of registration] 28.04.2005

[Number of appeal against examiner's decision of rejection] 2004-02865

[Date of requesting appeal against examiner's decision of rejection] 12.02.2004

[Date of extinction of right]

\* NOTICES \*

JPO and NCIP are not responsible for any damages caused by the use of this translation.

1. This document has been translated by computer. So the translation may not reflect the original precisely.
2. \*\*\*\* shows the word which can not be translated.
3. In the drawings, any words are not translated.

---

## CLAIMS

---

### [Claim(s)]

[Claim 1] Code data [ as opposed to / are the approach of elongating the compression image data containing the code data obtained by carrying out orthogonal transformation of the image data, quantizing, and carrying out entropy code modulation further, and / the pixel block of (A) plurality ], The 1st [ to the 1st pixel block located in the head of a pixel block of said plurality ] multiplier code, While generating the process for which the compression image data containing the 2nd [ to the 2nd pixel block of arbitration ] multiplier code is prepared, and the transform coefficient quantized by carrying out the entropy decryption of the (B) aforementioned code data the process which generates the 1st and the 2nd multiplier by carrying out the entropy decryption of said 1st and 2nd multiplier code, and (C) -- said the 1st and 2nd multiplier -- respectively -- \*\* -- by multiplying by the base quantity child-sized level of a predetermined base quantity child-sized table While reverse-quantizing the 1st, the process which creates the 2nd quantization table, and said quantized transform coefficient to the pixel block of the (D) above 1st on said 1st quantization table By reverse-quantizing said quantized transform coefficient to said 2nd pixel block on said 2nd quantization table the process which asks for the reverse-quantized transform coefficient, and (E) -- the expanding approach of the compression image data characterized by having the process which asks for the image data elongated by carrying out reverse orthogonal transformation of said transform coefficient by which reverse quantization was carried out.

[Claim 2] It is the expanding approach of compression image data according to claim 1. Said code data While two or more data units arranged according to the array sequence of a pixel block of said plurality are included The expanding approach of compression image data that said 1st multiplier code is arranged just before the 1st [ to said 1st pixel block ] data unit, and said 2nd multiplier code is arranged just before the 2nd [ to said 2nd pixel block ] data unit.

[Claim 3] It is the expanding approach of the compression image data which is an entropy symbolic language in the sign table on which it is the expanding approach of compression image data according to claim 2, and said 1st and 2nd multiplier code contains two or more entropy symbolic languages which receive the dc component of said the 1st, 2nd multiplier, and said transform coefficient.

[Claim 4] It is the expanding approach of compression image data including the process which is the expanding approach of compression image data according to claim 1, and reverse-quantizes said process (D) using said 1st quantization table to a series of pixel blocks from said 1st pixel block to the 3rd pixel block arranged just before said 2nd pixel block.

[Claim 5] It is the expanding approach of the compression image data which includes the process which sets up equally to said maximum the value of the quantization level on which said multiplication result exceeded said maximum when it is the expanding approach of compression image data according to claim 1 and, as for said process (C), the multiplication result of said the 1st and each of the 2nd multiplier, and the quantization level contained in said base quantity child-sized table exceeds predetermined maximum.

[Claim 6] It is the expanding approach of compression image data including the process made into the quantization level of said 1st and 2nd reverse quantization table as it is, without being the expanding

approach of compression image data according to claim 1, and said process (C) performing said multiplication substantially about the predetermined base quantity child-sized level in said base quantity child-sized table.

[Claim 7] It is the expanding approach of the compression image data which is the quantization level are the expanding approach of compression image data according to claim 6, and concerning [ said predetermined base quantity child-sized level ] the dc component of said transform coefficient.

[Claim 8] It is the expanding approach of compression image data according to claim 1. Said compression image data The same pattern block data which shows the number of two or more continuous pixel blocks which has the same image pattern is included. Further furthermore, said process (B) In case said quantized transform coefficient to said two or more pixel blocks expressed with said same pattern block data is created, while setting the predetermined component of said transform coefficient as the value specified beforehand The expanding approach of the compression image data which has the process created by setting the value of the component of said transform coefficients other than said predetermined component as zero.

[Claim 9] It is the expanding approach of compression image data that it is the expanding approach of compression image data according to claim 8, and said predetermined component is a dc component.

[Claim 10] It is the expanding approach of compression image data including the process which is the expanding approach of compression image data according to claim 8, and omits said reverse quantization about said transform coefficient by which said process (D) was decrypted from said same pattern block data.

[Claim 11] Code data [ as opposed to / are the approach of elongating the compression image data containing the code data obtained by carrying out orthogonal transformation of the image data, quantizing, and carrying out entropy code modulation further, and / the 1st pixel block of (A) plurality ], The process for which the compression image data containing the same pattern block data which shows the number of the 2nd continuous pixel block of plurality which has the same image pattern is prepared, (B) While creating the transform coefficient by which the 1st to said two or more 1st pixel blocks was quantized by carrying out the entropy decryption of said code data In case the transform coefficient by which the 2nd to said two or more 2nd pixel blocks expressed with said same pattern block data was quantized is created, while setting the predetermined component of said transform coefficient as the value specified beforehand The process which creates said transform coefficient by which the 2nd was quantized by setting the value of the component of said transform coefficients other than said predetermined component as zero, (C) by reverse-quantizing using a quantization table, said transform coefficient by which the 1st and the 2nd were quantized The expanding approach of the compression image data characterized by having the process which creates the 1st, the process which asks for the transform coefficient by which the 2nd was reverse-quantized, and the image data elongated by carrying out reverse orthogonal transformation of the transform coefficient by which the 2nd was reverse-quantized to the (D) above 1st.

[Claim 12] Code data [ as opposed to / are equipment which elongates the compression image data containing the code data obtained by carrying out orthogonal transformation of the image data, quantizing and carrying out entropy code modulation further, and / two or more pixel blocks ], The 1st [ to the 1st pixel block located in the head of a pixel block of said plurality ] multiplier code, While generating the transform coefficient quantized a storage means to memorize the compression image data containing the 2nd [ to the 2nd pixel block of arbitration ] multiplier code, and by carrying out the entropy decryption of said code data An entropy decryption means to generate the 1st and the 2nd multiplier by carrying out the entropy decryption of said 1st and 2nd multiplier code, By multiplying by said the 1st and each of the 2nd multiplier, and the base quantity child-sized level of a predetermined base quantity child-sized table While reverse-quantizing said quantized transform coefficient to the 1st, a reverse quantization table creation means to create the 2nd quantization table, and said 1st pixel block, on said 1st quantization table By reverse-quantizing said quantized transform coefficient to said 2nd pixel block on said 2nd quantization table The expanding approach of the compression image data characterized by having a reverse quantization means to ask for the reverse-quantized transform

coefficient, and a reverse orthogonal transformation means to ask for the image data elongated by carrying out reverse orthogonal transformation of said transform coefficient by which reverse quantization was carried out.

[Claim 13] It is expanding equipment of compression image data according to claim 12. Said code data While two or more data units arranged according to the array sequence of a pixel block of said plurality are included Expanding equipment of compression image data with which said 1st multiplier code is arranged just before the 1st [ to said 1st pixel block ] data unit, and said 2nd multiplier code is arranged just before the 2nd [ to said 2nd pixel block ] data unit.

[Claim 14] It is expanding equipment of the compression image data which is an entropy symbolic language in the sign table on which it is expanding equipment of compression image data according to claim 13, and said 1st and 2nd multiplier code contains two or more entropy symbolic languages which receive the dc component of said the 1st, 2nd multiplier, and said transform coefficient.

[Claim 15] It is expanding equipment of compression image data including a means to be expanding equipment of compression image data according to claim 12, and to reverse-quantize said reverse quantization means using said 1st quantization table to a series of pixel blocks to the 3rd pixel block arranged just before said 2nd pixel block from said 1st pixel block.

[Claim 16] It is expanding equipment of the compression image data which includes a clipping means set up equally to said maximum the value of the quantization level on which said multiplication result exceeded said maximum when it is expanding equipment of compression image data according to claim 12 and, as for said reverse quantization table creation means, the multiplication result of said the 1st and each of the 2nd multiplier, and the quantization level contained in said base quantity child-sized table exceeds predetermined maximum.

[Claim 17] It is expanding equipment of compression image data including the means which makes said multiplication the quantization level of said 1st and 2nd reverse quantization table as it is about base quantity child-sized level are expanding equipment of compression image data according to claim 12, and predetermined [ in said base quantity child-sized table ] in said reverse quantization table creation means, without carrying out substantially.

[Claim 18] It is expanding equipment of the compression image data which is the quantization level are expanding equipment of compression image data according to claim 17, and concerning [ said predetermined base quantity child-sized level ] the dc component of said transform coefficient.

[Claim 19] It is expanding equipment of compression image data according to claim 12. Said compression image data The same pattern block data which shows the number of two or more continuous pixel blocks which has the same image pattern is included. Further said entropy decryption means In case said quantized transform coefficient to said two or more pixel blocks expressed with said same pattern block data is created, while setting the predetermined component of said transform coefficient as the value specified beforehand Expanding equipment of the compression image data which has a means to create by setting the value of the component of said transform coefficients other than said predetermined component as zero.

[Claim 20] It is expanding equipment of the compression image data said whose predetermined component it is expanding equipment of compression image data according to claim 19, and is a dc component.

[Claim 21] It is expanding equipment of compression image data including a means to be expanding equipment of compression image data according to claim 19, and to bypass said transform coefficient by which said entropy decryption means was decrypted from said same pattern block data, and to supply said reverse quantization means for it to said reverse orthogonal transformation means.

[Claim 22] Code data [ as opposed to / are equipment which elongates the compression image data containing the code data obtained by carrying out orthogonal transformation of the image data, quantizing, and carrying out entropy code modulation further, and / two or more 1st pixel blocks ], The 1st storage means which memorizes the compression image data containing the same pattern block data which shows the number of the 2nd continuous pixel block of plurality which has the same image pattern, The 2nd storage means which memorizes the assignment value over the predetermined

component of the transform coefficient about said two or more 2nd pixel blocks, While creating the transform coefficient by which the 1st to said two or more 1st pixel blocks was quantized by carrying out the entropy decryption of said code data In case the transform coefficient by which the 2nd to said two or more 2nd pixel blocks expressed with said same pattern block data was quantized is created An entropy decryption means to create the transform coefficient by which the 2nd was quantized by setting the value of said transform coefficients other than said predetermined component as zero while setting said predetermined component as the assignment value memorized by said 2nd storage means, By reverse-quantizing using a quantization table, said transform coefficient by which the 1st and the 2nd were quantized Expanding equipment of the compression image data characterized by having a reverse orthogonal transformation means to create the image data elongated the 1st, a reverse quantization means to ask for the transform coefficient by which the 2nd was reverse-quantized, and by carrying out reverse orthogonal transformation of the transform coefficient by which the 2nd was reverse-quantized to said 1st [ the ].

[Claim 23] The process which is the compression approach of image data, carries out orthogonal transformation of the (A) image data for two or more pixel blocks of every in an image, and asks for a transform coefficient, (B) The 1st multiplier to the 1st pixel block located in the head of said pixel blocks of two or more, the process which specifies the 2nd multiplier to the 2nd pixel block of arbitration, and (C) -- said the 1st and 2nd multiplier -- respectively -- \*\* -- by multiplying by the base quantity child-sized level of a predetermined base quantity child-sized table The 1st, the process which creates the 2nd quantization table, and by quantizing said transform coefficient to the (D) above 1st and the 2nd pixel block on said 1st and 2nd quantization table, respectively While creating code data by carrying out entropy code modulation of the transform coefficient by which the (E) aforementioned quantization was carried out to the process which asks for the quantized transform coefficient The compression approach of the image data characterized by having the process which creates the compression image data which contains the 1st, the process which creates the 2nd multiplier code, and the (F) aforementioned code data, said 1st [ the ] and the 2nd multiplier code by carrying out entropy code modulation of said the 1st and 2nd multiplier.

[Claim 24] It is the compression approach of image data which is the compression approach of image data according to claim 23, and contains said 2nd multiplier code just before the 2nd [ to said 2nd pixel block ] data unit just before the 1st [ to said 1st pixel block ] data unit including said 1st multiplier code while said code data contains two or more data units arranged according to the array sequence of a pixel block of said plurality.

[Claim 25] It is the compression approach of the image data which is the compression approach of image data according to claim 24, and includes the process which carries out entropy code modulation of said the 1st and 2nd multiplier using the sign table on which said process (E) contains two or more entropy symbolic languages which receive the dc component of said the 1st, 2nd multiplier, and said transform coefficient.

[Claim 26] It is the compression approach of the image data which is the compression approach of image data according to claim 23, and includes the process which creates the compression image data in which said process (F) contains said same pattern block data with said code data and said 1st and 2nd multiplier code including the process at which said process (E) creates the same pattern block data which shows the number of two or more continuous pixel blocks which has the same image pattern.

[Claim 27] Code data [ as opposed to / are the approach of elongating the compression image data containing the code data obtained by carrying out orthogonal transformation of the image data, quantizing, and carrying out entropy code modulation further, and / the pixel block of (A) plurality ], The process for which the compression image data containing the 1st [ to the 1st pixel block of the arbitration of the pixel blocks of said plurality ] multiplier code is prepared, and by carrying out the entropy decryption of the (B) aforementioned code data The process which generates the 1st multiplier by carrying out the entropy decryption of said 1st multiplier code while generating the quantized transform coefficient, (C) by multiplying by said 1st multiplier and the base quantity child-sized level of a predetermined base quantity child-sized table By reverse-quantizing the process which creates the 1st

quantization table, and said quantized transform coefficient to at least one pixel block including the pixel block of the (D) above 1st on said 1st quantization table the process which asks for the reverse-quantized transform coefficient, and (E) -- the expanding approach of the compression image data characterized by having the process which asks for the image data elongated by carrying out reverse orthogonal transformation of said transform coefficient by which reverse quantization was carried out.

[Claim 28] It is the expanding approach of compression image data according to claim 27. Said compression image data The 2nd [ to the 2nd pixel block arranged after said 1st pixel block ] multiplier code is included. Furthermore, said process (B) The process which generates the 2nd multiplier by carrying out the entropy decryption of said 2nd multiplier code is included. Said process (C) The process which creates the 2nd quantization table by multiplying by said 2nd multiplier and the base quantity child-sized level of said base quantity child-sized table is included. Said process (D) While reverse-quantizing using said 1st quantization table, said quantized transform coefficient to a series of pixel blocks to the 3rd pixel block arranged from said 1st pixel block just before said 2nd pixel block The expanding approach of compression image data including the process which reverse-quantizes said quantized transform coefficient to at least one pixel block including said 2nd pixel block on said 2nd quantization table.

[Claim 29] It is the expanding approach of compression image data according to claim 28. Said code data While two or more data units arranged according to the array sequence of a pixel block of said plurality are included The expanding approach of compression image data that said 1st multiplier code is arranged just before the 1st [ to said 1st pixel block ] data unit, and said 2nd multiplier code is arranged just before the 2nd [ to said 2nd pixel block ] data unit.

---

[Translation done.]

\* NOTICES \*

JPO and NCIP are not responsible for any damages caused by the use of this translation.

1. This document has been translated by computer. So the translation may not reflect the original precisely.
2. \*\*\*\* shows the word which can not be translated.
3. In the drawings, any words are not translated.

---

## DETAILED DESCRIPTION

---

### [Detailed Description of the Invention]

[0001]

[Industrial Application] This invention relates to the approach of compressing and elongating image data, and the equipment for it.

[0002]

[Description of the Prior Art] Drawing 31 is the block diagram showing the configuration of compression/expanding equipment of the conventional image data. After it carries out orthogonal transformation of the subject-copy image data for every block of a MxN pixel in the orthogonal transformation section 542, image data compression equipment 540 quantizes in the quantization section 544, further, encodes in the entropy-code-modulation section 546, and creates compression image data. On the other hand, image data decompression equipment 550 restores image data by the reverse orthogonal transformation section 552, after decrypting compression image data in the entropy decryption section 556 first and reverse-quantizing in the reverse quantization section 554. In addition, the quantization section 544 and the reverse quantization section 554 use the same quantization table 562, and use the same sign table 564 also for the entropy-code-modulation section 546 and the entropy decryption section 556.

[0003]

[Problem(s) to be Solved by the Invention] By the way, the 1st part to restore by high definition in one image and the 2nd part which may be restored by low image quality may be included. In such a case, what is necessary is to quantize using the small quantization table of quantization level to the 1st part, and just to quantize using the big quantization table of quantization level to the 2nd part. A quantization table is the matrix of the same size as the pixel block of image data, i.e., the matrix of a M line N train. With the conventional image data compression / expanding equipment, when two or more quantization tables were used in one image, two or more quantization tables of a M line N train had to be transmitted to expanding equipment 550 from the compression equipment 540, and there was a problem of increasing the amount of data of compression image data.

[0004] Moreover, entropy code modulation of all the orthogonal transformation multipliers of the M line N train over those pixel blocks is carried out also with a simple image part in which the pixel block which has a uniform color continues into an image. Therefore, the compression image data showing such a simple image part also had the problem of becoming the remarkable amount of data.

[0005] This invention is made in order to solve the above-mentioned technical problem in the conventional technique, and it aims at offering the technique in which the amount of data of compression image data can be reduced compared with the former.

[0006]

[Means for Solving the Problem and its Function] In order to solve an above-mentioned technical problem, the compression image data decompression approach by this invention (A) The code data to two or more pixel blocks, and the 1st [ to the 1st pixel block located in the head of said pixel blocks of two or more ] multiplier code, While generating the process for which the compression image data

containing the 2nd [ to the 2nd pixel block of arbitration ] multiplier code is prepared, and the transform coefficient quantized by carrying out the entropy decryption of the (B) aforementioned code data the process which generates the 1st and the 2nd multiplier by carrying out the entropy decryption of said 1st and 2nd multiplier code, and (C) -- said the 1st and 2nd multiplier -- respectively -- \*\* -- by multiplying by the base quantity child-sized level of a predetermined base quantity child-sized table While reverse-quantizing the 1st, the process which creates the 2nd quantization table, and said quantized transform coefficient to the pixel block of the (D) above 1st on said 1st quantization table By reverse-quantizing said quantized transform coefficient to said 2nd pixel block on said 2nd quantization table It has the process which asks for the elongated image data by carrying out reverse orthogonal transformation of the transform coefficient by which the (E) aforementioned reverse quantization was carried out to the process which asks for the reverse-quantized transform coefficient.

[0007] Since the 1st and 2nd quantization table on which quantization level differs can be created only by giving the 1st and the 2nd multiplier, compared with the case where two quantization tables are included in compression image data, the amount of data of compression image data can be reduced.

[0008] While said code data contains two or more data units arranged according to the array sequence of a pixel block of said plurality, said 1st multiplier code is arranged just before the 1st [ to said 1st pixel block ] data unit, and said 2nd multiplier code may be made to be arranged just before the 2nd [ to said 2nd pixel block ] data unit.

[0009] If it carries out like this, the 1st and the 2nd multiplier are easily matched with the 1st and 2nd pixel block, respectively.

[0010] Moreover, as for said 1st and 2nd multiplier code, it is desirable that it is an entropy symbolic language in the sign table containing two or more entropy symbolic languages which receive the dc component of said the 1st, 2nd multiplier, and said transform coefficient.

[0011] if it carries out like this -- the 1st and the 2nd multiplier -- a meaning -- and it can decode in an instant.

[0012] As for said process (D), it is desirable to make it include the process reverse-quantized using said 1st quantization table to a series of pixel blocks from said 1st pixel block to the 3rd pixel block arranged just before said 2nd pixel block.

[0013] Since the 1st multiplier is applied to two or more pixel blocks until the 2nd multiplier appears, it is not necessary to include the 1st multiplier before each code data of a pixel block of these plurality.

[0014] When the multiplication result of said the 1st and each of the 2nd multiplier, and the quantization level contained in said base quantity child-sized table exceeds predetermined maximum, you may make it said process (C) include the process which sets up equally to said maximum the value of the quantization level for which said multiplication result exceeded said maximum.

[0015] If it carries out like this, the number of bits of quantization level can be maintained below at constant value.

[0016] You may make it said process (C) include the process made into the quantization level of said 1st and 2nd reverse quantization table as it is about the predetermined base quantity child-sized level in said base quantity child-sized table, without performing said multiplication substantially.

[0017] If it carries out like this, the error generated in case the transform coefficient of a position is reverse-quantized can be reduced.

[0018] Said predetermined base quantity child-sized level has [ making it be the quantization level about the dc component of said transform coefficient ] a desirable thing.

[0019] If it is made not to change the quantization level to the dc component of a transform coefficient, the error accompanying reverse quantization of a dc component can be reduced.

[0020] Said compression image data contains the same pattern block data which shows the number of two or more continuous pixel blocks which has the still more nearly same image pattern. Further said process (B) In case said quantized transform coefficient to said two or more pixel blocks expressed with said same pattern block data is created, while setting the predetermined component of said transform coefficient as the value specified beforehand You may make it have the process created by setting the value of the component of said transform coefficients other than said predetermined component as zero.

- [0021] If it carries out like this, comparatively little data can express two or more continuous pixel blocks which have the same image pattern of each other.
- [0022] As for said predetermined component, it is desirable that it is a dc component.
- [0023] If it carries out like this, two or more continuous pixel blocks can be applied in the same color.
- [0024] You may make it said process (D) include the process which omits said reverse quantization about said transform coefficient decrypted from said same pattern block data.
- [0025] If it carries out like this, generating of the error by reverse quantization can be prevented.
- [0026] Other compression image data decompression approaches by this invention (A) The process for which the compression image data containing the code data to two or more 1st pixel blocks and the same pattern block data which shows the number of two or more 2nd continuous pixel blocks which has the same image pattern is prepared, (B) While creating the transform coefficient by which the 1st to said two or more 1st pixel blocks was quantized by carrying out the entropy decryption of said code data In case the transform coefficient by which the 2nd to said two or more 2nd pixel blocks expressed with said same pattern block data was quantized is created, while setting the predetermined component of said transform coefficient as the value specified beforehand The process which creates said transform coefficient by which the 2nd was quantized by setting the value of the component of said transform coefficients other than said predetermined component as zero, (C) by reverse-quantizing using a quantization table, said transform coefficient by which the 1st and the 2nd were quantized It has the process which creates the elongated image data the 1st, the process which asks for the transform coefficient by which the 2nd was reverse-quantized, and by carrying out reverse orthogonal transformation of the transform coefficient by which the 2nd was reverse-quantized to the (D) above 1st.
- [0027] Comparatively little data can express two or more continuous pixel blocks which have the color of the same pattern mutually by this approach.
- [0028] Code data [ as opposed to two or more pixel blocks in the compression image data decompression equipment by this invention ], The 1st [ to the 1st pixel block located in the head of a pixel block of said plurality ] multiplier code, While generating the transform coefficient quantized a storage means to memorize the compression image data containing the 2nd [ to the 2nd pixel block of arbitration ] multiplier code, and by carrying out the entropy decryption of said code data An entropy decryption means to generate the 1st and the 2nd multiplier by carrying out the entropy decryption of said 1st and 2nd multiplier code, By multiplying by said the 1st and each of the 2nd multiplier, and the base quantity child-sized level of a predetermined base quantity child-sized table While reverse-quantizing said quantized transform coefficient to the 1st, a reverse quantization table creation means to create the 2nd quantization table, and said 1st pixel block, on said 1st quantization table By reverse-quantizing said quantized transform coefficient to said 2nd pixel block on said 2nd quantization table It has a reverse quantization means to ask for the reverse-quantized transform coefficient, and a reverse orthogonal transformation means to ask for the image data elongated by carrying out reverse orthogonal transformation of said transform coefficient by which reverse quantization was carried out.
- [0029] Moreover, other compression image data decompression equipments by this invention The 1st storage means which memorizes the compression image data containing the code data to two or more 1st pixel blocks, and the same pattern block data which shows the number of the 2nd continuous pixel block of plurality which has the same image pattern, The 2nd storage means which memorizes the assignment value over the predetermined component of the transform coefficient about said two or more 2nd pixel blocks, While creating the transform coefficient by which the 1st to said two or more 1st pixel blocks was quantized by carrying out the entropy decryption of said code data In case the transform coefficient by which the 2nd to said two or more 2nd pixel blocks expressed with said same pattern block data was quantized is created An entropy decryption means to create the transform coefficient by which the 2nd was quantized by setting the value of said transform coefficients other than said predetermined component as zero while setting said predetermined component as the assignment value memorized by said 2nd storage means, By reverse-quantizing using a quantization table, said transform coefficient by which the 1st and the 2nd were quantized It has a reverse orthogonal transformation means to create the

elongated image data, the 1st, a reverse quantization means to ask for the transform coefficient by which the 2nd was reverse-quantized, and by carrying out reverse orthogonal transformation of the transform coefficient by which the 2nd was reverse-quantized to said 1st [ the ].

[0030] The process which the compression approach of the image data based on this invention carries out orthogonal transformation of the (A) image data for two or more pixel blocks of every in an image, and asks for a transform coefficient, (B) The 1st multiplier to the 1st pixel block located in the head of said pixel blocks of two or more, the process which specifies the 2nd multiplier to the 2nd pixel block of arbitration, and (C) -- said the 1st and 2nd multiplier -- respectively -- \*\* -- by multiplying by the base quantity child-sized level of a predetermined base quantity child-sized table The 1st, the process which creates the 2nd quantization table, and by quantizing said transform coefficient to the (D) above 1st and the 2nd pixel block on said 1st and 2nd quantization table, respectively While creating code data by carrying out entropy code modulation of the transform coefficient by which the (E) aforementioned quantization was carried out to the process which asks for the quantized transform coefficient It has the process which creates the compression image data containing the 1st, the process which creates the 2nd multiplier code, and the (F) aforementioned code data, said 1st [ the ] and the 2nd multiplier code by carrying out entropy code modulation of said the 1st and 2nd multiplier.

[0031] Since the 1st and 2nd quantization table on which quantization level differs can be created only by giving the 1st and the 2nd multiplier, compared with the case where two quantization tables are included in compression image data, the amount of data of compression image data can be reduced.

[0032] The approach of elongating other compression image data to the pan by this invention (A) The process for which the compression image data containing the code data to two or more pixel blocks and the 1st [ to the 1st pixel block of the arbitration of said two or more pixel blocks ] multiplier code is prepared, (B) While generating the transform coefficient quantized by carrying out the entropy decryption of said code data By multiplying by the process which generates the 1st multiplier by carrying out the entropy decryption of said 1st multiplier code, the 1st multiplier of (C) above, and the base quantity child-sized level of a predetermined base quantity child-sized table By reverse-quantizing the process which creates the 1st quantization table, and said quantized transform coefficient to at least one pixel block including the pixel block of the (D) above 1st on said 1st quantization table It has the process which asks for the elongated image data by carrying out reverse orthogonal transformation of the transform coefficient by which the (E) aforementioned reverse quantization was carried out to the process which asks for the reverse-quantized transform coefficient.

[0033] [0034] which can reduce the amount of data of compression image data compared with the case where two quantization tables are included in compression image data since the 1st and 2nd quantization table on which quantization level differs can be created only by giving the 1st and the 2nd multiplier, if it carries out like this Said compression image data contains further the 2nd [ to the 2nd pixel block arranged after said 1st pixel block ] multiplier code. Said process (B) The process which generates the 2nd multiplier by carrying out the entropy decryption of said 2nd multiplier code is included. Said process (C) The process which creates the 2nd quantization table by multiplying by said 2nd multiplier and the base quantity child-sized level of said base quantity child-sized table is included. Said process (D) While reverse-quantizing using said 1st quantization table, said quantized transform coefficient to a series of pixel blocks to the 3rd pixel block arranged from said 1st pixel block just before said 2nd pixel block It is desirable to make it include the process which reverse-quantizes said quantized transform coefficient to at least one pixel block including said 2nd pixel block on said 2nd quantization table.

[0035] If it carries out like this, it can quantize by using two quantization tables at the time of expanding of compression image data only by including the 1st and 2nd multiplier code in compression image data.

[0036]

[Example] Below, sequential explanation is performed about each following item.

- A. The whole compression/expanding equipment configuration and basic actuation;
- B. Adjustment of the quantization table QT by the quantization level multiplier QCx;
- C. Huffman coding and configuration of compressed data;

D. Detail configuration of the reverse quantization table creation section;

E. null -- a decryption [0037] of run data A. The configuration and actuation of compression/expanding equipment : drawing 1 is the block diagram showing the function of the compression equipment 100 of image data, and expanding equipment 200 which applied one example of this invention.

[0038] Image data compression equipment 100 is equipped with the DCT section 110 which performs discrete cosine conversion to the subject-copy image data  $f(x, y)$ , the quantization section 120 which quantizes the transform coefficient  $F(u, v)$  obtained by DCT conversion, the Huffman coding section 130 which carries out Huffman coding of the quantized transform coefficient  $QF(u, v)$ , and creates the compression image data ZZ, the quantization table creation section 140, and the Huffman-coding table memory 150. The quantization table creation section 140 creates the quantization table QT based on the base quantity child-sized table BQT and the quantization level multiplier QCx so that it may mention later. The compression image data ZZ is memorized by storages, such as CD-ROM, and is supplied to image data decompression equipment 200 from image data compression equipment 100.

[0039] Image data decompression equipment 200 is equipped with the Huffman decryption section 210 which carries out the Huffman decryption of the compression image data ZZ, the reverse quantization section 220 which reverse-quantizes the transform coefficient  $QF(u, v)$  after the decoded quantization, the IDCT section 230 which performs discrete cosine inverse transformation to the reverse-quantized transform coefficient  $FF(u, v)$ , and obtains image data  $ff(x, y)$ , the Huffman-coding table memory 240, and the reverse quantization table creation section 250. The reverse quantization table creation section 250 creates the quantization table QT for the base quantity child-sized table BQT and the quantization level multiplier QCx which were decoded from the compression image data ZZ based on a receipt and these from the Huffman decryption section 210. This quantization table QT is the same as the quantization table QT used with the compression equipment 100. Moreover, the Huffman-coding table HT memorized by the Huffman-coding table memory 240 is the same as what is memorized by the Huffman-coding table memory 150 of a compression equipment 100.

[0040] Drawing 2 is the block diagram showing the concrete configuration of image data compression equipment 100. This image data compression equipment 100 is equipped with CPU101, main memory 102, a keyboard 103, a mouse 104, a magnetic disk drive 105, and optical-magnetic disc equipment 106. Each processing sections 110-140 of the image data compression equipment 100 shown in drawing 1 are realized by the software program memorized by main memory 102. Moreover, the base quantity child-sized table BQT and the Huffman-coding table HT are memorized by the magnetic disk drive 105. This image data compression equipment 100 is a workstation for creating video game, and various kinds of programs for creating the video game other than the software program for image data compression are performed by CPU101. The completed game program is stored in optical-magnetic disc equipment 106 with compression image data. And CD-ROM containing the program and compression image data of video game is manufactured using this magneto-optic disk.

[0041] Drawing 3 is the block diagram showing the configuration of the video game equipment 20 containing image data decompression equipment 200. The microprocessor 40 which manages in generalization all processings about others, an image processing, and this (henceforth MPU), [ CD-ROM drive / 32 / to which this video game equipment 20 was connected through the SCSI bus 36 ] The main memory 41 by which direct continuation was carried out to this MPU40 (it is hereafter called M-RAM), The various units connected to the bus (M-BUS) 43 of ROM42 and MPU40 which similarly memorized the BIOS program, That is, it has the picture signal control unit 45, the image data decompression unit 200, the VDP unit 49 that outputs a specific picture signal, the video encoder unit 50 which performs composition and the output of a video signal, and the voice data output unit 52 treating voice data.

[0042] Moreover, the memory connected to the local bus (K-BUS) 54 of the picture signal control unit 45 into this video game equipment 20 The memory connected to the local bus of 55 and the image data decompression unit 200 (It is hereafter called K-RAM) The video memory connected to the local bus of 251 and the VDP unit 49 (It is hereafter called R-RAM) (It is hereafter referred to as V=RAM) The output signal from 59 and the video encoder unit 50 is changed into the usual video signal (NTSC), and it has the NTSC converter 60 outputted to color television 28.

[0043] The picture signal control unit 45, the image data decompression unit 200, the video encoder unit 50, and the voice data output unit 52 are constituted by the logical circuit, respectively.

[0044] Drawing 4 (A) is the top view showing an example of the subject-copy image used as the background image of a game. This subject-copy image is an image with which volcanic natural drawing was inserted in the background BG applied in the uniform color. Drawing 4 (B) expands and shows some subject-copy images including one pixel block PB. Generally, the pixel block PB can be set up so that the pixel PX of a MxN individual may be included. As a value of two integers M and N, 8 or 16 are desirable and is M=N=8 in this example. In addition, integers M and N may be set as a different value. The amount of [ which expresses Background BG in the compression image data ZZ ] data division have a special data format (null run data) showing that the pixel block PB of a uniform color continues so that it may mention later.

[0045] The DCT section 110 of image data compression equipment 100 performs two-dimensional DCT conversion for every pixel block PB according to the following formula 1.

[Equation 1]

$$F(u, v) = \frac{1}{4} C(u)C(v) \sum_{x=0}^7 \sum_{y=0}^7 f(x, y) \cos \frac{\pi(2x+1)u}{16} \cos \frac{\pi(2y+1)v}{16}$$

$$\begin{aligned} C(u), C(v) &= \frac{1}{\sqrt{2}} \quad (u, v = 0 \text{ の時}) \\ &= 1 \quad (u, v \neq 0 \text{ の時}) \end{aligned}$$

Here, the array of a transform coefficient, and u and v of the coordinate and F (u, v) the array of 8x8 image data by which f (x y) is contained in one pixel block PB, x, and y indicate the location of each pixel within each pixel block PB to be are the coordinates of frequency space.

[0046] Drawing 5 is the explanatory view showing the array of a transform coefficient F (u, v). A transform coefficient F (u, v) is the same array of 8x8 as the pixel block PB. The transform coefficient F of an upper left edge (0 0) is called DC component (or DC multiplier), and other transform coefficients are called AC component (or AC multiplier). DC component shows the average of the image data in the pixel block PB. Moreover, AC component shows change of the image data within the pixel block PB. Since the image data of the adjoining pixel has a certain amount of correlation, the value of a low-frequency component is comparatively large in AC multiplier, and the value of a high frequency component is comparatively small. Moreover, the effect which it has on image quality has a comparatively small high frequency component.

[0047] Drawing 6 is the explanatory view showing basic actuation of image data compression equipment 100 and image data decompression equipment 200. The DCT section 110 creates the DCT multiplier F (u, v) shown in drawing 6 (a).

[0048] The quantization table creation section 140 creates the quantization table QT (drawing 6 (d)) by multiplying by the base quantity child-sized table BQT (drawing 6 (c)) and the quantization level multiplier QCx, as shown in the following formula 2.

[Equation 2]

$$QT(u, v) = QCx \times BQT(u, v)$$

[0049] Since it is QCx=1 in the example of drawing 6, the quantization table QT is the same as the base quantity child-sized table BQT.

[0050] The quantization section 120 asks for the quantized DCT multiplier QF (u, v) which is shown in drawing 6 (b) by carrying out linear quantization of the DCT multiplier F (u, v) on the quantization table QT. Linear quantization is processing which does a division and rounds off the division result for an integer.

[0051] The Huffman coding section 130 creates the compression image data ZZ (drawing 6 (e)) by carrying out Huffman coding of this DCT multiplier QF (u, v). In addition, about the approach of

Huffman coding, it mentions later further. The compression image data ZZ contains the 1st data showing the base quantity child-sized table BQT, and the 2nd data showing the quantization level multiplier QCx and a transform coefficient QF (u, v) so that it may mention later.

[0052] If the compression image data ZZ is given to image data decompression equipment 200, the Huffman decryption section 210 will decrypt the compression image data ZZ, and it will ask for the DCT multiplier QF (u, v) (drawing 6 (f)). Since Huffman coding is reversible coding, this DCT multiplier QF (u, v) is the same as the DCT multiplier QF (u, v) (drawing 6 (b)) after the quantization called for by the quantization section 120 of image data compression equipment 100. In addition, the Huffman decryption section 210 also decrypts the base quantity child-sized table BQT (drawing 6 (c)) and the quantization level multiplier QCx which are contained in the compression image data ZZ other than the DCT multiplier QF (u, v), and is given to the reverse quantization table creation section 250.

[0053] The reverse quantization table creation section 250 creates the quantization table QT (drawing 6 (d)) by carrying out the multiplication of the base quantity child-sized table BQT and the quantization level multiplier QCx. The reverse quantization section 220 carries out the multiplication of this quantization table QT and the DCT multiplier QF (u, v), and asks for the decoded DCT multiplier FF (u, v) which is shown in drawing 6 (g).

[0054] The IDCT section 230 performs two-dimensional DCT inverse transformation shown in the following formula 3 to this DCT multiplier FF (u, v), and creates the restored image data ff (x y).

[Equation 3]

$$ff(x,y) = \frac{1}{4} \sum_{u=0}^7 \sum_{v=0}^7 C(u)C(v) FF(u,v) \cos \frac{\pi(2x+1)u}{16} \cos \frac{\pi(2y+1)v}{16}$$

$$C(u), C(v) = \frac{1}{\sqrt{2}} \quad (u, v = 0 \text{ の時})$$

$$= 1 \quad (u, v \neq 0 \text{ の時})$$

[0055] B. Adjustment of the quantization table QT by the quantization level multiplier QCx : since it is created by carrying out the multiplication of the base quantity child-sized table BQT and the quantization level multiplier QCx according to said formula 2, the quantization table QT can enlarge each quantization level in the quantization table QT, if the value of the quantization level multiplier QCx is enlarged. In case the value of the quantization level multiplier QCx compresses image data in image data compression equipment 100, an operator chooses it from two or more values (0-15) defined beforehand.

[0056] Drawing 7 is the explanatory view showing compression/expanding actuation at the time of specifying the quantization level multiplier QCx as 4. The quantization table creation section 140 and the reverse quantization table creation section 250 create the quantization table QT shown in drawing 7 (d) according to the above-mentioned formula 2. However, in this example, the maximum of quantization level is restricted to 15 and all the values of the quantization level from which the result of multiplication becomes 15 or more are compulsorily set as 15.

[0057] If linear quantization of the DCT multiplier F (u, v) of drawing 7 (a) is carried out using the quantization table QT shown in drawing 7 (d), as shown in drawing 7 (b), DC component will be 1 and the DCT multiplier QF (u, v) all whose AC components are 0 will be obtained. Thus, if the value of the quantization level multiplier QCx is enlarged, since the number of 0 in the quantized DCT multiplier QF (u, v) will increase, a data compression rate can be raised. However, since the decoded DCT multiplier FF (u, v) which is shown in drawing 7 (g) shows the DCT multiplier F (u, v) and a value which becomes and is different of the origin shown in drawing 7 (a), it is larger than the case (drawing 6) where the quantization level multiplier QCx is equal to 1. [ of degradation of image quality ]

[0058] By the way, since DC component of a DCT multiplier shows the average of the image data within the pixel block PB, its effect to image quality is quite large. Therefore, it is desirable to keep the quantization level for DC components the same as the value in the base quantity child-sized table BQT

irrespective of the value of the quantization level multiplier QCx. Drawing 8 is the explanatory view showing the processing in such a case, and the quantization level QT (0 0) for DC components is kept compulsory at 1. In the case of QT(0 0)=1, since DC component of the decoded DCT multiplier FF (u, v) which is shown in drawing 8 (g) is maintained at the same value as DC component of the original DCT multiplier F (u, v) shown in drawing 8 (a), degradation of image quality can be suppressed comparatively small with compressibility comparable as the case of drawing 7. In addition, it is not necessary to necessarily set quantization level QT (0 0) for DC components to 1, and it may set up other any value.

[0059] It is also possible to choose 0 as a quantization level multiplier QCx. As shown at drawing 9 in the case of QCx=0, all the quantization level in the quantization table QT is set as 1. Since the quantized DCT multiplier QF (u, v) is the same as the original DCT multiplier F (u, v), although compressibility is small, compression/expanding can be performed by high definition.

[0060] If the above is summarized, the quantization table creation section 140 and the reverse quantization table creation section 250 have the following descriptions.

- (1) Create the quantization table QT by carrying out the multiplication of the base quantity child-sized table BQT and the quantization level multiplier QCx according to a formula 2.
- (2) The result of multiplication sets up equally to maximum the thing more than the maximum (= 15) of quantization level (drawing 6 - drawing 8).
- (3) Keep the quantization level QT (0 0) for DC components the same as that of the value in the base quantity child-sized table BQT irrespective of the value of the quantization level multiplier QCx (drawing 8).

(4) When the quantization level multiplier QCx is 0, set all quantization level as 1 (drawing 9).

[0061] The above-mentioned operation which asks for the quantization table QT is performed by the software program in image data compression equipment 100, and it is carried out by the reverse quantization table creation section 250 which consisted of hardware of dedication in image data decompression equipment 200. About the concrete circuitry of the reverse quantization table creation section 250, it mentions later further.

[0062] In addition, in case the quantization level multiplier QCx compresses image data with image data compression equipment 100, it can specify a value which uses a keyboard 103 and a mouse 104 and is different for every pixel block PB.

[0063] C. Huffman coding and the configuration of compressed data; the Huffman coding section 130 (drawing 1) of image data compression equipment 100 consists of the DC multiplier coding section and the AC multiplier coding section. Drawing 10 (A) is the block diagram showing the function of DC multiplier coding section. the block delay section 131 and an adder 132 are shown in drawing 10 (B) -- as -- DC multiplier DCi of each pixel block PB DC multiplier DCi-1 of the pixel block PB before one difference -- \*\*DC is computed.

[0064] the category-sized table showing the categorizing processing section 133 in drawing 11 -- following -- the difference of DC multiplier -- it asks for Category SSSS and the discernment data ID corresponding to \*\*DC. Category SSSS -- the difference of DC multiplier -- it is the number which shows the range of \*\*DC. two or more difference as which the discernment data ID are specified by Category SSSS -- it is data in which it is shown from the smaller one in \*\*DC the value of what position it is.

[0065] Category SSSS is changed into the Huffman-coding word HFDC for DC multipliers in the 1 more-dimensional Huffman coding section 134 (drawing 10). Drawing 12 is the explanatory view showing an example of the Huffman-coding table HTDC used by the 1-dimensional Huffman coding section 134. In this example, the subject-copy image data f (x y) shall be expressed by the YUV signal (a luminance signal Y and two color-difference signals U and V). The Huffman-coding table for DC multipliers of U signal / V signal common use only contains the symbolic language of the category SSSS of 0-9. On the other hand, the Huffman-coding table for DC multipliers for Y signals contains the symbolic language of the category SSSS of 15-31 other than the symbolic language of the category SSSS of 0-9. the null which mentions the Huffman-coding word of SSSS=15 later -- it is shown that it is

run data. null -- run data are data in which it is shown that the pixel block PB of a uniform color continues. Moreover, the Huffman-coding word of SSSS=16-31 is a sign which shows the value of the quantization level multiplier QCx. For example, the Huffman-coding word "111110000" over SSSS=16 shows QCx=0, and the Huffman-coding word "111111111" over SSSS=31 shows QCx=15. In addition, meaning decode is possible for the Huffman-coding word of drawing 12 about all of category SSSS=1-9, and 15-31, and instant decode is possible for it.

[0066] Drawing 13 is the block diagram showing the function of AC multiplier coding section in the Huffman coding section 130. List direct [ of the array F of AC multiplier (u, v) (except for u=v=0) ] is first carried out to one dimension by the zigzag scan section 135. Drawing 14 is the explanatory view showing the usual route of a zigzag scan.

[0067] The value of AC multiplier by which list direct [ of the judgment section 136 ] was carried out to one dimension judges whether it is 0. If the value of AC multiplier is 0, the run length counter 137 will change AC multiplier of continuous 0 into the zero run length NNNN. If AC multiplier is not 0, the value of the AC multiplier will be changed into Category SSSS and the discernment data ID by the categorizing section 138. Under the present circumstances, the category-sized table shown in drawing 11 is referred to.

[0068] The zero run length NNNN and Category SSSS are changed into the Huffman-coding word HFAC for AC multipliers in the two-dimensional Huffman coding section 139. Drawing 15 is the explanatory view showing the two-dimensional Huffman-coding table HTAC for AC multipliers.

Moreover, drawing 16 shows an example of the Huffman-coding word of the parts (topmost part of two lines in drawing 15 ) of NNNN=0 and NNNN=1 in the Huffman-coding table HTAC. In addition, the Huffman-coding word "11111" of NNNN/SSSS=0/0 shows termination of the code data to one pixel block.

[0069] Drawing 17 is the explanatory view showing an example of Huffman coding. Drawing 17 (B) shows coding of DC multiplier. When the value of DC multiplier in the pixel block before one is assumed to be 0, it is \*\*DC=F(0 0)=12. According to the category-sized table of drawing 11, the category SSSS of \*\*DC=12 is 4 and the discernment data ID are "1100." Moreover, according to the Huffman-coding table for DC multipliers of drawing 12, the Huffman-coding word HFDC of category SSSS=4 is "011." In addition, the Huffman-coding table for Y signals is used here. Huffman coding (HF+ID) to DC multiplier is set to "0111100" as shown in drawing 17 (B).

[0070] Drawing 17 (C) shows coding of AC multiplier. First, AC multiplier is put in order by the array of a single dimension with a zigzag scan. This array is changed into the category SSSS of the value which are not the zero run length NNNN and zero (refer to drawing 11). The combination of the zero run length NNNN and Category SSSS is changed into the Huffman-coding word HFAC by the Huffman-coding table for AC multipliers shown in drawing 15 and drawing 16, and it is combined with the discernment data ID of AC multiplier which is not zero, and as shown in drawing 17 (C), Huffman coding (HFAC+ID) is created.

[0071] Drawing 18 is the explanatory view showing the configuration of compressed data. The whole compressed data consists of a header unit, the compressed data section, and the dummy section, as shown in drawing 18 (A). The header unit has 1 byte of four data DFH, DFL, DLH, and DLL, respectively. The first two data DFH and DFL show the class of data contained in the compressed data section. There is a class of the data of the base quantity child-sized table BQT, full color natural picture compression data, run length picture compression data, etc. of the data of the compressed data section. Data (DLH+DLL) of 16 bits of posterior parts of a header show the data length of the sum total of the compressed data section and the dummy section. Since the compressed data section is variable-length data containing Huffman coding, the data length of the sum total with the dummy section is adjusted so that it may become the die length of the integral multiple of WORD (= 2 bytes).

[0072] Drawing 18 (B) shows the configuration of the compressed data showing the base quantity child-sized table BQT. This one-set compressed data contains the data showing the base quantity child-sized table BQT for Y signals, and the data showing the base quantity child-sized table BQT of U signal / V signal common use. In addition, it is not necessary to carry out Huffman coding of the data showing the

base quantity child-sized table BQT.

[0073] Drawing 18 (C) shows the configuration of the compressed data of a full color natural image. the code data (symbolic language of category SSSS=16-31 in drawing 12 ) which expresses the quantization level multiplier QCx to the compressed data section, the block data which is code data of each pixel block, and the null which shows two or more pixel blocks of a uniform color -- run data are included.

[0074] As shown in drawing 18 (D), the block data of one unit consists of 4 sets of data for Y signals, 1 set of U signal data, and 1 set of data for V signals. Drawing 19 is the explanatory view showing the relation of a block of each signal of YUV. As shown in drawing 19 (A), one screen in this example has the magnitude of the 256 pixel x240 scanning line. About a Y signal, DCT conversion is performed for every 8x8-pixel pixel block, without culling out. On the other hand, about U signal and V signal, as shown in drawing 19 (B), it is thinned out and (subsampling) made one half in a longitudinal direction and a lengthwise direction, and DCT conversion is performed to the 8x8-pixel block after thinning out. Therefore, as shown in drawing 19 (C), the field of four pixel blocks Y1-Y4 of a Y signal is equivalent to the field of one pixel block of U signal and V signal. In addition, it is because it is comparatively insensible to change (change of U signal and V signal) of a color to operate U signal and V signal on a curtailed schedule, without operating a Y signal on a curtailed schedule although human being's eyes are comparatively sensitive to change (change of a Y signal) of brightness. Compressibility can be raised by thinning out only U signal and V signal, without degrading image quality too much. In addition, the block data of one unit shown in drawing 18 (D) arranges in order the Huffman-coding data of each field shown in drawing 19 (C).

[0075] The code data to one pixel block in block data consists of one Huffman-coding data of DC multiplier, and two or more Huffman-coding data of AC multiplier, as shown in drawing 18 (F). The Huffman-coding data of DC multiplier consist of a Huffman-coding word HFDC of Category SSSS, and discernment data ID, as mentioned above ( drawing 18 (G)). Moreover, the Huffman-coding data of AC multiplier consist of the Huffman-coding words HFAC and the discernment data ID to combination with Category SSSS with the zero run length NNNN ( drawing 18 (H)).

[0076] The code data of the quantization level multiplier QCx is inserted just before the block data of the pixel block which wants to change the head of the compressed data section, and the value of the quantization level multiplier QCx. To two or more pixel blocks before the 2nd quantization level multiplier QCx is inserted, the top quantization level multiplier QCx is used in common. Moreover, to two or more blocks before the 3rd quantization level multiplier QCx (not shown) is inserted, the 2nd quantization level multiplier QCx is used in common.

[0077] In addition, when the symbolic language of the quantization level multiplier QCx is not contained at the head of the compressed data section, it is considered that it is QCx=1. Therefore, also when the quantization level multiplier QCx is not inserted in the head of the compressed data section but the quantization level multiplier QCx is inserted only once on the way, it is equivalent to two quantization level multipliers QCx being specified.

[0078] Since Huffman coding showing the quantization level multiplier QCx is inserted between block data, it can apply this new quantization level multiplier QCx easily to the next block data at the time of the new quantization level multiplier QCx being decrypted. Moreover, since the code data of the quantization level multiplier QCx is expressed with the Huffman-coding word for DC multipliers as shown in drawing 12 , even if this is inserted between block data, it is possible to judge immediately whether this code data is code data of DC multiplier for block Y1 or it is the code data of the quantization level multiplier QCx.

[0079] the null contained in the compressed data section -- run data are shown in drawing 18 (E) -- as -- null -- it consists of discernment data ID with the symbolic language for DC multipliers "NRL" which shows that it is run data, and the block count.

[0080] drawing 20 -- null -- it is the explanatory view showing the image expressed by run data. The background BG of the subject-copy image of drawing 20 (A) is applied in the uniform color. The part of the ellipse of drawing 20 (A) is assumed to be what 18 pixel blocks which have the image data value ( $f(x, y) = 12$ ) with all the same pixels as shown in drawing 20 (B) are following. the null to which drawing

20 (C) expresses these pixel blocks -- run data are shown. this null -- run data -- the 1st null for a 16-pixel block -- the run data NRD1 and the 2nd null for a 2-pixel block -- the run data NRD2 are included.

[0081] each -- null -- the head of the run data NRD1 and NRD2 -- null -- it has the symbolic language for DC multipliers "NRL" (symbolic language of category SSSS=15 of drawing 12 "1111011") which shows that it is run data. decrypting the symbolic language for DC multipliers in a head, since Huffman coding of DC multiplier is arranged at the head of the usual block data as shown in drawing 18 (F) -- null -- run data, block data, and the code data of the quantization level multiplier QCx -- a meaning -- and it is discriminable in an instant.

[0082] The block count is expressed with the Huffman-coding word for AC multipliers as shown in drawing 20 (C). drawing 21 -- the inside of the Huffman-coding table for AC multipliers ( drawing 15 ) -- null -- it is drawing showing the part used for run data. null -- when used for run data, if the zero run length NNNN is equal to ([block count]-1), he will be set up. Moreover, the Huffman-coding word of category SSSS=1 is used noting that the value of AC multiplier is 1. the 1st null shown in drawing 20 (C) -- the data (NNNN/SSSS=15/1) of the block count in the run data NRD1 show that 16 pixel blocks of a uniform color are continuing. moreover, the 2nd null -- the data (NNNN/SSSS=1/1) of the block count in the run data NRD2 show that two pixel blocks of a uniform color are continuing.

[0083] each -- null -- the discernment data ID are added to the back end of the run data NRD1 and NRD2. It is fixed to ID=1 in this example.

[0084] null -- run data can express that two or more pixel blocks which continued with about 20-bit data are uniform colors in this way for the usual block data's, expressing the one-set block (the Y signal shown in drawing 19 -- a 4-pixel block, U signal, and V signal -- an every 1pixelblock -- it contains) of a uniform color on the other hand -- about 300- about 400 bits is required. And also when it is shown that a two or more sets pixel block is a uniform color, it is [ about 300 - 400 bits of abbreviation ] required about each set. therefore, null -- if run data are used, it is possible to reduce considerably the amount of data of the compressed data showing the pixel block of a majority of continuous uniform colors.

[0085] in addition, null -- the value of the luminance signal Y of a pixel block of a uniform color and color-difference signals U and V expressed by run data is not included in compressed data, but is specified in the software program which describes video game. He specifies the brightness and color tone of these blocks using a keyboard 103 or a mouse 104 while he specifies the range of the field ( drawing 20 (A) the background BG) of a pixel block of a uniform color with a mouse 104, in case an operator creates the software program for video game. When carrying out like this and a specific event occurs while having performed the game, for example using video game equipment 20 ( drawing 3 ), the special visual effectiveness of changing the color of Background BG in time can be produced. in addition, null -- about the concrete circuitry which decrypts run data, it mentions later further.

[0086] D. The detail configuration of the reverse quantization table creation section : drawing 22 is the block diagram showing the internal configuration of the reverse quantization table creation section 250 shown in drawing 1 . The reverse quantization table creation section 250 is equipped with the multiplication unit 254 which carries out the multiplication of the address-generation circuit 252 which generates the address of RAM251 and RAM251 which memorizes the base quantity child-sized table BQT, the latch circuit 253 holding the quantization level multiplier QCx, and the quantization level multiplier QCx and the base quantity child-sized table BQT, and generates the quantization table QT. The quantization table QT which was caused multiplication unit 254 and created is supplied to the reverse quantization section 220.

[0087] First, if the compression image data ZZ contained by CD-ROM is given to the Huffman decryption section 210, the code data of the base quantity child-sized table BQT will be decrypted first, and will be supplied to RAM251. This base quantity child-sized table BQT is memorized by RAM251 according to the light address given from the address-generation circuit 252. The base quantity child-sized table BQT memorized by RAM251 is used to all pixel blocks.

[0088] The address-generation circuit 252 generates the lead address synchronizing with the DCT multiplier data QF (u, v) outputted from the Huffman decryption section 210, and the base quantity

child-sized table BQT is read from RAM251 according to this lead address. On the other hand, the quantization level multiplier QCx decrypted in the Huffman decryption section 210 is latched by the latch circuit 253, and it is saved at a latch circuit 253 until the following quantization level multiplier QCx is given. Therefore, the same quantization level multiplier QCx is used in common to two or more pixel blocks until the quantization level multiplier QCx is newly supplied.

[0089] Drawing 23 is the block diagram showing the internal configuration of the latch circuit 253 contained in the reverse quantization table creation section 250 (drawing 22), and the multiplication unit 254. The latch circuit 253 consists of two latches 402,404. The multiplication unit 254 has the synchronous-clock creation circuit 412, AND circuit 414, U signal start detector 416, V signal start detector 418, NAND circuit 420, the selector 422, the multiplier 424, the clipping circuit 426, and the zero value correction circuit 428.

[0090] Drawing 24 is a timing chart which shows actuation of the circuit shown in drawing 23. If the quantization level multiplier QCx is decrypted in the Huffman decryption section 210 (drawing 1), the data of the quantization level multiplier QCx will be given to a latch circuit 253 with an enable signal QEN (drawing 24 (a), (b)). The 1st latch 402 latches the quantization level multiplier QCx by the rising edge of an enable signal QEN, and supplies an output Q1 to the 2nd latch 404 (drawing 24 (c)).

[0091] As shown in drawing 24 (d), after the base quantity child-sized table BQT for Y signals is read 4 times, the base quantity child-sized table BQT of U signal / V signal common use is read twice from RAM251 (drawing 22).

[0092] Synchronizing with the base quantity child-sized table BQT read from RAM251, it is set to L level in the period of a Y signal, and block recognition signal UV/Y used as H level is given to the synchronous-clock creation circuit 412 (drawing 23) from the address-generation circuit 252 in the period of U signal and V signal (drawing 24 (e)). Moreover, an enable signal EN is also given to the synchronous-clock creation circuit 412 in this case. The synchronous-clock creation circuit 412 reverses block recognition signal UV/Y, generates a synchronizing clock signal SCK (drawing 24 (f)), and supplies this to the 2nd latch's 404 clock input terminal. The 2nd latch 404 latches the 1st latch's 402 output Q1 by the rising edge of this synchronizing clock signal SCK, and supplies that output Q2 (drawing 24 (g)) to the data input terminal of a selector 422. In addition, the fixed value "1" is given to other data input terminals of a selector 422.

[0093] U signal start detector 416 generates U start signal USTRT (drawing 24 (h)) which shows the start time of day of the base quantity child-sized table BQT for U signals based on the block recognition signals UV/Y and the 10MHz basic clock signal CLK. This U start signal USTRT is a signal set to L level from the rising edge of block recognition signal UV/Y only for 100 nanoseconds.

[0094] Six blocks Y1-Y4, and the block change signal SWTCH (drawing 24 (i)) and the block recognition signals UV/Y with which level changes by turns in each period of U and V are given to AND circuit 414. V signal start detector 418 generates V start signal VSTRT (drawing 24 (j)) which shows the start time of day of the base quantity child-sized table BQT for V signals based on the output of AND circuit 414, and the 10MHz basic clock signal CLK. Block recognition signal UV/Y of this V start signal VSTRT is the signal set to L level from the rising edge of the block change signal SWTCH only for 100 nanoseconds in the period of H level.

[0095] Drawing 25 is the block diagram showing the internal configuration of U signal start detector 416 and V signal start detector 418. These circuits 416,418 are constituted from a D flip-flop and a NAND circuit by both. The block recognition signals UV/Y are supplied to D input terminal of D flip-flop 432 of U signal start detector 416, and the 10MHz basic clock signal CLK is inputted into the clock input terminal. Block recognition signal UV/Y and the reversal output of D flip-flop 432 are given to the input terminal of NAND circuit 434. In addition, block recognition signal UV/Y synchronizes with the basic clock signal CLK.

[0096] Since the block recognition signals UV/Y is [the reversal output of D flip-flop 432] H level between L level, the output USTRT of NAND circuit 434 is maintained at H level (refer to drawing 24 (h)). If the output (U start signal USTRT) of NAND circuit 434 is set to L level and D flip-flop 432 latches an input with the edge of the basic clock signal CLK after 100 nanoseconds immediately after

block recognition signal UV/Y changes from L level to H level, the output USTRT of NAND circuit 434 will return to H level again.

[0097] Actuation of V signal start detector 418 is the same as actuation of U signal start detector 416. However, since the AND of block recognition signal UV/Y and the block change signal SWTCH is given to D input terminal of D flip-flop 436, in the period of H level, V start signal VSTRT is set to L level from the rising edge of the block change signal SWTCH by block recognition signal UV/Y only for 100 nanoseconds.

[0098] U start signal USTRT and V start signal VSTRT are inputted into NAND circuit 420 (drawing 23), and the output (selection signal SEL) is supplied to the selection input terminal of a selector 422. A selection signal SEL (drawing 24 (k)) is set to H level only for 100 nanoseconds at the beginning of the period the start of the period for U signals, and for V signals. A selector 422 outputs the output Q2 given from the latch circuit 253 as it is, when a selection signal SEL is L level, and on the other hand, when a selection signal SEL is H level, it outputs a fixed value "1." The multiplication of the output Q3 of a selector 422 is carried out to the base quantity child-sized table BQT by the multiplier 424.

[0099] As shown in drawing 24 (l), the output Q3 of a selector 422 is surely set to "1" irrespective of the value of the quantization level multiplier QCx at the beginning of the period the start of the period for U signals, and for V signals. The period for 100 nanoseconds which the period the object for U signals and for V signals begins is a period for computing the quantization level for DC multipliers. Therefore, if the circuit of drawing 23 is used, it can avoid performing substantially the multiplication of the quantization level for DC multipliers used for U signal and V signal, and the specified quantization level multiplier QCx. If it puts in another way, the circuit shown in drawing 23 will be a circuit which realizes the operation shown in drawing 8 (c) and (d).

[0100] As shown in drawing 23, the output of a multiplier 424 is corrected by a clipping circuit 426 and the zero value correction circuit 428, and serves as the final quantization table QT. Drawing 26 is the block diagram showing the internal configuration of these two circuits 426, 428.

[0101] The clipping circuit 426 consists of 4 input OR circuit 450 and eight 2 input OR circuits 452. These circuits are circuits in the case of expressing quantization level by 9 bits (the most significant bit being a sign bit). D9-D12 are inputted into 4 input OR circuit 450 4 bits of high orders except a sign bit D13 among the 14-bit data outputted from the multiplier 424 (drawing 23). The output of 4 input OR circuit 450 is given to one input terminal of eight 2 input OR circuits 452, and D1-D8 are given to the input terminal of another side 8 bits of low order of the output of a multiplier 424. When at least one value of D9-D12 is "1" 4 bits of high-orders, all of the output of eight 2 input OR circuits 452 are set to "1." Therefore, the output of a clipping circuit 426 is set to the case of 255 or more for the output of a multiplier 424 by 255 with a decimal number.

[0102] In the zero value correction circuit 428, the output of seven 2 input OR circuits 452 seven bits D2 in a clipping circuit 426 - for D8 is outputted as it is. Moreover, the output and sign bit D13 of these seven 2 input OR circuits 452 are given to eight inverters 460, respectively. The output of eight inverters 460 is given to 8 input AND circuit 462, and the output of this AND circuit 462 is supplied to 2 input OR circuit 464. The output of 2 input OR circuit 452 for least significant bit D1 is given to this 2 input OR circuit 464. Consequently, when [ of the output of a multiplier 424 ] all of 13 bits of values of D1-D13 are "0", as for the zero value correction circuit 428, the value of 8 bits of others [ value / of the least significant bit ] outputs the quantization level QT of "0" by "1." If it puts in another way, the zero value correction circuit 420 will have realized the operation shown in drawing 9 (c) and (d).

[0103] E. null -- decryption [ of run data ]: -- drawing 27 is the block diagram showing the internal configuration of the Huffman decryption section 210 (drawing 1) in image data decompression equipment 200. The Huffman decryption section 210 is equipped with the decryption section 470 which carries out the Huffman decryption of the compression image data ZZ, the control section 472, the selector 474, and DC multiplier register 476.

[0104] the class of compressed data with which the decryption section 470 was given -- the base quantity child-sized table BQT, the quantization level multiplier QCx, block data, and null -- it judges any of run data they are, and the condition signal SS which shows the class of compressed data is

supplied to a control section 472. A control section 472 supplies control signals CTL1, CTL2, and CTL3 to the decryption section 470, a selector 474, and MPU40 (drawing 3) of video game equipment according to this condition signal SS, respectively. The base quantity child-sized table BQT and the quantization level multiplier QCx which were decrypted are supplied to the reverse quantization table creation section 250 from the decryption section 470. The DCT multiplier QF (u, v) which quantized after decode is supplied to a selector 474 from the decryption section 470.

[0105] The DC multiplier QF (0 0) registered into the zero data and DC multiplier register 476 other than the DCT multiplier QF (u, v) given from the decryption section 470 is given to the data input terminal of a selector 474. The value of the DC multiplier QF (0 0) of a pixel block of the uniform color described in the software program of a game is written in DC multiplier register 476 by MPU40 of video game equipment. In addition, the value from which the DC multiplier QF (0 0) differs to a YUV signal, respectively is registered.

[0106] two null shown in drawing 20 (C) here -- the case where the run data NRD1 and NRD2 are decrypted is considered. the 1st null -- if the data NRL of the head of the run data NRD1 are detected by the decryption section 470 -- null -- the condition signal SS which tells that it is run data is outputted to a control section 472 from the decryption section 470. A control section 472 is controlled to output control signals CTL1 and CTL2 to the decryption section 470 and MPU40 immediately according to the condition signal SS, respectively, and to stop decryption actuation. Moreover, a control section 472 supplies a control signal CTL3 to a selector 474, and makes the DC multiplier QF (0 0) registered into DC multiplier register 476 choose as a DC multiplier of the first block. As shown in drawing 20 (B), when the value of the subject-copy image data  $f(x y)$  is 12,  $QF(0 0) = 12$  are registered into DC multiplier register 476. A control section 472 controls a selector 474 as an AC multiplier of 63 pieces further to choose zero data altogether. Drawing 28 shows the DCT multiplier QF (u, v) created in this way.

[0107] the 1st null -- since the run data NRD1 show that 16 blocks of a uniform color are continuing, the DCT multiplier QF (u, v) shown in drawing 28 is created about each of 16 pixel blocks. the 2nd null -- the DCT multiplier QF (u, v) shown in drawing 28 is similarly created [ each / of two blocks ] about the run data NRD2.

[0108] null -- after processing of run data is completed, a control section 472 outputs control signals CTL1 and CTL2 to the decryption section 470 and MPU40, and it controls to resume decryption actuation. null -- the value of the luminance signal Y of a pixel block and color-difference signals U and V as which it is specified that it is a uniform color can be easily changed with run data by changing the value of DC multiplier written in DC multiplier register 476 from MPU40. if it puts in another way -- null -- if run data are used, the color of the image field of a uniform color can be changed into a desired color according to data other than compression image data. In this example, the value of the luminance signal Y of a block of a uniform color and color-difference signals U and V is specified in the software program which describes video game.

[0109] Drawing 29 is the block diagram showing other configurations of the Huffman decryption section. In this Huffman decryption section 210a, the DCT multiplier QF (u, v) outputted from the decryption section 470 bypasses a selector 474, and is directly given to the reverse quantization section 220. A selector 474 chooses one side of the DC multiplier QF (0 0) given from DC multiplier register 476, and zero data, bypasses the reverse quantization section 220, and supplies it to the IDCT section 230 directly. The control section 478 is outputting the 4th control signal CTL4 other than the three same control signals CTL1-CTL3 as drawing 27 to the IDCT section 230.

[0110] A control section 478 outputs control signals CTRL1, CTL2, and CTL4 to the decryption section 470, MPU40, and the IDCT section 230 according to the condition signal SS which shows the class of compressed data, respectively. Decode of the block data usual in the decryption section 470 supplies the decrypted DCT multiplier QF (u, v) to the reverse quantization section 220.

[0111] the 1st null shown in drawing 20 (C) -- if the data NRL of the head of the run data NDR 1 are detected by the decryption section 470 -- null -- the condition signal SS which tells that it is run data is outputted to a control section 478 from the decryption section 470. A control section 478 is controlled to

output control signals CTL1 and CTL2 to the decryption section 470 and MPU40 immediately according to the condition signal SS, respectively; and to stop decryption actuation. Moreover, a control section 478 switches the level of the control signal CTL3 given to a selector 474, and makes the DC multiplier QF (0 0) registered into DC multiplier register 476 choose. A control section 478 controls a selector 474 as an AC multiplier of 63 pieces further to choose zero data altogether. A control section 478 outputs a control signal CTL4 to coincidence to the IDCT section 230, and controls it to choose and carry out inverse transformation of the output of a selector 474.

[0112] null -- after processing of run data is completed, control signals CTL1 and CTL2 are outputted to the decryption section 470 and MPU40 from a control section 478, and it controls to resume decode actuation.

[0113] thus -- the circuit shown in drawing 29 -- null -- in the case of processing of run data, since the DCT multiplier QF (u, v) outputted from a selector 474 bypasses the reverse quantization section 220 and is directly supplied to the IDCT section 230, there is an advantage that the operation error by reverse quantization does not arise. For example, since the operation error by quantization is pressed down by min when the video encoder unit 50 (drawing 3) performs chroma-key processing which detects the part of a specific color and is made into a transparent plane color, it is possible to make the pixel block of a desired color into a transparent plane color certainly.

[0114] the above-mentioned example -- setting -- null -- although it enabled it to set only DC multiplier as arbitration with run data, it is also possible to enable it to set up a desired value to the predetermined part (for example, QF (1 0) and QF (0 1)) of AC multiplier. in this case, null -- run data express the number of a series of pixel blocks which has the same image pattern.

[0115] F. modification: -- the range which this invention is not restricted to the above-mentioned example, and does not deviate from that summary in addition -- setting -- various voice -- it is possible to set like and to carry out, for example, the following deformation is also possible.

[0116] (1) Although two-dimensional DCT conversion was taken up as orthogonal transformation in the above-mentioned example, to this invention, the orthogonal transformation (for example, K-L conversion and a Hadamard transform) of arbitration is available. Moreover, as entropy code modulation, coding (for example, algebraic-sign-izing and MEL coding) of arbitration other than Huffman coding can be used.

[0117] (2) You may realize by hardware and image data compression equipment 100 may realize image data decompression equipment 200 by software. Drawing 30 is a flow chart which shows the procedure of the processing which elongates the compressed data shown in drawing 18 with software. At step S1, the contents of compressed data are judged from the value of a header unit. When compressed data expresses the base quantity child-sized table BQT, in step S2, the base quantity child-sized table BQT is memorized in memory, and it returns to step S1.

[0118] When compressed data expresses the image, the class of initial data of each data unit contained in compressed data circles in step S3 is judged. the code data of the quantization level multiplier QCx indicated to be a data unit to drawing 18 (C) here, block data, and null -- each of run data is meant. the Huffman-coding word which expresses the quantization level multiplier QCx as initial data of a data unit as shown in drawing 18 (C), (D), (E), and (F), the Huffman-coding word of DC multiplier \*\*DC of block data, and null -- there are three kinds of the Huffman-coding words NRL showing a run.

[0119] When initial data are the Huffman-coding words of the quantization level multiplier QCx, in step S4, the quantization table QT is created by multiplying by this quantization level multiplier QCx and the base quantity child-sized table BQT. In addition, when the value of that the description mentioned above on the occasion of this multiplication, i.e., (1) multiplication result, sets the quantization level more than maximum as maximum, not quantizing (2) DC multiplier, and (3) quantization level multiplier QCx is 0, setting-as 1-all quantization level \*\* is realized. If the quantization table QT is created in step S4, it will return to step S3.

[0120] When it is judged that initial data are the Huffman-coding words of DC multiplier \*\*DC in step S3, the data for a 1-pixel block are decrypted in step S5, and the DCT multiplier QF (u, v) is obtained. Reverse quantization is performed at step S6, and two-dimensional DCT inverse transformation is

performed at step S7. If processing of steps S5-S7 is repeated about all pixel blocks included in one-set block data (drawing 18 (D)), it will return from step S8 to step S3.

[0121] step S3 -- setting -- initial data -- null -- when it is judged that it is the Huffman-coding word NRL which shows that it is run data, while DC component of a 1-pixel block is set as the value specified beforehand in step S9, all of 63 AC components are set as 0. At step S10, two-dimensional DCT inverse transformation is performed to the DCT multiplier created in this way. null -- only the block count specified by run data will return from step S11 to step S3, if step S9 and processing of S10 are repeated. In this way, image data  $ff(x, y)$  is restored by covering all the compressed data sections and performing processing of steps S3-S11.

[0122] (3) Although the above-mentioned example showed the example which applied this invention to video game equipment, this invention can be applied to all kinds of image processing system.

[0123] (4) Although the quantization level multiplier QCx and each quantization level in the quantization table QT were made into the integer in the above-mentioned example, a number including a decimal is sufficient as these.

[0124]

[Effect of the Invention] Since the 1st and 2nd quantization table on which quantization level differs can be created only by giving the 1st and the 2nd multiplier according to invention indicated to claims 1, 12, and 23 as explained above, it is effective in the ability to reduce the amount of data of compression image data compared with the case where two quantization tables are included in compression image data.

[0125] according to claims 2, 13, and 24 and invention boiled and indicated, it is effective in matching the 1st and the 2nd multiplier with the 1st and 2nd pixel block easily, respectively.

[0126] according to invention indicated to claims 3, 14, and 25 -- the 1st and the 2nd multiplier -- a meaning -- and it is effective in the ability to decode in an instant.

[0127] Since the 1st multiplier is applied to two or more pixel blocks until the 2nd multiplier appears according to invention indicated to claims 4 and 15, it is not necessary to include the 1st multiplier before each code data of a pixel block of these plurality, therefore is effective in the ability to reduce the amount of data of compression image data.

[0128] According to invention indicated to claims 5 and 16, it is effective in the ability to maintain the number of bits of quantization level below at constant value.

[0129] According to invention indicated to claims 6 and 17, it is effective in the ability to reduce the error generated in case the transform coefficient of a position is reverse-quantized.

[0130] If it is made not to change the quantization level to the dc component of a transform coefficient according to invention indicated to claims 7 and 18, it is effective in the ability to reduce the error accompanying reverse quantization of a dc component.

[0131] According to invention indicated to claims 8 and 19, it is effective in the ability of comparatively little data to express two or more continuous pixel blocks which have the same image pattern of each other.

[0132] According to invention indicated to claims 9 and 20, it is effective in the ability to apply two or more continuous pixel blocks in the same color.

[0133] According to invention indicated to claims 10 and 21, it is effective in the ability to prevent generating of the error by reverse quantization.

[0134] According to invention indicated to claim 11, it is effective in the ability of comparatively little data to express two or more continuous pixel blocks which have the color of the same pattern mutually.

[0135] According to invention indicated to claim 26, it is effective in the ability of comparatively little data to express two or more continuous pixel blocks which have the same image pattern of each other.

[0136] Since what is necessary is according to invention indicated to claim 27 just to include the 1st multiplier code in compression image data when using a new quantization table to at least one pixel block including the 1st pixel block, it is effective in the ability to reduce the amount of data of compression image data.

[0137] Since the 1st and 2nd quantization table on which quantization level differs can be created only

by giving the 1st and the 2nd multiplier according to invention indicated to claim 28, it is effective in the ability to reduce the amount of data of compression image data compared with the case where two quantization tables are included in compression image data.

[0138] Moreover, according to invention indicated to claim 29, it is effective in matching the 1st and the 2nd multiplier with the 1st and 2nd pixel block easily, respectively.

---

[Translation done.]

\* NOTICES \*

JPO and NCIP are not responsible for any  
damages caused by the use of this translation.

1. This document has been translated by computer. So the translation may not reflect the original precisely.
2. \*\*\*\* shows the word which can not be translated.
3. In the drawings, any words are not translated.

---

TECHNICAL FIELD

---

[Industrial Application] This invention relates to the approach of compressing and elongating image data, and the equipment for it.

---

[Translation done.]

\* NOTICES \*

JPO and NCIP are not responsible for any  
damages caused by the use of this translation.

1. This document has been translated by computer. So the translation may not reflect the original precisely.
2. \*\*\*\* shows the word which can not be translated.
3. In the drawings, any words are not translated.

---

PRIOR ART

---

[Description of the Prior Art] Drawing 31 is the block diagram showing the configuration of compression/expanding equipment of the conventional image data. After it carries out orthogonal transformation of the subject-copy image data for every block of a MxN pixel in the orthogonal transformation section 542, image data compression equipment 540 quantizes in the quantization section 544, further, encodes in the entropy-code-modulation section 546, and creates compression image data. On the other hand, image data decompression equipment 550 restores image data by the reverse orthogonal transformation section 552, after decrypting compression image data in the entropy decryption section 556 first and reverse-quantizing in the reverse quantization section 554. In addition, the quantization section 544 and the reverse quantization section 554 use the same quantization table 562, and use the same sign table 564 also for the entropy-code-modulation section 546 and the entropy decryption section 556.

---

[Translation done.]

\* NOTICES \*

JPO and NCIPI are not responsible for any damages caused by the use of this translation.

1. This document has been translated by computer. So the translation may not reflect the original precisely.
2. \*\*\*\* shows the word which can not be translated.
3. In the drawings, any words are not translated.

---

## EFFECT OF THE INVENTION

---

[Effect of the Invention] Since the 1st and 2nd quantization table on which quantization level differs can be created only by giving the 1st and the 2nd multiplier according to invention indicated to claims 1, 12, and 23 as explained above, it is effective in the ability to reduce the amount of data of compression image data compared with the case where two quantization tables are included in compression image data.

[0125] according to claims 2, 13, and 24 and invention boiled and indicated, it is effective in matching the 1st and the 2nd multiplier with the 1st and 2nd pixel block easily, respectively.

[0126] according to invention indicated to claims 3, 14, and 25 -- the 1st and the 2nd multiplier -- a meaning -- and it is effective in the ability to decode in an instant.

[0127] Since the 1st multiplier is applied to two or more pixel blocks until the 2nd multiplier appears according to invention indicated to claims 4 and 15, it is not necessary to include the 1st multiplier before each code data of a pixel block of these plurality, therefore is effective in the ability to reduce the amount of data of compression image data.

[0128] According to invention indicated to claims 5 and 16, it is effective in the ability to maintain the number of bits of quantization level below at constant value.

[0129] According to invention indicated to claims 6 and 17, it is effective in the ability to reduce the error generated in case the transform coefficient of a position is reverse-quantized.

[0130] If it is made not to change the quantization level to the dc component of a transform coefficient according to invention indicated to claims 7 and 18, it is effective in the ability to reduce the error accompanying reverse quantization of a dc component.

[0131] According to invention indicated to claims 8 and 19, it is effective in the ability of comparatively little data to express two or more continuous pixel blocks which have the same image pattern of each other.

[0132] According to invention indicated to claims 9 and 20, it is effective in the ability to apply two or more continuous pixel blocks in the same color.

[0133] According to invention indicated to claims 10 and 21, it is effective in the ability to prevent generating of the error by reverse quantization.

[0134] According to invention indicated to claim 11, it is effective in the ability of comparatively little data to express two or more continuous pixel blocks which have the color of the same pattern mutually.

[0135] According to invention indicated to claim 26, it is effective in the ability of comparatively little data to express two or more continuous pixel blocks which have the same image pattern of each other.

[0136] Since what is necessary is according to invention indicated to claim 27 just to include the 1st multiplier code in compression image data when using a new quantization table to at least one pixel block including the 1st pixel block, it is effective in the ability to reduce the amount of data of compression image data.

[0137] Since the 1st and 2nd quantization table on which quantization level differs can be created only by giving the 1st and the 2nd multiplier according to invention indicated to claim 28, it is effective in the ability to reduce the amount of data of compression image data compared with the case where two

quantization tables are included in compression image data.

[0138] Moreover, according to invention indicated to claim 29, it is effective in matching the 1st and the 2nd multiplier with the 1st and 2nd pixel block easily, respectively.

---

[Translation done.]

\* NOTICES \*

JPO and NCIPPI are not responsible for any damages caused by the use of this translation.

1. This document has been translated by computer. So the translation may not reflect the original precisely.
2. \*\*\*\* shows the word which can not be translated.
3. In the drawings, any words are not translated.

---

## TECHNICAL PROBLEM

---

[Problem(s) to be Solved by the Invention] By the way, the 1st part to restore by high definition in one image and the 2nd part which may be restored by low image quality may be included. In such a case, what is necessary is to quantize using the small quantization table of quantization level to the 1st part, and just to quantize using the big quantization table of quantization level to the 2nd part. A quantization table is the matrix of the same size as the pixel block of image data, i.e., the matrix of a M line N train. With the conventional image data compression / expanding equipment, when two or more quantization tables were used in one image, two or more quantization tables of a M line N train had to be transmitted to expanding equipment 550 from the compression equipment 540, and there was a problem of increasing the amount of data of compression image data.

[0004] Moreover, entropy code modulation of all the orthogonal transformation multipliers of the M line N train over those pixel blocks is carried out also with a simple image part in which the pixel block which has a uniform color continues into an image. Therefore, the compression image data showing such a simple image part also had the problem of becoming the remarkable amount of data.

[0005] This invention is made in order to solve the above-mentioned technical problem in the conventional technique, and it aims at offering the technique in which the amount of data of compression image data can be reduced compared with the former.

---

[Translation done.]

\* NOTICES \*

JPO and NCIPI are not responsible for any damages caused by the use of this translation.

1. This document has been translated by computer. So the translation may not reflect the original precisely.
2. \*\*\*\* shows the word which can not be translated.
3. In the drawings, any words are not translated.

---

## OPERATION

---

[Means for Solving the Problem and its Function] In order to solve an above-mentioned technical problem, the compression image data decompression approach by this invention (A) The code data to two or more pixel blocks, and the 1st [ to the 1st pixel block located in the head of said pixel blocks of two or more ] multiplier code, While generating the process for which the compression image data containing the 2nd [ to the 2nd pixel block of arbitration ] multiplier code is prepared, and the transform coefficient quantized by carrying out the entropy decryption of the (B) aforementioned code data the process which generates the 1st and the 2nd multiplier by carrying out the entropy decryption of said 1st and 2nd multiplier code, and (C) -- said the 1st and 2nd multiplier -- respectively -- \*\* -- by multiplying by the base quantity child-sized level of a predetermined base quantity child-sized table While reverse-quantizing the 1st, the process which creates the 2nd quantization table, and said quantized transform coefficient to the pixel block of the (D) above 1st on said 1st quantization table By reverse-quantizing said quantized transform coefficient to said 2nd pixel block on said 2nd quantization table It has the process which asks for the elongated image data by carrying out reverse orthogonal transformation of the transform coefficient by which the (E) aforementioned reverse quantization was carried out to the process which asks for the reverse-quantized transform coefficient.

[0007] Since the 1st and 2nd quantization table on which quantization level differs can be created only by giving the 1st and the 2nd multiplier, compared with the case where two quantization tables are included in compression image data, the amount of data of compression image data can be reduced.

[0008] While said code data contains two or more data units arranged according to the array sequence of a pixel block of said plurality, said 1st multiplier code is arranged just before the 1st [ to said 1st pixel block ] data unit, and said 2nd multiplier code may be made to be arranged just before the 2nd [ to said 2nd pixel block ] data unit.

[0009] If it carries out like this, the 1st and the 2nd multiplier are easily matched with the 1st and 2nd pixel block, respectively.

[0010] Moreover, as for said 1st and 2nd multiplier code, it is desirable that it is an entropy symbolic language in the sign table containing two or more entropy symbolic languages which receive the dc component of said the 1st, 2nd multiplier, and said transform coefficient.

[0011] if it carries out like this -- the 1st and the 2nd multiplier -- a meaning -- and it can decode in an instant.

[0012] As for said process (D), it is desirable to make it include the process reverse-quantized using said 1st quantization table to a series of pixel blocks from said 1st pixel block to the 3rd pixel block arranged just before said 2nd pixel block.

[0013] Since the 1st multiplier is applied to two or more pixel blocks until the 2nd multiplier appears, it is not necessary to include the 1st multiplier before each code data of a pixel block of these plurality.

[0014] When the multiplication result of said the 1st and each of the 2nd multiplier, and the quantization level contained in said base quantity child-sized table exceeds predetermined maximum, you may make it said process (C) include the process which sets up equally to said maximum the value of the quantization level for which said multiplication result exceeded said maximum.

- [0015] If it carries out like this, the number of bits of quantization level can be maintained below at constant value.
- [0016] You may make it said process (C) include the process made into the quantization level of said 1st and 2nd reverse quantization table as it is about the predetermined base quantity child-sized level in said base quantity child-sized table, without performing said multiplication substantially.
- [0017] If it carries out like this, the error generated in case the transform coefficient of a position is reverse-quantized can be reduced.
- [0018] Said predetermined base quantity child-sized level has [ making it be the quantization level about the dc component of said transform coefficient ] a desirable thing.
- [0019] If it is made not to change the quantization level to the dc component of a transform coefficient, the error accompanying reverse quantization of a dc component can be reduced.
- [0020] Said compression image data contains the same pattern block data which shows the number of two or more continuous pixel blocks which has the still more nearly same image pattern. Further said process (B) In case said quantized transform coefficient to said two or more pixel blocks expressed with said same pattern block data is created, while setting the predetermined component of said transform coefficient as the value specified beforehand You may make it have the process created by setting the value of the component of said transform coefficients other than said predetermined component as zero.
- [0021] If it carries out like this, comparatively little data can express two or more continuous pixel blocks which have the same image pattern of each other.
- [0022] As for said predetermined component, it is desirable that it is a dc component.
- [0023] If it carries out like this, two or more continuous pixel blocks can be applied in the same color.
- [0024] You may make it said process (D) include the process which omits said reverse quantization about said transform coefficient decrypted from said same pattern block data.
- [0025] If it carries out like this, generating of the error by reverse quantization can be prevented.
- [0026] Other compression image data decompression approaches by this invention (A) The process for which the compression image data containing the code data to two or more 1st pixel blocks and the same pattern block data which shows the number of two or more 2nd continuous pixel blocks which has the same image pattern is prepared, (B) While creating the transform coefficient by which the 1st to said two or more 1st pixel blocks was quantized by carrying out the entropy decryption of said code data In case the transform coefficient by which the 2nd to said two or more 2nd pixel blocks expressed with said same pattern block data was quantized is created, while setting the predetermined component of said transform coefficient as the value specified beforehand The process which creates said transform coefficient by which the 2nd was quantized by setting the value of the component of said transform coefficients other than said predetermined component as zero, (C) by reverse-quantizing using a quantization table, said transform coefficient by which the 1st and the 2nd were quantized It has the process which creates the elongated image data the 1st, the process which asks for the transform coefficient by which the 2nd was reverse-quantized, and by carrying out reverse orthogonal transformation of the transform coefficient by which the 2nd was reverse-quantized to the (D) above 1st.
- [0027] Comparatively little data can express two or more continuous pixel blocks which have the color of the same pattern mutually by this approach.
- [0028] Code data [ as opposed to two or more pixel blocks in the compression image data decompression equipment by this invention ], The 1st [ to the 1st pixel block located in the head of a pixel block of said plurality ] multiplier code, While generating the transform coefficient quantized a storage means to memorize the compression image data containing the 2nd [ to the 2nd pixel block of arbitration ] multiplier code, and by carrying out the entropy decryption of said code data An entropy decryption means to generate the 1st and the 2nd multiplier by carrying out the entropy decryption of said 1st and 2nd multiplier code, By multiplying by said the 1st and each of the 2nd multiplier, and the base quantity child-sized level of a predetermined base quantity child-sized table While reverse-quantizing said quantized transform coefficient to the 1st, a reverse quantization table creation means to create the 2nd quantization table, and said 1st pixel block, on said 1st quantization table By reverse-quantizing said

quantized transform coefficient to said 2nd pixel block on said 2nd quantization table It has a reverse quantization means to ask for the reverse-quantized transform coefficient, and a reverse orthogonal transformation means to ask for the image data elongated by carrying out reverse orthogonal transformation of said transform coefficient by which reverse quantization was carried out.

[0029] Moreover, other compression image data decompression equipments by this invention The 1st storage means which memorizes the compression image data containing the code data to two or more 1st pixel blocks, and the same pattern block data which shows the number of the 2nd continuous pixel block of plurality which has the same image pattern, The 2nd storage means which memorizes the assignment value over the predetermined component of the transform coefficient about said two or more 2nd pixel blocks, While creating the transform coefficient by which the 1st to said two or more 1st pixel blocks was quantized by carrying out the entropy decryption of said code data In case the transform coefficient by which the 2nd to said two or more 2nd pixel blocks expressed with said same pattern block data was quantized is created An entropy decryption means to create the transform coefficient by which the 2nd was quantized by setting the value of said transform coefficients other than said predetermined component as zero while setting said predetermined component as the assignment value memorized by said 2nd storage means, By reverse-quantizing using a quantization table, said transform coefficient by which the 1st and the 2nd were quantized It has a reverse orthogonal transformation means to create the elongated image data, the 1st, a reverse quantization means to ask for the transform coefficient by which the 2nd was reverse-quantized, and by carrying out reverse orthogonal transformation of the transform coefficient by which the 2nd was reverse-quantized to said 1st [ the ].

[0030] The process which the compression approach of the image data based on this invention carries out orthogonal transformation of the (A) image data for two or more pixel blocks of every in an image, and asks for a transform coefficient, (B) The 1st multiplier to the 1st pixel block located in the head of said pixel blocks of two or more, the process which specifies the 2nd multiplier to the 2nd pixel block of arbitration, and (C) -- said the 1st and 2nd multiplier -- respectively -- \*\* -- by multiplying by the base quantity child-sized level of a predetermined base quantity child-sized table The 1st, the process which creates the 2nd quantization table, and by quantizing said transform coefficient to the (D) above 1st and the 2nd pixel block on said 1st and 2nd quantization table, respectively While creating code data by carrying out entropy code modulation of the transform coefficient by which the (E) aforementioned quantization was carried out to the process which asks for the quantized transform coefficient It has the process which creates the compression image data containing the 1st, the process which creates the 2nd multiplier code, and the (F) aforementioned code data, said 1st [ the ] and the 2nd multiplier code by carrying out entropy code modulation of said the 1st and 2nd multiplier.

[0031] Since the 1st and 2nd quantization table on which quantization level differs can be created only by giving the 1st and the 2nd multiplier, compared with the case where two quantization tables are included in compression image data, the amount of data of compression image data can be reduced.

[0032] The approach of elongating other compression image data to the pan by this invention (A) The process for which the compression image data containing the code data to two or more pixel blocks and the 1st [ to the 1st pixel block of the arbitration of said two or more pixel blocks ] multiplier code is prepared, (B) While generating the transform coefficient quantized by carrying out the entropy decryption of said code data By multiplying by the process which generates the 1st multiplier by carrying out the entropy decryption of said 1st multiplier code, the 1st multiplier of (C) above, and the base quantity child-sized level of a predetermined base quantity child-sized table By reverse-quantizing the process which creates the 1st quantization table, and said quantized transform coefficient to at least one pixel block including the pixel block of the (D) above 1st on said 1st quantization table It has the process which asks for the elongated image data by carrying out reverse orthogonal transformation of the transform coefficient by which the (E) aforementioned reverse quantization was carried out to the process which asks for the reverse-quantized transform coefficient.

[0033] [0034] which can reduce the amount of data of compression image data compared with the case where two quantization tables are included in compression image data since the 1st and 2nd quantization table on which quantization level differs can be created only by giving the 1st and the 2nd multiplier, if

it carries out like this Said compression image data contains further the 2nd [ to the 2nd pixel block arranged after said 1st pixel block ] multiplier code. Said process (B) The process which generates the 2nd multiplier by carrying out the entropy decryption of said 2nd multiplier code is included. Said process (C) The process which creates the 2nd quantization table by multiplying by said 2nd multiplier and the base quantity child-sized level of said base quantity child-sized table is included. Said process (D) While reverse-quantizing using said 1st quantization table, said quantized transform coefficient to a series of pixel blocks to the 3rd pixel block arranged from said 1st pixel block just before said 2nd pixel block It is desirable to make it include the process which reverse-quantizes said quantized transform coefficient to at least one pixel block including said 2nd pixel block on said 2nd quantization table. [0035] If it carries out like this, it can quantize by using two quantization tables at the time of expanding of compression image data only by including the 1st and 2nd multiplier code in compression image data.

---

[Translation done.]

\* NOTICES \*

JPO and NCIP are not responsible for any damages caused by the use of this translation.

1. This document has been translated by computer. So the translation may not reflect the original precisely.
2. \*\*\*\* shows the word which can not be translated.
3. In the drawings, any words are not translated.

---

EXAMPLE

---

[Example] Below, sequential explanation is performed about each following item.

- A. The whole compression/expanding equipment configuration and basic actuation;
- B. Adjustment of the quantization table QT by the quantization level multiplier QCx;
- C. Huffman coding and configuration of compressed data;
- D. Detail configuration of the reverse quantization table creation section;
- E. null -- a decryption [0037] of run data A. The configuration and actuation of compression/expanding equipment : drawing 1 is the block diagram showing the function of the compression equipment 100 of image data, and expanding equipment 200 which applied one example of this invention.

[0038] Image data compression equipment 100 is equipped with the DCT section 110 which performs discrete cosine conversion to the subject-copy image data  $f(x, y)$ , the quantization section 120 which quantizes the transform coefficient  $F(u, v)$  obtained by DCT conversion, the Huffman coding section 130 which carries out Huffman coding of the quantized transform coefficient  $QF(u, v)$ , and creates the compression image data ZZ, the quantization table creation section 140, and the Huffman-coding table memory 150. The quantization table creation section 140 creates the quantization table QT based on the base quantity child-sized table BQT and the quantization level multiplier QCx so that it may mention later. The compression image data ZZ is memorized by storages, such as CD-ROM, and is supplied to image data decompression equipment 200 from image data compression equipment 100.

[0039] Image data decompression equipment 200 is equipped with the Huffman decryption section 210 which carries out the Huffman decryption of the compression image data ZZ, the reverse quantization section 220 which reverse-quantizes the transform coefficient  $QF(u, v)$  after the decoded quantization, the IDCT section 230 which performs discrete cosine inverse transformation to the reverse-quantized transform coefficient  $FF(u, v)$ , and obtains image data  $ff(x, y)$ , the Huffman-coding table memory 240, and the reverse quantization table creation section 250. The reverse quantization table creation section 250 creates the quantization table QT for the base quantity child-sized table BQT and the quantization level multiplier QCx which were decoded from the compression image data ZZ based on a receipt and these from the Huffman decryption section 210. This quantization table QT is the same as the quantization table QT used with the compression equipment 100. Moreover, the Huffman-coding table HT memorized by the Huffman-coding table memory 240 is the same as what is memorized by the Huffman-coding table memory 150 of a compression equipment 100.

[0040] Drawing 2 is the block diagram showing the concrete configuration of image data compression equipment 100. This image data compression equipment 100 is equipped with CPU101, main memory 102, a keyboard 103, a mouse 104, a magnetic disk drive 105, and optical-magnetic disc equipment 106. Each processing sections 110-140 of the image data compression equipment 100 shown in drawing 1 are realized by the software program memorized by main memory 102. Moreover, the base quantity child-sized table BQT and the Huffman-coding table HT are memorized by the magnetic disk drive 105. This image data compression equipment 100 is a workstation for creating video game, and various kinds of programs for creating the video game other than the software program for image data compression are performed by CPU101. The completed game program is stored in optical-magnetic disc equipment 106

with compression image data. And CD-ROM containing the program and compression image data of video game is manufactured using this magneto-optic disk.

[0041] Drawing 3 is the block diagram showing the configuration of the video game equipment 20 containing image data decompression equipment 200. The microprocessor 40 which manages in generalization all processings about others, an image processing, and this (henceforth MPU), [ CD-ROM drive / 32 / to which this video game equipment 20 was connected through the SCSI bus 36 ] The main memory 41 by which direct continuation was carried out to this MPU40 (it is hereafter called M-RAM), The various units connected to the bus (M-BUS) 43 of ROM42 and MPU40 which similarly memorized the BIOS program, That is, it has the picture signal control unit 45, the image data decompression unit 200, the VDP unit 49 that outputs a specific picture signal, the video encoder unit 50 which performs composition and the output of a video signal, and the voice data output unit 52 treating voice data.

[0042] Moreover, the memory connected to the local bus (K-BUS) 54 of the picture signal control unit 45 into this video game equipment 20 The memory connected to the local bus of 55 and the image data decompression unit 200 (It is hereafter called K-RAM) The video memory connected to the local bus of 251 and the VDP unit 49 (It is hereafter called R-RAM) (It is hereafter referred to as V=RAM) The output signal from 59 and the video encoder unit 50 is changed into the usual video signal (NTSC), and it has the NTSC converter 60 outputted to color television 28.

[0043] The picture signal control unit 45, the image data decompression unit 200, the video encoder unit 50, and the voice data output unit 52 are constituted by the logical circuit, respectively.

[0044] Drawing 4 (A) is the top view showing an example of the subject-copy image used as the background image of a game. This subject-copy image is an image with which volcanic natural drawing was inserted in the background BG applied in the uniform color. Drawing 4 (B) expands and shows some subject-copy images including one pixel block PB. Generally, the pixel block PB can be set up so that the pixel PX of a MxN individual may be included. As a value of two integers M and N, 8 or 16 are desirable and is M=N=8 in this example. In addition, integers M and N may be set as a different value. The amount of [ which expresses Background BG in the compression image data ZZ ] data division have a special data format (null run data) showing that the pixel block PB of a uniform color continues so that it may mention later.

[0045] The DCT section 110 of image data compression equipment 100 performs two-dimensional DCT conversion for every pixel block PB according to the following formula 1.

[Equation 1]

$$F(u, v) = \frac{1}{4} C(u)C(v) \sum_{x=0}^7 \sum_{y=0}^7 f(x, y) \cos \frac{\pi(2x+1)u}{16} \cos \frac{\pi(2y+1)v}{16}$$

$$C(u), C(v) = \frac{1}{\sqrt{2}} \quad (u, v = 0 \text{ の時})$$

$$= 1 \quad (u, v \neq 0 \text{ の時})$$

Here, the array of a transform coefficient, and u and v of the coordinate and F (u, v) the array of 8x8 image data by which f (x y) is contained in one pixel block PB, x, and y indicate the location of each pixel within each pixel block PB to be are the coordinates of frequency space.

[0046] Drawing 5 is the explanatory view showing the array of a transform coefficient F (u, v). A transform coefficient F (u, v) is the same array of 8x8 as the pixel block PB. The transform coefficient F of an upper left edge (0 0) is called DC component (or DC multiplier), and other transform coefficients are called AC component (or AC multiplier). DC component shows the average of the image data in the pixel block PB. Moreover, AC component shows change of the image data within the pixel block PB. Since the image data of the adjoining pixel has a certain amount of correlation, the value of a low-frequency component is comparatively large in AC multiplier, and the value of a high frequency component is comparatively small. Moreover, the effect which it has on image quality has a comparatively small high frequency component.

[0047] Drawing 6 is the explanatory view showing basic actuation of image data compression equipment 100 and image data decompression equipment 200. The DCT section 110 creates the DCT multiplier F (u, v) shown in drawing 6 (a).

[0048] The quantization table creation section 140 creates the quantization table QT ( drawing 6 (d)) by multiplying by the base quantity child-sized table BQT ( drawing 6 (c)) and the quantization level multiplier QCx, as shown in the following formula 2.

[Equation 2]

$$QT(u, v) = QCx \times BQT(u, v)$$

[0049] Since it is QCx=1 in the example of drawing 6 , the quantization table QT is the same as the base quantity child-sized table BQT.

[0050] The quantization section 120 asks for the quantized DCT multiplier QF (u, v) which is shown in drawing 6 (b) by carrying out linear quantization of the DCT multiplier F (u, v) on the quantization table QT. Linear quantization is processing which does a division and rounds off the division result for an integer.

[0051] The Huffman coding section 130 creates the compression image data ZZ ( drawing 6 (e)) by carrying out Huffman coding of this DCT multiplier QF (u, v). In addition, about the approach of Huffman coding, it mentions later further. The compression image data ZZ contains the 1st data showing the base quantity child-sized table BQT, and the 2nd data showing the quantization level multiplier QCx and a transform coefficient QF (u, v) so that it may mention later.

[0052] If the compression image data ZZ is given to image data decompression equipment 200, the Huffman decryption section 210 will decrypt the compression image data ZZ, and it will ask for the DCT multiplier QF (u, v) ( drawing 6 (f)). Since Huffman coding is reversible coding, this DCT multiplier QF (u, v) is the same as the DCT multiplier QF (u, v) ( drawing 6 (b)) after the quantization called for by the quantization section 120 of image data compression equipment 100. In addition, the Huffman decryption section 210 also decrypts the base quantity child-sized table BQT ( drawing 6 (c)) and the quantization level multiplier QCx which are contained in the compression image data ZZ other than the DCT multiplier QF (u, v), and is given to the reverse quantization table creation section 250.

[0053] The reverse quantization table creation section 250 creates the quantization table QT ( drawing 6 (d)) by carrying out the multiplication of the base quantity child-sized table BQT and the quantization level multiplier QCx. The reverse quantization section 220 carries out the multiplication of this quantization table QT and the DCT multiplier QF (u, v), and asks for the decoded DCT multiplier FF (u, v) which is shown in drawing 6 (g).

[0054] The IDCT section 230 performs two-dimensional DCT inverse transformation shown in the following formula 3 to this DCT multiplier FF (u, v), and creates the restored image data ff (x y).

[Equation 3]

$$ff(x, y) = \frac{1}{4} \sum_{u=0}^7 \sum_{v=0}^7 C(u)C(v) FF(u, v) \cos \frac{\pi(2x+1)u}{16} \cos \frac{\pi(2y+1)v}{16}$$

$$C(u), C(v) = \frac{1}{\sqrt{2}} \quad (u, v = 0 の時)$$

$$= 1 \quad (u, v \neq 0 の時)$$

[0055] B. Adjustment of the quantization table QT by the quantization level multiplier QCx : since it is created by carrying out the multiplication of the base quantity child-sized table BQT and the quantization level multiplier QCx according to said formula 2, the quantization table QT can enlarge each quantization level in the quantization table QT, if the value of the quantization level multiplier QCx is enlarged. In case the value of the quantization level multiplier QCx compresses image data in image data compression equipment 100, an operator chooses it from two or more values (0-15) defined beforehand.

[0056] Drawing 7 is the explanatory view showing compression/expanding actuation at the time of specifying the quantization level multiplier QCx as 4. The quantization table creation section 140 and the reverse quantization table creation section 250 create the quantization table QT shown in drawing 7 (d) according to the above-mentioned formula 2. However, in this example, the maximum of quantization level is restricted to 15 and all the values of the quantization level from which the result of multiplication becomes 15 or more are compulsorily set as 15.

[0057] If linear quantization of the DCT multiplier F (u, v) of drawing 7 (a) is carried out using the quantization table QT shown in drawing 7 (d), as shown in drawing 7 (b), DC component will be 1 and the DCT multiplier QF (u, v) all whose AC components are 0 will be obtained. Thus, if the value of the quantization level multiplier QCx is enlarged, since the number of 0 in the quantized DCT multiplier QF (u, v) will increase, a data compression rate can be raised. However, since the decoded DCT multiplier FF (u, v) which is shown in drawing 7 (g) shows the DCT multiplier F (u, v) and a value which becomes and is different of the origin shown in drawing 7 (a), it is larger than the case ( drawing 6 ) where the quantization level multiplier QCx is equal to 1. [ of degradation of image quality ]

[0058] By the way, since DC component of a DCT multiplier shows the average of the image data within the pixel block PB, its effect to image quality is quite large. Therefore, it is desirable to keep the quantization level for DC components the same as the value in the base quantity child-sized table BQT irrespective of the value of the quantization level multiplier QCx. Drawing 8 is the explanatory view showing the processing in such a case, and the quantization level QT (0 0) for DC components is kept compulsory at 1. In the case of QT(0 0)=1, since DC component of the decoded DCT multiplier FF (u, v) which is shown in drawing 8 (g) is maintained at the same value as DC component of the original DCT multiplier F (u, v) shown in drawing 8 (a), degradation of image quality can be suppressed comparatively small with compressibility comparable as the case of drawing 7 . In addition, it is not necessary to necessarily set quantization level QT (0 0) for DC components to 1, and it may set up other any value.

[0059] It is also possible to choose 0 as a quantization level multiplier QCx. As shown at drawing 9 in the case of QCx=0, all the quantization level in the quantization table QT is set as 1. Since the quantized DCT multiplier QF (u, v) is the same as the original DCT multiplier F (u, v), although compressibility is small, compression/expanding can be performed by high definition.

[0060] If the above is summarized, the quantization table creation section 140 and the reverse quantization table creation section 250 have the following descriptions.

- (1) Create the quantization table QT by carrying out the multiplication of the base quantity child-sized table BQT and the quantization level multiplier QCx according to a formula 2.
- (2) The result of multiplication sets up equally to maximum the thing more than the maximum (= 15) of quantization level ( drawing 6 - drawing 8 ).
- (3) Keep the quantization level QT (0 0) for DC components the same as that of the value in the base quantity child-sized table BQT irrespective of the value of the quantization level multiplier QCx ( drawing 8 ).
- (4) When the quantization level multiplier QCx is 0, set all quantization level as 1 ( drawing 9 ).

[0061] The above-mentioned operation which asks for the quantization table QT is performed by the software program in image data compression equipment 100, and it is carried out by the reverse quantization table creation section 250 which consisted of hardware of dedication in image data decompression equipment 200. About the concrete circuitry of the reverse quantization table creation section 250, it mentions later further.

[0062] In addition, in case the quantization level multiplier QCx compresses image data with image data compression equipment 100, it can specify a value which uses a keyboard 103 and a mouse 104 and is different for every pixel block PB.

[0063] C. Huffman coding and the configuration of compressed data; the Huffman coding section 130 ( drawing 1 ) of image data compression equipment 100 consists of the DC multiplier coding section and the AC multiplier coding section. Drawing 10 (A) is the block diagram showing the function of DC multiplier coding section. the block delay section 131 and an adder 132 are shown in drawing 10 (B) --

as -- DC multiplier DCi of each pixel block PB DC multiplier DCi-1 of the pixel block PB before one difference -- \*\*DC is computed.

[0064] the category-ized table showing the categorizing processing section 133 in drawing 11 -- following -- the difference of DC multiplier -- it asks for Category SSSS and the discernment data ID corresponding to \*\*DC. Category SSSS -- the difference of DC multiplier -- it is the number which shows the range of \*\*DC. two or more difference as which the discernment data ID are specified by Category SSSS -- it is data in which it is shown from the smaller one in \*\*DC the value of what position it is.

[0065] Category SSSS is changed into the Huffman-coding word HFDC for DC multipliers in the 1 more-dimensional Huffman coding section 134 ( drawing 10 ). Drawing 12 is the explanatory view showing an example of the Huffman-coding table HTDC used by the 1-dimensional Huffman coding section 134. In this example, the subject-copy image data  $f(x, y)$  shall be expressed by the YUV signal (a luminance signal Y and two color-difference signals U and V). The Huffman-coding table for DC multipliers of U signal / V signal common use only contains the symbolic language of the category SSSS of 0-9. On the other hand, the Huffman-coding table for DC multipliers for Y signals contains the symbolic language of the category SSSS of 15-31 other than the symbolic language of the category SSSS of 0-9. the null which mentions the Huffman-coding word of SSSS=15 later -- it is shown that it is run data. null -- run data are data in which it is shown that the pixel block PB of a uniform color continues. Moreover, the Huffman-coding word of SSSS=16-31 is a sign which shows the value of the quantization level multiplier QCx. For example, the Huffman-coding word "111110000" over SSSS=16 shows QCx=0, and the Huffman-coding word "111111111" over SSSS=31 shows QCx=15. In addition, meaning decode is possible for the Huffman-coding word of drawing 12 about all of category SSSS=1-9, and 15-31, and instant decode is possible for it.

[0066] Drawing 13 is the block diagram showing the function of AC multiplier coding section in the Huffman coding section 130. List direct [ of the array F of AC multiplier (u, v) (except for  $u=v=0$ ) ] is first carried out to one dimension by the zigzag scan section 135. Drawing 14 is the explanatory view showing the usual route of a zigzag scan.

[0067] The value of AC multiplier by which list direct [ of the judgment section 136 ] was carried out to one dimension judges whether it is 0. If the value of AC multiplier is 0, the run length counter 137 will change AC multiplier of continuous 0 into the zero run length NNNN. If AC multiplier is not 0, the value of the AC multiplier will be changed into Category SSSS and the discernment data ID by the categorizing section 138. Under the present circumstances, the category-ized table shown in drawing 11 is referred to.

[0068] The zero run length NNNN and Category SSSS are changed into the Huffman-coding word HFAC for AC multipliers in the two-dimensional Huffman coding section 139. Drawing 15 is the explanatory view showing the two-dimensional Huffman-coding table HTAC for AC multipliers. Moreover, drawing 16 shows an example of the Huffman-coding word of the parts (topmost part of two lines in drawing 15 ) of NNNN=0 and NNNN=1 in the Huffman-coding table HTAC. In addition, the Huffman-coding word "11111" of NNNN/SSSS=0/0 shows termination of the code data to one pixel block.

[0069] Drawing 17 is the explanatory view showing an example of Huffman coding. Drawing 17 (B) shows coding of DC multiplier. When the value of DC multiplier in the pixel block before one is assumed to be 0, it is  $**DC=F(0, 0)=12$ . According to the category-ized table of drawing 11 , the category SSSS of  $**DC=12$  is 4 and the discernment data ID are "1100." Moreover, according to the Huffman-coding table for DC multipliers of drawing 12 , the Huffman-coding word HFDC of category SSSS=4 is "011." In addition, the Huffman-coding table for Y signals is used here. Huffman coding (HF+ID) to DC multiplier is set to "0111100" as shown in drawing 17 (B).

[0070] Drawing 17 (C) shows coding of AC multiplier. First, AC multiplier is put in order by the array of a single dimension with a zigzag scan. This array is changed into the category SSSS of the value which are not the zero run length NNNN and zero (refer to drawing 11 ). The combination of the zero run length NNNN and Category SSSS is changed into the Huffman-coding word HFAC by the

Huffman-coding table for AC multipliers shown in drawing 15 and drawing 16, and it is combined with the discernment data ID of AC multiplier which is not zero, and as shown in drawing 17 (C), Huffman coding (HFAC+ID) is created.

[0071] Drawing 18 is the explanatory view showing the configuration of compressed data. The whole compressed data consists of a header unit, the compressed data section, and the dummy section, as shown in drawing 18 (A). The header unit has 1 byte of four data DFH, DFL, DLH, and DLL, respectively. The first two data DFH and DFL show the class of data contained in the compressed data section. There is a class of the data of the base quantity child-sized table BQT, full color natural picture compression data, run length picture compression data, etc. of the data of the compressed data section. Data (DLH+DLL) of 16 bits of posterior parts of a header show the data length of the sum total of the compressed data section and the dummy section. Since the compressed data section is variable-length data containing Huffman coding, the data length of the sum total with the dummy section is adjusted so that it may become the die length of the integral multiple of WORD (= 2 bytes).

[0072] Drawing 18 (B) shows the configuration of the compressed data showing the base quantity child-sized table BQT. This one-set compressed data contains the data showing the base quantity child-sized table BQT for Y signals, and the data showing the base quantity child-sized table BQT of U signal / V signal common use. In addition, it is not necessary to carry out Huffman coding of the data showing the base quantity child-sized table BQT.

[0073] Drawing 18 (C) shows the configuration of the compressed data of a full color natural image. the code data (symbolic language of category SSSS=16-31 in drawing 12) which expresses the quantization level multiplier QCx to the compressed data section, the block data which is code data of each pixel block, and the null which shows two or more pixel blocks of a uniform color -- run data are included.

[0074] As shown in drawing 18 (D), the block data of one unit consists of 4 sets of data for Y signals, 1 set of U signal data, and 1 set of data for V signals. Drawing 19 is the explanatory view showing the relation of a block of each signal of YUV. As shown in drawing 19 (A), one screen in this example has the magnitude of the 256 pixel x240 scanning line. About a Y signal, DCT conversion is performed for every 8x8-pixel pixel block, without culling out. On the other hand, about U signal and V signal, as shown in drawing 19 (B), it is thinned out and (subsampling) made one half in a longitudinal direction and a lengthwise direction, and DCT conversion is performed to the 8x8-pixel block after thinning out. Therefore, as shown in drawing 19 (C), the field of four pixel blocks Y1-Y4 of a Y signal is equivalent to the field of one pixel block of U signal and V signal. In addition, it is because it is comparatively insensible to change (change of U signal and V signal) of a color to operate U signal and V signal on a curtailed schedule, without operating a Y signal on a curtailed schedule although human being's eyes are comparatively sensitive to change (change of a Y signal) of brightness. Compressibility can be raised by thinning out only U signal and V signal, without degrading image quality too much. In addition, the block data of one unit shown in drawing 18 (D) arranges in order the Huffman-coding data of each field shown in drawing 19 (C).

[0075] The code data to one pixel block in block data consists of one Huffman-coding data of DC multiplier, and two or more Huffman-coding data of AC multiplier, as shown in drawing 18 (F). The Huffman-coding data of DC multiplier consist of a Huffman-coding word HFDC of Category SSSS, and discernment data ID, as mentioned above (drawing 18 (G)). Moreover, the Huffman-coding data of AC multiplier consist of the Huffman-coding words HFAC and the discernment data ID to combination with Category SSSS with the zero run length NNNN (drawing 18 (H)).

[0076] The code data of the quantization level multiplier QCx is inserted just before the block data of the pixel block which wants to change the head of the compressed data section, and the value of the quantization level multiplier QCx. To two or more pixel blocks before the 2nd quantization level multiplier QCx is inserted, the top quantization level multiplier QCx is used in common. Moreover, to two or more blocks before the 3rd quantization level multiplier QCx (not shown) is inserted, the 2nd quantization level multiplier QCx is used in common.

[0077] In addition, when the symbolic language of the quantization level multiplier QCx is not contained at the head of the compressed data section, it is considered that it is QCx=1. Therefore, also

when the quantization level multiplier QCx is not inserted in the head of the compressed data section but the quantization level multiplier QCx is inserted only once on the way, it is equivalent to two quantization level multipliers QCx being specified.

[0078] Since Huffman coding showing the quantization level multiplier QCx is inserted between block data, it can apply this new quantization level multiplier QCx easily to the next block data at the time of the new quantization level multiplier QCx being decrypted. Moreover, since the code data of the quantization level multiplier QCx is expressed with the Huffman-coding word for DC multipliers as shown in drawing 12, even if this is inserted between block data, it is possible to judge immediately whether this code data is code data of DC multiplier for block Y1 or it is the code data of the quantization level multiplier QCx.

[0079] the null contained in the compressed data section -- run data are shown in drawing 18 (E) -- as -- null -- it consists of discernment data ID with the symbolic language for DC multipliers "NRL" which shows that it is run data, and the block count.

[0080] drawing 20 -- null -- it is the explanatory view showing the image expressed by run data. The background BG of the subject-copy image of drawing 20 (A) is applied in the uniform color. The part of the ellipse of drawing 20 (A) is assumed to be what 18 pixel blocks which have the image data value (f(x,y)=12) with all the same pixels as shown in drawing 20 (B) are following. the null to which drawing 20 (C) expresses these pixel blocks -- run data are shown. this null -- run data -- the 1st null for a 16-pixel block -- the run data NRD1 and the 2nd null for a 2-pixel block -- the run data NRD2 are included.

[0081] each -- null -- the head of the run data NRD1 and NRD2 -- null -- it has the symbolic language for DC multipliers "NRL" (symbolic language of category SSSS=15 of drawing 12 "1111011") which shows that it is run data. decrypting the symbolic language for DC multipliers in a head, since Huffman coding of DC multiplier is arranged at the head of the usual block data as shown in drawing 18 (F) -- null -- run data, block data, and the code data of the quantization level multiplier QCx -- a meaning -- and it is discriminable in an instant.

[0082] The block count is expressed with the Huffman-coding word for AC multipliers as shown in drawing 20 (C). drawing 21 -- the inside of the Huffman-coding table for AC multipliers (drawing 15) -- null -- it is drawing showing the part used for run data. null -- when used for run data, if the zero run length NNNN is equal to ([block count]-1), he will be set up. Moreover, the Huffman-coding word of category SSSS=1 is used noting that the value of AC multiplier is 1. the 1st null shown in drawing 20 (C) -- the data (NNNN/SSSS=15/1) of the block count in the run data-NRD1 show that 16 pixel blocks of a uniform color are continuing. moreover, the 2nd null -- the data (NNNN/SSSS=1/1) of the block count in the run data NRD2 show that two pixel blocks of a uniform color are continuing.

[0083] each -- null -- the discernment data ID are added to the back end of the run data NRD1 and NRD2. It is fixed to ID=1 in this example.

[0084] null -- run data can express that two or more pixel blocks which continued with about 20-bit data are uniform colors in this way for the usual block data's, expressing the one-set block (the Y signal shown in drawing 19 -- a 4-pixel block, U signal, and V signal -- an every 1pixelblock -- it contains) of a uniform color on the other hand -- about 300- about 400 bits is required. And also when it is shown that a two or more sets pixel block is a uniform color, it is [ about 300 - 400 bits of abbreviation ] required about each set. therefore, null -- if run data are used, it is possible to reduce considerably the amount of data of the compressed data showing the pixel block of a majority of continuous uniform colors.

[0085] in addition, null -- the value of the luminance signal Y of a pixel block of a uniform color and color-difference signals U and V expressed by run data is not included in compressed data, but is specified in the software program which describes video game. He specifies the brightness and color tone of these blocks using a keyboard 103 or a mouse 104 while he specifies the range of the field (drawing 20 (A) the background BG) of a pixel block of a uniform color with a mouse 104, in case an operator creates the software program for video game. When carrying out like this and a specific event occurs while having performed the game, for example using video game equipment 20 (drawing 3), the special visual effectiveness of changing the color of Background BG in time can be produced. in

addition, null -- about the concrete circuitry which decrypts run data, it mentions later further.

[0086] D. The detail configuration of the reverse quantization table creation section : drawing 22 is the block diagram showing the internal configuration of the reverse quantization table creation section 250 shown in drawing 1. The reverse quantization table creation section 250 is equipped with the multiplication unit 254 which carries out the multiplication of the address-generation circuit 252 which generates the address of RAM251 and RAM251 which memorizes the base quantity child-sized table BQT, the latch circuit 253 holding the quantization level multiplier QCx, and the quantization level multiplier QCx and the base quantity child-sized table BQT, and generates the quantization table QT. The quantization table QT which was caused multiplication unit 254 and created is supplied to the reverse quantization section 220.

[0087] First, if the compression image data ZZ contained by CD-ROM is given to the Huffman decryption section 210, the code data of the base quantity child-sized table BQT will be decrypted first, and will be supplied to RAM251. This base quantity child-sized table BQT is memorized by RAM251 according to the light address given from the address-generation circuit 252. The base quantity child-sized table BQT memorized by RAM251 is used to all pixel blocks.

[0088] The address-generation circuit 252 generates the lead address synchronizing with the DCT multiplier data QF (u, v) outputted from the Huffman decryption section 210, and the base quantity child-sized table BQT is read from RAM251 according to this lead address. On the other hand, the quantization level multiplier QCx decrypted in the Huffman decryption section 210 is latched by the latch circuit 253, and it is saved at a latch circuit 253 until the following quantization level multiplier QCx is given. Therefore, the same quantization level multiplier QCx is used in common to two or more pixel blocks until the quantization level multiplier QCx is newly supplied.

[0089] Drawing 23 is the block diagram showing the internal configuration of the latch circuit 253 contained in the reverse quantization table creation section 250 (drawing 22), and the multiplication unit 254. The latch circuit 253 consists of two latches 402, 404. The multiplication unit 254 has the synchronous-clock creation circuit 412, AND circuit 414, U signal start detector 416, V signal start detector 418, NAND circuit 420, the selector 422, the multiplier 424, the clipping circuit 426, and the zero value correction circuit 428.

[0090] Drawing 24 is a timing chart which shows actuation of the circuit shown in drawing 23. If the quantization level multiplier QCx is decrypted in the Huffman decryption section 210 (drawing 1), the data of the quantization level multiplier QCx will be given to a latch circuit 253 with an enable signal QEN (drawing 24 (a), (b)). The 1st latch 402 latches the quantization level multiplier QCx by the rising edge of an enable signal QEN, and supplies an output Q1 to the 2nd latch 404 (drawing 24 (c)).

[0091] As shown in drawing 24 (d), after the base quantity child-sized table BQT for Y signals is read 4 times, the base quantity child-sized table BQT of U signal / V signal common use is read twice from RAM251 (drawing 22).

[0092] Synchronizing with the base quantity child-sized table BQT read from RAM251, it is set to L level in the period of a Y signal, and block recognition signal UV/Y used as H level is given to the synchronous-clock creation circuit 412 (drawing 23) from the address-generation circuit 252 in the period of U signal and V signal (drawing 24 (e)). Moreover, an enable signal EN is also given to the synchronous-clock creation circuit 412 in this case. The synchronous-clock creation circuit 412 reverses block recognition signal UV/Y, generates a synchronizing clock signal SCK (drawing 24 (f)), and supplies this to the 2nd latch's 404 clock input terminal. The 2nd latch 404 latches the 1st latch's 402 output Q1 by the rising edge of this synchronizing clock signal SCK, and supplies that output Q2 (drawing 24 (g)) to the data input terminal of a selector 422. In addition, the fixed value "1" is given to other data input terminals of a selector 422.

[0093] U signal start detector 416 generates U start signal USTRT (drawing 24 (h)) which shows the start time of day of the base quantity child-sized table BQT for U signals based on the block recognition signals UV/Y and the 10MHz basic clock signal CLK. This U start signal USTRT is a signal set to L level from the rising edge of block recognition signal UV/Y only for 100 nanoseconds.

[0094] Six blocks Y1-Y4, and the block change signal SWTCH (drawing 24 (i)) and the block

recognition signals UV/Y with which level changes by turns in each period of U and V are given to AND circuit 414. V signal start detector 418 generates V start signal VSTRT (drawing 24 (j)) which shows the start time of day of the base quantity child-sized table BQT for V signals based on the output of AND circuit 414, and the 10MHz basic clock signal CLK. Block recognition signal UV/Y of this V start signal VSTRT is the signal set to L level from the rising edge of the block change signal SWTCH only for 100 nanoseconds in the period of H level.

[0095] Drawing 25 is the block diagram showing the internal configuration of U signal start detector 416 and V signal start detector 418. These circuits 416,418 are constituted from a D flip-flop and a NAND circuit by both. The block recognition signals UV/Y are supplied to D input terminal of D flip-flop 432 of U signal start detector 416, and the 10MHz basic clock signal CLK is inputted into the clock input terminal. Block recognition signal UV/Y and the reversal output of D flip-flop 432 are given to the input terminal of NAND circuit 434. In addition, block recognition signal UV/Y synchronizes with the basic clock signal CLK.

[0096] Since the block recognition signals UV/Y is [ the reversal output of D flip-flop 432 ] H level between L level, the output USTRT of NAND circuit 434 is maintained at H level (refer to drawing 24 (h)). If the output (U start signal USTRT) of NAND circuit 434 is set to L level and D flip-flop 432 latches an input with the edge of the basic clock signal CLK after 100 nanoseconds immediately after block recognition signal UV/Y changes from L level to H level, the output USTRT of NAND circuit 434 will return to H level again.

[0097] Actuation of V signal start detector 418 is the same as actuation of U signal start detector 416. However, since the AND of block recognition signal UV/Y and the block change signal SWTCH is given to D input terminal of D flip-flop 436, in the period of H level, V start signal VSTRT is set to L level from the rising edge of the block change signal SWTCH by block recognition signal UV/Y only for 100 nanoseconds.

[0098] U start signal USTRT and V start signal VSTRT are inputted into NAND circuit 420 (drawing 23), and the output (selection signal SEL) is supplied to the selection input terminal of a selector 422. A selection signal SEL (drawing 24 (k)) is set to H level only for 100 nanoseconds at the beginning of the period the start of the period for U signals, and for V signals. A selector 422 outputs the output Q2 given from the latch circuit 253 as it is, when a selection signal SEL is L level, and on the other hand, when a selection signal SEL is H level, it outputs a fixed value "1." The multiplication of the output Q3 of a selector 422 is carried out to the base quantity child-sized table BQT by the multiplier 424.

[0099] As shown in drawing 24 (l), the output Q3 of a selector 422 is surely set to "1" irrespective of the value of the quantization level multiplier QCx at the beginning of the period the start of the period for U signals, and for V signals. The period for 100 nanoseconds which the period the object for U signals and for V signals begins is a period for computing the quantization level for DC multipliers. Therefore, if the circuit of drawing 23 is used, it can avoid performing substantially the multiplication of the quantization level for DC multipliers used for U signal and V signal, and the specified quantization level multiplier QCx. If it puts in another way, the circuit shown in drawing 23 will be a circuit which realizes the operation shown in drawing 8 (c) and (d).

[0100] As shown in drawing 23, the output of a multiplier 424 is corrected by a clipping circuit 426 and the zero value correction circuit 428, and serves as the final quantization table QT. Drawing 26 is the block diagram showing the internal configuration of these two circuits 426,428.

[0101] The clipping circuit 426 consists of 4 input OR circuit 450 and eight 2 input OR circuits 452. These circuits are circuits in the case of expressing quantization level by 9 bits (the most significant bit being a sign bit). D9-D12 are inputted into 4 input OR circuit 450 4 bits of high orders except a sign bit D13 among the 14-bit data outputted from the multiplier 424 (drawing 23). The output of 4 input OR circuit 450 is given to one input terminal of eight 2 input OR circuits 452, and D1-D8 are given to the input terminal of another side 8 bits of low order of the output of a multiplier 424. When at least one value of D9-D12 is "1" 4 bits of high orders; all of the output of eight 2 input OR circuits 452 are set to "1." Therefore, the output of a clipping circuit 426 is set to the case of 255 or more for the output of a multiplier 424 by 255 with a decimal number.

[0102] In the zero value correction circuit 428, the output of seven 2 input OR circuits 452 seven bits D2 in a clipping circuit 426 - for D8 is outputted as it is. Moreover, the output and sign bit D13 of these seven 2 input OR circuits 452 are given to eight inverters 460, respectively. The output of eight inverters 460 is given to 8 input AND circuit 462, and the output of this AND circuit 462 is supplied to 2 input OR circuit 464. The output of 2 input OR circuit 452 for least significant bit D1 is given to this 2 input OR circuit 464. Consequently, when [ of the output of a multiplier 424 ] all of 13 bits of values of D1-D13 are "0", as for the zero value correction circuit 428, the value of 8 bits of others [ value / of the least significant bit ] outputs the quantization level QT of "0" by "1." If it puts in another way, the zero value correction circuit 420 will have realized the operation shown in drawing 9 (c) and (d).

[0103] E. null -- decryption [ of run data ]: -- drawing 27 is the block diagram showing the internal configuration of the Huffman decryption section 210 (drawing 1) in image data decompression equipment 200. The Huffman decryption section 210 is equipped with the decryption section 470 which carries out the Huffman decryption of the compression image data ZZ, the control section 472, the selector 474, and DC multiplier register 476.

[0104] the class of compressed data with which the decryption section 470 was given -- the base quantity child-sized table BQT, the quantization level multiplier QCx, block data, and null -- it judges any of run data they are, and the condition signal SS which shows the class of compressed data is supplied to a control section 472. A control section 472 supplies control signals CTL1, CTL2, and CTL3 to the decryption section 470, a selector 474, and MPU40 (drawing 3) of video game equipment according to this condition signal SS, respectively. The base quantity child-sized table BQT and the quantization level multiplier QCx which were decrypted are supplied to the reverse quantization table creation section 250 from the decryption section 470. The DCT multiplier QF (u, v) which quantized after decode is supplied to a selector 474 from the decryption section 470.

[0105] The DC multiplier QF (0 0) registered into the zero data and DC multiplier register 476 other than the DCT multiplier QF (u, v) given from the decryption section 470 is given to the data input terminal of a selector 474. The value of the DC multiplier QF (0 0) of a pixel block of the uniform color described in the software program of a game is written in DC multiplier register 476 by MPU40 of video game equipment. In addition, the value from which the DC multiplier QF (0 0) differs to a YUV signal, respectively is registered.

[0106] two null shown in drawing 20 (C) here -- the case where the run data NRD1 and NRD2 are decrypted is considered. the 1st null -- if the data NRL of the head of the run data NRD1 are detected by the decryption section 470 -- null -- the condition signal SS which tells that it is run data is outputted to a control section 472 from the decryption section 470. A control section 472 is controlled to output control signals CTL1 and CTL2 to the decryption section 470 and MPU40 immediately according to the condition signal SS, respectively, and to stop decryption actuation. Moreover, a control section 472 supplies a control signal CTL3 to a selector 474, and makes the DC multiplier QF (0 0) registered into DC multiplier register 476 choose as a DC multiplier of the first block. As shown in drawing 20 (B), when the value of the subject-copy image data f (x y) is 12, QF(0 0)=12 are registered into DC multiplier register 476. A control section 472 controls a selector 474 as an AC multiplier of 63 pieces further to choose zero data altogether. Drawing 28 shows the DCT multiplier QF (u, v) created in this way.

[0107] the 1st null -- since the run data NRD1 show that 16 blocks of a uniform color are continuing, the DCT multiplier QF (u, v) shown in drawing 28 is created about each of 16 pixel blocks. the 2nd null -- the DCT multiplier QF (u, v) shown in drawing 28 is similarly created [ each / of two blocks ] about the run data NRD2.

[0108] null -- after processing of run data is completed, a control section 472 outputs control signals CTL1 and CTL2 to the decryption section 470 and MPU40, and it controls to resume decryption actuation. null -- the value of the luminance signal Y of a pixel block and color-difference signals U and V as which it is specified that it is a uniform color can be easily changed with run data by changing the value of DC multiplier written in DC multiplier register 476 from MPU40. if it puts in another way -- null -- if run data are used, the color of the image field of a uniform color can be changed into a desired

color according to data other than compression image data. In this example, the value of the luminance signal Y of a block of a uniform color and color-difference signals U and V is specified in the software program which describes video game.

[0109] Drawing 29 is the block diagram showing other configurations of the Huffman decryption section. In this Huffman decryption section 210a, the DCT multiplier QF (u, v) outputted from the decryption section 470 bypasses a selector 474, and is directly given to the reverse quantization section 220. A selector 474 chooses one side of the DC multiplier QF (0 0) given from DC multiplier register 476, and zero data, bypasses the reverse quantization section 220, and supplies it to the IDCT section 230 directly. The control section 478 is outputting the 4th control signal CTL4 other than the three same control signals CTL1-CTL3 as drawing 27 to the IDCT section 230.

[0110] A control section 478 outputs control signals CTRL1, CTL2, and CTL4 to the decryption section 470, MPU40, and the IDCT section 230 according to the condition signal SS which shows the class of compressed data, respectively. Decode of the block data usual in the decryption section 470 supplies the decrypted DCT multiplier QF (u, v) to the reverse quantization section 220.

[0111] the 1st null shown in drawing 20 (C) -- if the data NRL of the head of the run data NDR 1 are detected by the decryption section 470 -- null -- the condition signal SS which tells that it is run data is outputted to a control section 478 from the decryption section 470. A control section 478 is controlled to output control signals CTL1 and CTL2 to the decryption section 470 and MPU40 immediately according to the condition signal SS, respectively, and to stop decryption actuation. Moreover, a control section 478 switches the level of the control signal CTL3 given to a selector 474, and makes the DC multiplier QF (0 0) registered into DC multiplier register 476 choose. A control section 478 controls a selector 474 as an AC multiplier of 63 pieces further to choose zero data altogether. A control section 478 outputs a control signal CTL4 to coincidence to the IDCT section 230, and controls it to choose and carry out inverse transformation of the output of a selector 474.

[0112] null -- after processing of run data is completed, control signals CTL1 and CTL2 are outputted to the decryption section 470 and MPU40 from a control section 478, and it controls to resume decode actuation.

[0113] thus -- the circuit shown in drawing 29 -- null -- in the case of processing of run data, since the DCT multiplier QF (u, v) outputted from a selector 474 bypasses the reverse quantization section 220 and is directly supplied to the IDCT section 230, there is an advantage that the operation error by reverse quantization does not arise. For example, since the operation error by quantization is pressed down by min when the video encoder unit 50 ( drawing 3 ) performs chroma-key processing which detects the part of a specific color and is made into a transparent plane color, it is possible to make the pixel block of a desired color into a transparent plane color certainly.

[0114] the above-mentioned example -- setting -- null -- although it enabled it to set only DC multiplier as arbitration with run data, it is also possible to enable it to set up a desired value to the predetermined part (for example, QF (1 0) and QF (0 1)) of AC multiplier. in this case, null -- run data express the number of a series of pixel blocks which has the same image pattern.

[0115] F. modification: -- the range which this invention is not restricted to the above-mentioned example, and does not deviate from that summary in addition -- setting -- various voice -- it is possible to set like and to carry out, for example, the following deformation is also possible.

[0116] (1) Although two-dimensional DCT conversion was taken up as orthogonal transformation in the above-mentioned example, to this invention, the orthogonal transformation (for example, K-L conversion and a Hadamard transform) of arbitration is available. Moreover, as entropy code modulation, coding (for example, algebraic-sign-izing and MEL coding) of arbitration other than Huffman coding can be used.

[0117] (2) You may realize by hardware and image data compression equipment 100 may realize image data decompression equipment 200 by software. Drawing 30 is a flow chart which shows the procedure of the processing which elongates the compressed data shown in drawing 18 with software. At step S1, the contents of compressed data are judged from the value of a header unit. When compressed data expresses the base quantity child-sized table BQT, in step S2, the base quantity child-sized table BQT is

memorized in memory, and it returns to step S1.

[0118] When compressed data expresses the image, the class of initial data of each data unit contained in compressed data circles in step S3 is judged. the code data of the quantization level multiplier QCx indicated to be a data unit to drawing 18 (C) here, block data, and null -- each of run data is meant. the Huffman-coding word which expresses the quantization level multiplier QCx as initial data of a data unit as shown in drawing 18 (C), (D), (E), and (F), the Huffman-coding word of DC multiplier \*\*DC of block data, and null -- there are three kinds of the Huffman-coding words NRL showing a run.

[0119] When initial data are the Huffman-coding words of the quantization level multiplier QCx, in step S4, the quantization table QT is created by multiplying by this quantization level multiplier QCx and the base quantity child-sized table BQT. In addition, when the value of that the description mentioned above on the occasion of this multiplication, i.e., (1) multiplication result, sets the quantization level more than maximum as maximum, not quantizing (2) DC multiplier, and (3) quantization level multiplier QCx is 0, setting-as 1-all quantization level \*\* is realized. If the quantization table QT is created in step S4, it will return to step S3.

[0120] When it is judged that initial data are the Huffman-coding words of DC multiplier \*\*DC in step S3, the data for a 1-pixel block are decrypted in step S5, and the DCT multiplier QF (u, v) is obtained. Reverse quantization is performed at step S6, and two-dimensional DCT inverse transformation is performed at step S7. If processing of steps S5-S7 is repeated about all pixel blocks included in one-set block data (drawing 18 (D)), it will return from step S8 to step S3.

[0121] step S3 -- setting -- initial data -- null -- when it is judged that it is the Huffman-coding word NRL which shows that it is run data, while DC component of a 1-pixel block is set as the value specified beforehand in step S9, all of 63 AC components are set as 0. At step S10, two-dimensional DCT inverse transformation is performed to the DCT multiplier created in this way. null -- only the block count specified by run data will return from step S11 to step S3, if step S9 and processing of S10 are repeated. In this way, image data ff (x y) is restored by covering all the compressed data sections and performing processing of steps S3-S11.

[0122] (3) Although the above-mentioned example showed the example which applied this invention to video game equipment, this invention can be applied to all kinds of image processing system.

[0123] (4) Although the quantization level multiplier QCx and each quantization level in the quantization table QT were made into the integer in the above-mentioned example, a number including a decimal is sufficient as these.

---

[Translation done.]

\* NOTICES \*

JPO and NCIP are not responsible for any damages caused by the use of this translation.

1. This document has been translated by computer. So the translation may not reflect the original precisely.
2. \*\*\*\* shows the word which can not be translated.
3. In the drawings, any words are not translated.

---

## DESCRIPTION OF DRAWINGS

---

### [Brief Description of the Drawings]

[Drawing 1] The block diagram showing the function of the compression equipment of image data, and expanding equipment which applied one example of this invention.

[Drawing 2] The block diagram showing the concrete configuration of image data compression equipment 100.

[Drawing 3] The block diagram showing the concrete configuration of image data decompression equipment 200.

[Drawing 4] The top view showing a subject-copy image.

[Drawing 5] The explanatory view showing the array of the DCT multiplier F (u, v).

[Drawing 6] The explanatory view showing basic actuation of compression/expanding.

[Drawing 7] The explanatory view showing compression/expanding actuation in case the quantization level multiplier QCx is 4.

[Drawing 8] The explanatory view showing the actuation when not changing DC component in quantization.

[Drawing 9] The explanatory view showing compression/expanding actuation in case the quantization level multiplier QCx is 0.

[Drawing 10] The block diagram showing the function of DC multiplier coding section.

[Drawing 11] Drawing showing the category-ized table in Huffman coding.

[Drawing 12] The explanatory view showing an example of the Huffman-coding table HTDC for DC multipliers.

[Drawing 13] The block diagram showing the function of AC multiplier coding section.

[Drawing 14] The explanatory view showing the usual route of a zigzag scan of AC multiplier.

[Drawing 15] The explanatory view showing the two-dimensional Huffman-coding table for AC multipliers.

[Drawing 16] Drawing showing the contents of the Huffman-coding table.

[Drawing 17] The explanatory view showing an example of Huffman coding.

[Drawing 18] The explanatory view showing the configuration of compressed data.

[Drawing 19] The explanatory view showing the relation of a block of each signal of YUV.

[Drawing 20] null -- the explanatory view showing the image expressed by run data.

[Drawing 21] Drawing showing other parts of the Huffman-coding table for AC multipliers.

[Drawing 22] The block diagram showing the internal configuration of the reverse quantization table creation section 250.

[Drawing 23] The block diagram showing the internal configuration of a latch circuit 253 and the multiplication unit 254.

[Drawing 24] The timing chart which shows actuation of the circuit shown in drawing 23 .

[Drawing 25] The block diagram showing the internal configuration of U signal start detector 416 and V signal start detector 418.

[Drawing 26] The block diagram showing the internal configuration of a clipping circuit 426 and the

zero value correction circuit 428.

[Drawing 27] The block diagram showing the internal configuration of the Huffman decryption section 210.

[Drawing 28] null -- drawing showing the DCT multiplier which decrypted run data and was obtained.

[Drawing 29] The block diagram showing other configurations of the Huffman decryption section.

[Drawing 30] The flow chart which shows the procedure of the processing which decodes compressed data with software.

[Drawing 31] The block diagram showing conventional image data compression equipment and expanding equipment.

[Description of Notations]

20 -- Video game equipment

21 -- ROM

24 26 -- Gamepad

28 -- Color television

30 -- Video signal cable

32 -- CD-ROM drive

34 -- Loudspeaker

36 -- SCSI bus

38 -- Loudspeaker

40 -- MPU

40a -- Operation part

40b -- Controller

41 -- Main memory

42 -- ROM

43 -- M-BUS

45 -- Picture signal control unit

45 a--MPUI/F

45 b--SCSI controller

45 c--AFFIN converter

45d -- Graphic controller

45e -- Sound controller

49 -- VDP unit

50 -- Video encoder unit

50a -- Look-up table

50b -- Priority setting section

50c -- Mixer

50c -- Image merge section

50d -- Superimposition section

50 e--DAC

52 -- Voice output unit

The 52 a--ADPCM sections

55 -- RAM

59 -- RAM

60 -- NTSC converter

100 -- Image data compression equipment

101 -- CPU

102 -- Main memory

103 -- Keyboard

104 -- Mouse

105 -- Magnetic disk drive

106 -- Optical-magnetic disc equipment

110 -- The DCT section  
120 -- Quantization section  
130 -- Huffman coding section  
131 -- Block delay section  
132 -- Adder  
133 -- Categorizing processing section  
134--1-dimensional Huffman coding section  
135 -- Zigzag scan section  
136 -- Judgment section  
137 -- Run-length counter  
138 -- Categorizing section  
140 -- Quantization table creation section  
150 -- Huffman-coding table memory  
200 -- Image data decompression equipment  
210 -- Huffman decryption section  
220 -- Reverse quantization section  
230 -- The IDCT section  
240 -- Huffman-coding table memory  
250 -- Reverse quantization table creation section  
251 -- RAM  
252 -- Address-generation circuit  
253 -- Latch circuit  
254 -- Multiplication unit  
402,404 -- Latch  
412 -- Synchronous-clock creation circuit  
414 -- AND circuit  
416 -- U signal start detector  
418 -- V signal start detector  
420 -- NAND circuit  
422 -- Selector  
424 -- Multiplier  
426 -- Clipping circuit  
428 -- Zero value correction circuit  
432,436 -- D flip-flop  
434,438 -- NAND circuit  
450 -- 4 input OR circuit  
452 -- 2 input OR circuit  
460 -- Inverter  
462 -- AND circuit  
464 -- OR circuit  
470 -- Decryption section  
472 -- Control section  
474 -- Selector  
476 -- DC multiplier register  
478 -- Control section  
540 -- Image data compression equipment  
542 -- Orthogonal transformation section  
544 -- Quantization section  
546 -- Entropy-code-modulation section  
550 -- Image data decompression equipment  
550 -- Expanding equipment

552 -- Reverse orthogonal transformation section  
554 -- Reverse quantization section  
556 -- Entropy decryption section  
562 -- Quantization table  
564 -- Sign table  
BG -- Background  
BQT -- Base quantity child-sized table  
CLK -- Basic clock signal  
D1-D13 -- Output of a multiplier 424  
D13 -- Sign bit  
EN -- Enable signal  
 $f(x, y)$  -- Subject-copy image data  
Image data ff(ed) -- (x, y) decoded  
 $F(u, v)$  -- The original DCT multiplier  
 $FF(u, v)$  -- Decoded DCT multiplier  
The Huffman-coding word for HFAC--AC multipliers  
The Huffman-coding word for HFDC--DC multipliers  
The Huffman-coding table for HTAC--AC multipliers  
The Huffman-coding table for HTDC--DC multipliers  
HT -- Huffman-coding table  
ID -- Discernment data  
NNNN -- Zero run length  
NRL -- null -- run data  
PB -- Pixel block  
PX -- Pixel  
QCx -- Quantization level multiplier  
QEN -- Enable signal  
The DCT multiplier QF(ed) -- quantized  
QT -- Quantization table  
SCK -- Synchronizing clock signal  
SEL -- Selection signal  
SS -- Condition signal  
SSSS -- Category  
SWTCH -- Block change signal  
U, V -- Color-difference signal  
USTRT -- U start signal  
VSTR -- V start signal  
Y -- Luminance signal  
ZZ -- Compression image data

---

[Translation done.]

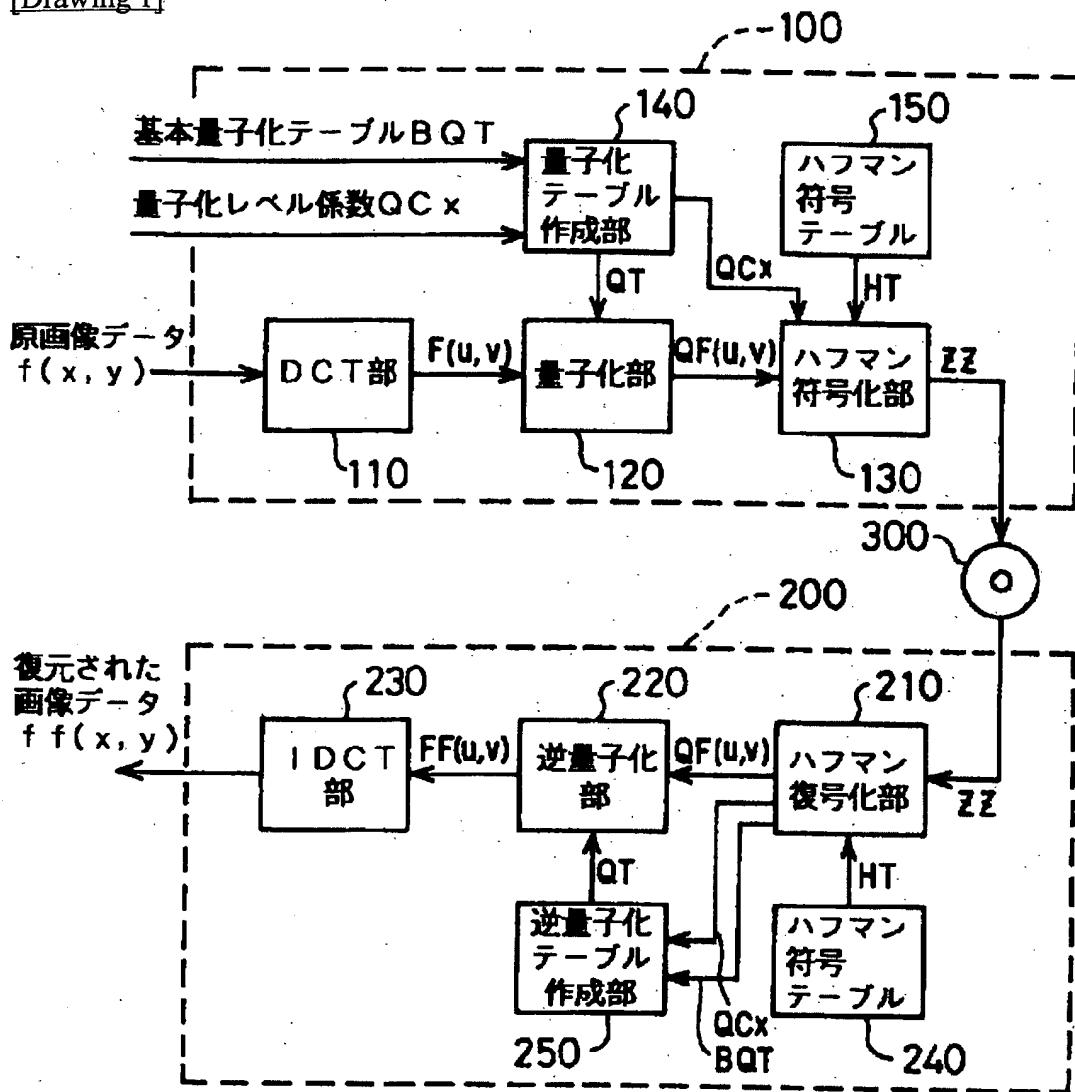
\* NOTICES \*

JPO and NCIPPI are not responsible for any damages caused by the use of this translation.

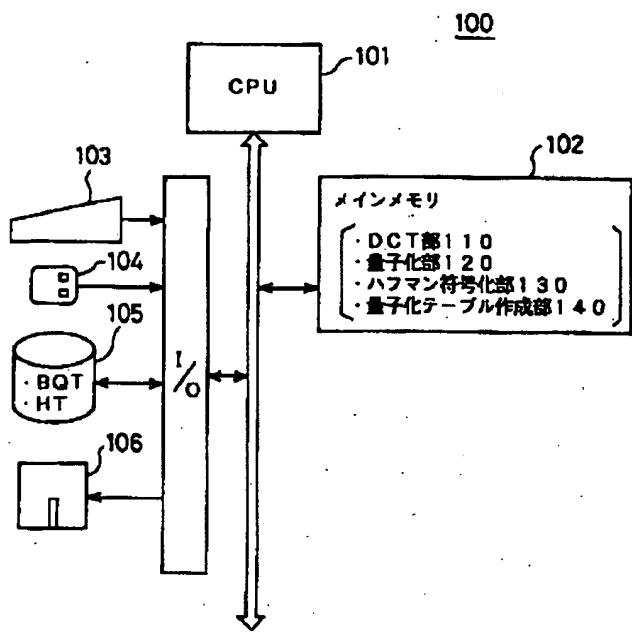
1. This document has been translated by computer. So the translation may not reflect the original precisely.
2. \*\*\*\* shows the word which can not be translated.
3. In the drawings, any words are not translated.

DRAWINGS

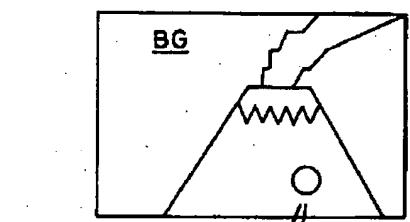
[Drawing 1]



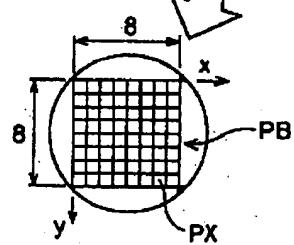
[Drawing 2]



[Drawing 4]  
(A)

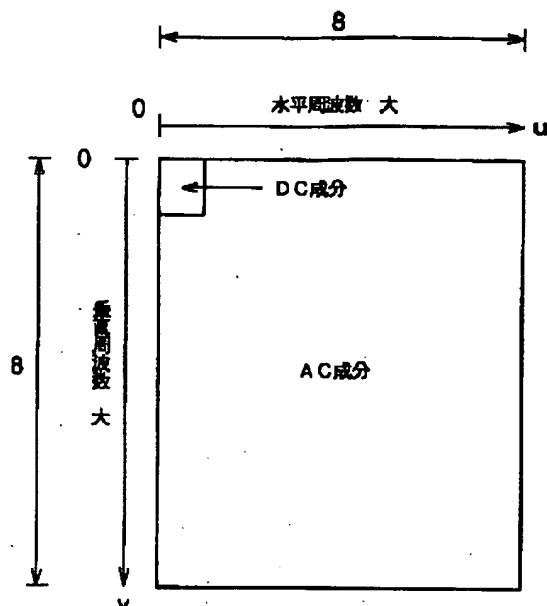


(B)



[Drawing 5]

変換係数  $F(u, v)$  の配列

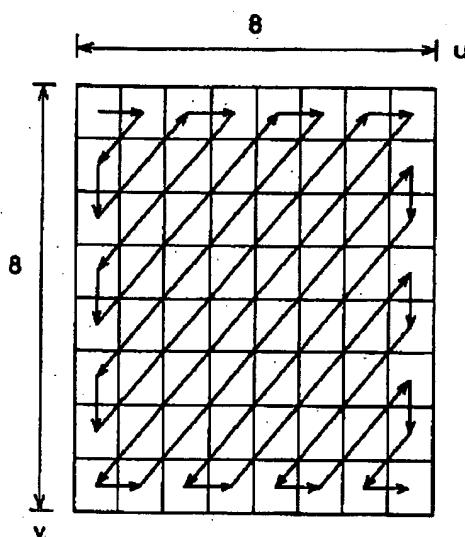


[Drawing 11]

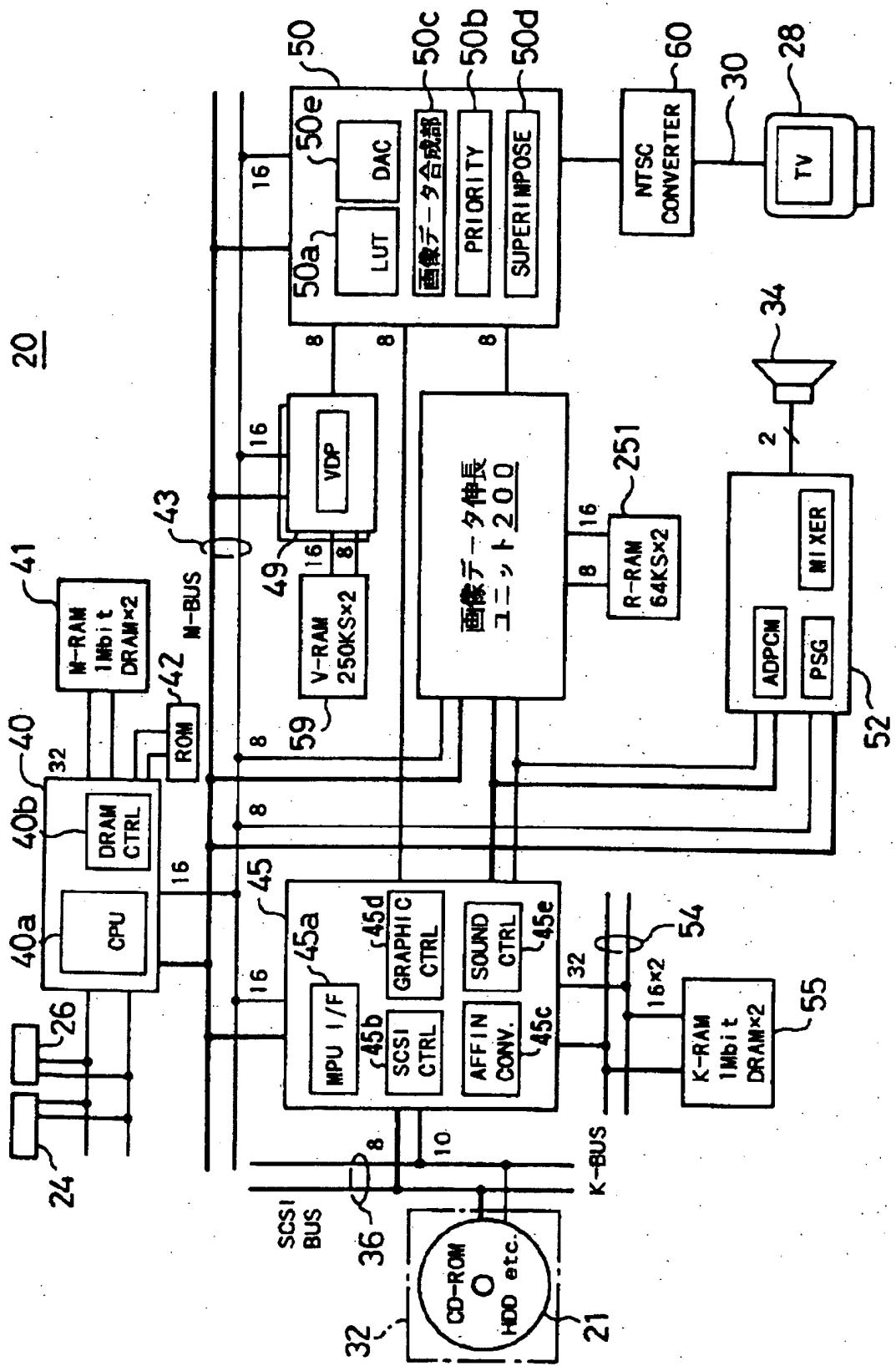
カテゴリ SSSS	$\Delta DC / AC$ 係数値	識別データ ビット数
0	0	0
1	-1, 1	1
2	-3, -2, 2, 3	2
3	-7, -4, 4, 7	3
4	-15, -8, 8, 15	3
5	-31, -16, 16, 31	5
6	-63, -32, 32, 63	6
7	-127, -64, 64, 127	7
8	-255, -128, 128, 255	8
9	-511, -256, 256, 511	9

[Drawing 14]

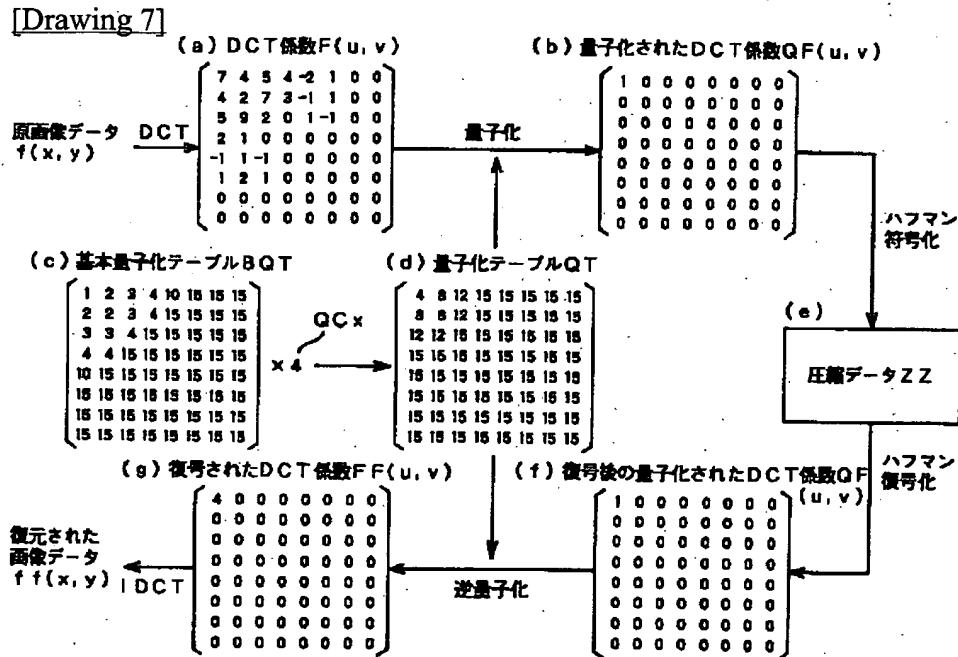
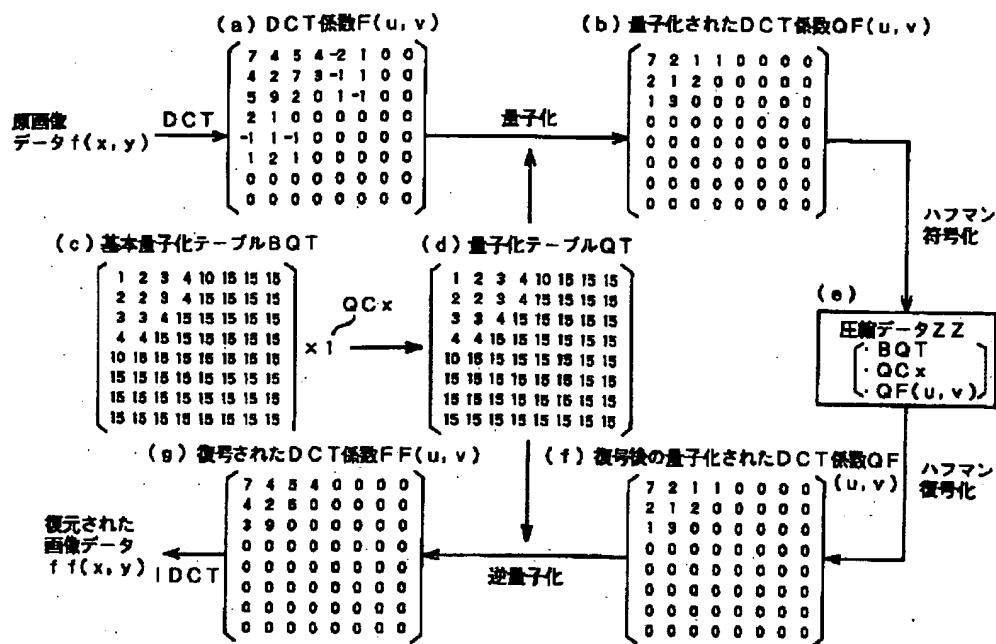
ジグザグスキャン



[Drawing 3]



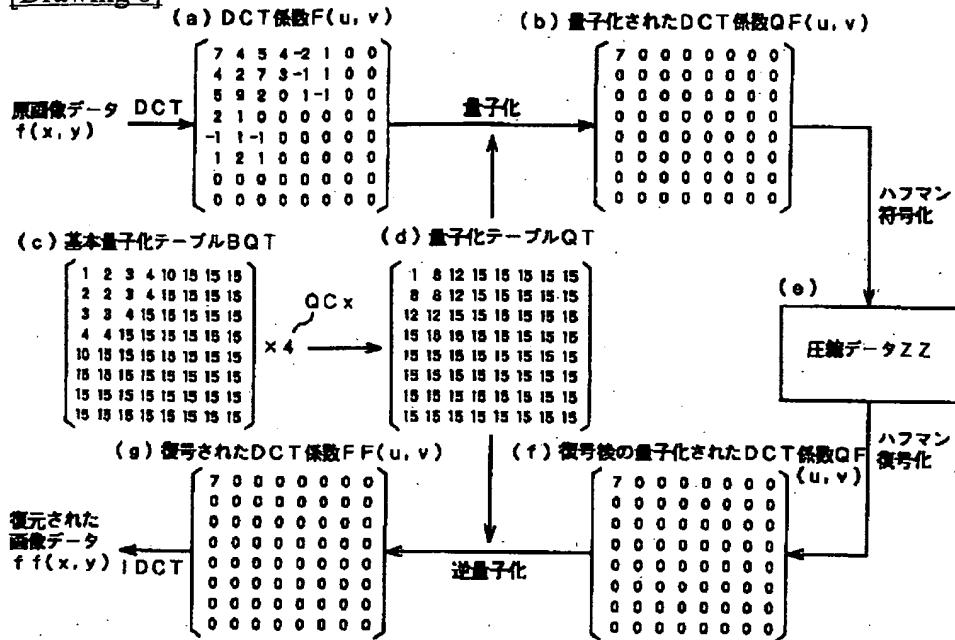
[Drawing 6]



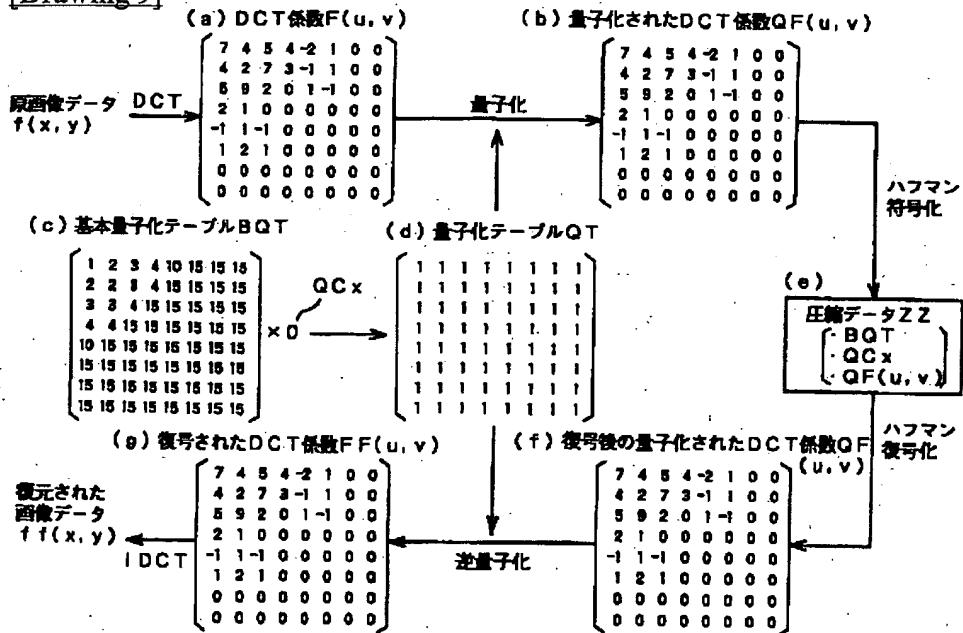
[Drawing 15]  
AC係数用ハフマン符号テーブル

		カテゴリ SSSS			
		0	1	-----	9
ゼロ 長 NNNN	0	EOB	AC係数 ハフマン符号器 HFac		
	1	ZRL			
	2	-			
	3	-			
	15	-			

### [Drawing 8]

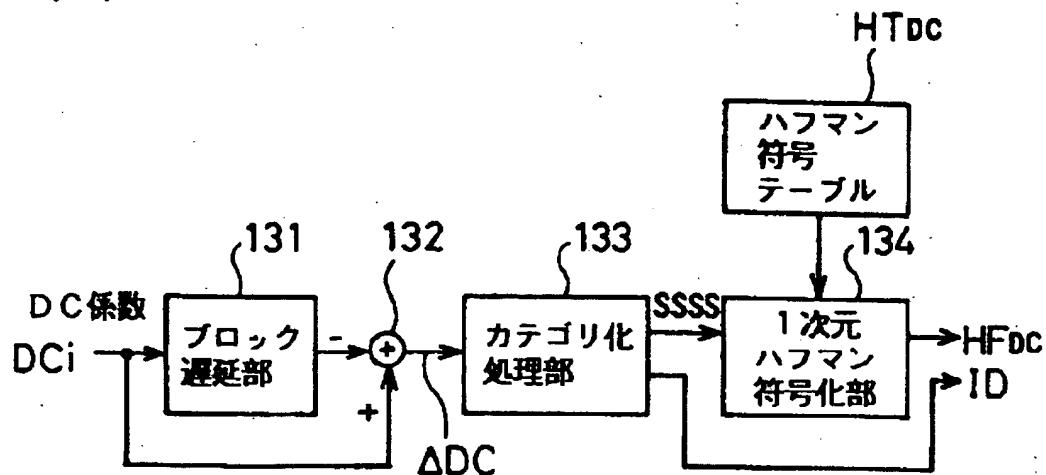


### [Drawing 9]

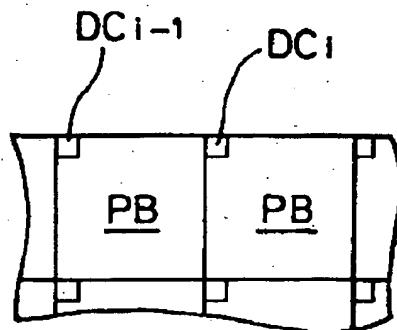


### [Drawing 10]

(A)

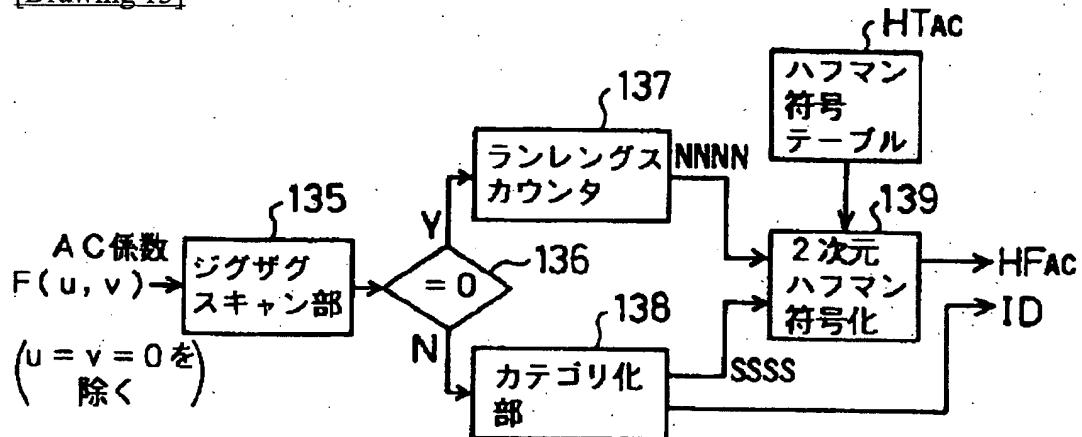


(B)



$$\Delta DC = DC_i - DC_{i-1}$$

[Drawing 13]



[Drawing 12]

カテゴリ SSSS	$\Delta$ DC 係数 ハフマン符号語 HF <sub>DC</sub>	
	Y信号用	U信号、V信号用
0	000	00
1	1110	01
2	001	10
3	010	110
4	011	1110
5	100	11110
6	101	111110
7	110	1111110
8	111100	11111110
9	1111010	11111111
10		
11		
12		
13		
14		
15(NUL)	1111011	
16(Q0)	111110000	
17(Q1)	111110001	
18(Q2)	111110010	
19(Q3)	111110011	
20(Q4)	111110100	
21(Q5)	111110101	
22(Q6)	111110110	
23(Q7)	111110111	
24(Q8)	111111000	
25(Q9)	111111001	
26(Q10)	111111010	
27(Q11)	111111011	
28(Q12)	111111100	
29(Q13)	111111101	
30(Q14)	111111110	
31(Q15)	111111111	

[Drawing 16]

ゼロラン長 NNNN	AC 係数ハフマン符号語 HF <sub>AC</sub>	
	Y信号用	U信号、V信号用
0/0 (EOB)	11111	11111
0/1	00	00
0/2	01	01
0/3	100	1010
0/4	1010	11000
0/5	11000	1110010
0/6	110110	111100010000
0/7	111011100	111100010001
0/8	111100010000	111100010010
0/9	11101111000	11101111000
1/0 (ZRL)	11110111111	111101111111
1/1	1011	100
1/2	11001	11001
1/3	1110100	1110011
1/4	111100010001	111100010011
1/5	111100010010	111100010100
1/6	111100010011	111100010101
1/7	111100010100	111100010110
1/8	111100010101	111100010111
1/9	11101111001	11101111001

[Drawing 17]

(A) DCT係数F(u, v)

DCi	12	0	4	0	0	0	1	0
	-4	3	0	-2	0	0	0	0
	0	5	3	0	1	0	0	0
	-4	0	0	0	0	0	0	0
	0	2	0	0	0	0	0	0
	0	0	0	0	0	1	0	0
	0	0	0	0	0	0	0	0
	0	0	0	0	0	0	0	0

(B) DC係数符号化(DC<sub>i</sub>-1=0)

$$\Delta DC = 12 \rightarrow SSSS = 4 \rightarrow HF_{DC} + ID = \underline{\underline{0111100}} \\ HF_{DC} \quad ID$$

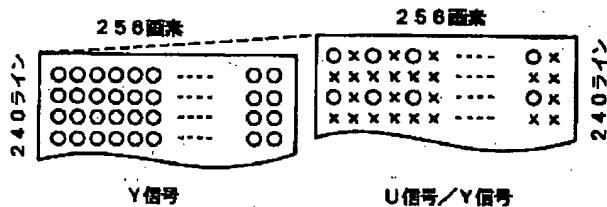
(C) AC係数符号化

- ・ジグザグスキャン : 0, -4, 0, 3, 4, 0, 0, 6
- ・NNNN/SSSS = 1/3 1/2 0/3 2/3
- ・HFAC + ID = 11101000111100111  
HFAC ID HFAC ID

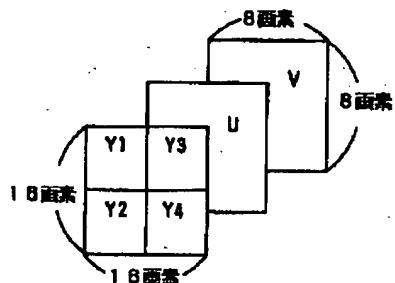
[Drawing 19]

(A)

(B)

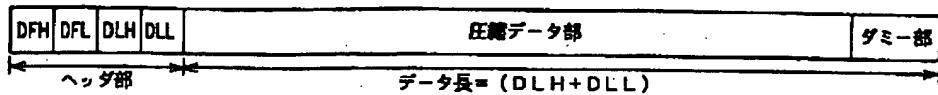


(C)

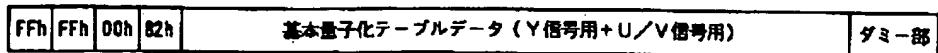


[Drawing 18]

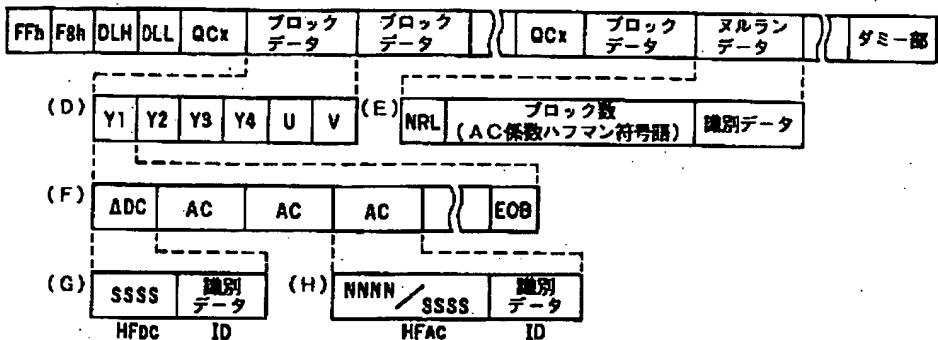
(A) 圧縮データ基本構造



(B) 基本量子化テーブルデータ

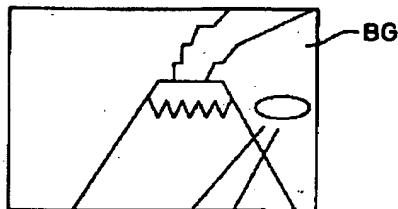


(C) フルカラー自然画像圧縮データ

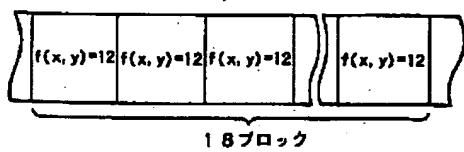


[Drawing 20]

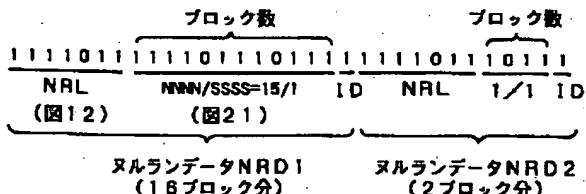
(A)



(B) 原画像データ



(C) ヌルランデータ



[Drawing 21]

ゼロラン長 NNNN カテゴリ SSSS	AC係数ハフマン符号語HFAc	
	Y信号用	U信号、V信号用
0/1	00	00
1/1	1011	100
2/1	11010	1011
3/1	110111	11010
4/1	111000	11011
5/1	111001	111000
6/1	1110110	1110101
7/1	111011101	1110110
8/1	11110011111	111011101
9/1	111101000111	111101000111
10/1	111101001111	111101001111
11/1	111101010111	111101010111
12/1	111101011111	111101011111
13/1	111101100111	111101100111
14/1	111101101111	111101101111
15/1	111101110111	111101110111

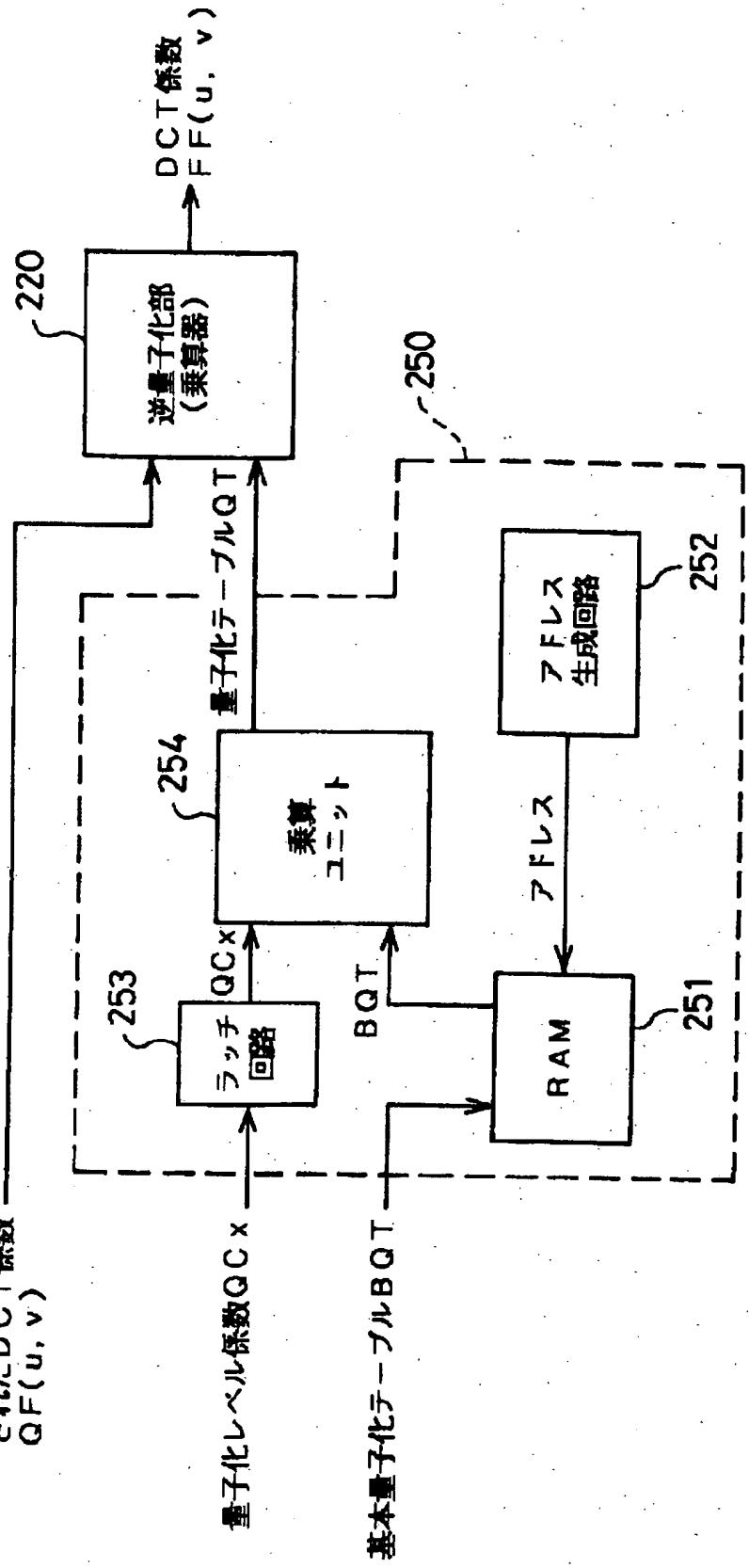
ヌルランデータに使用する場合：

NNNN = 【ブロック数】 - 1

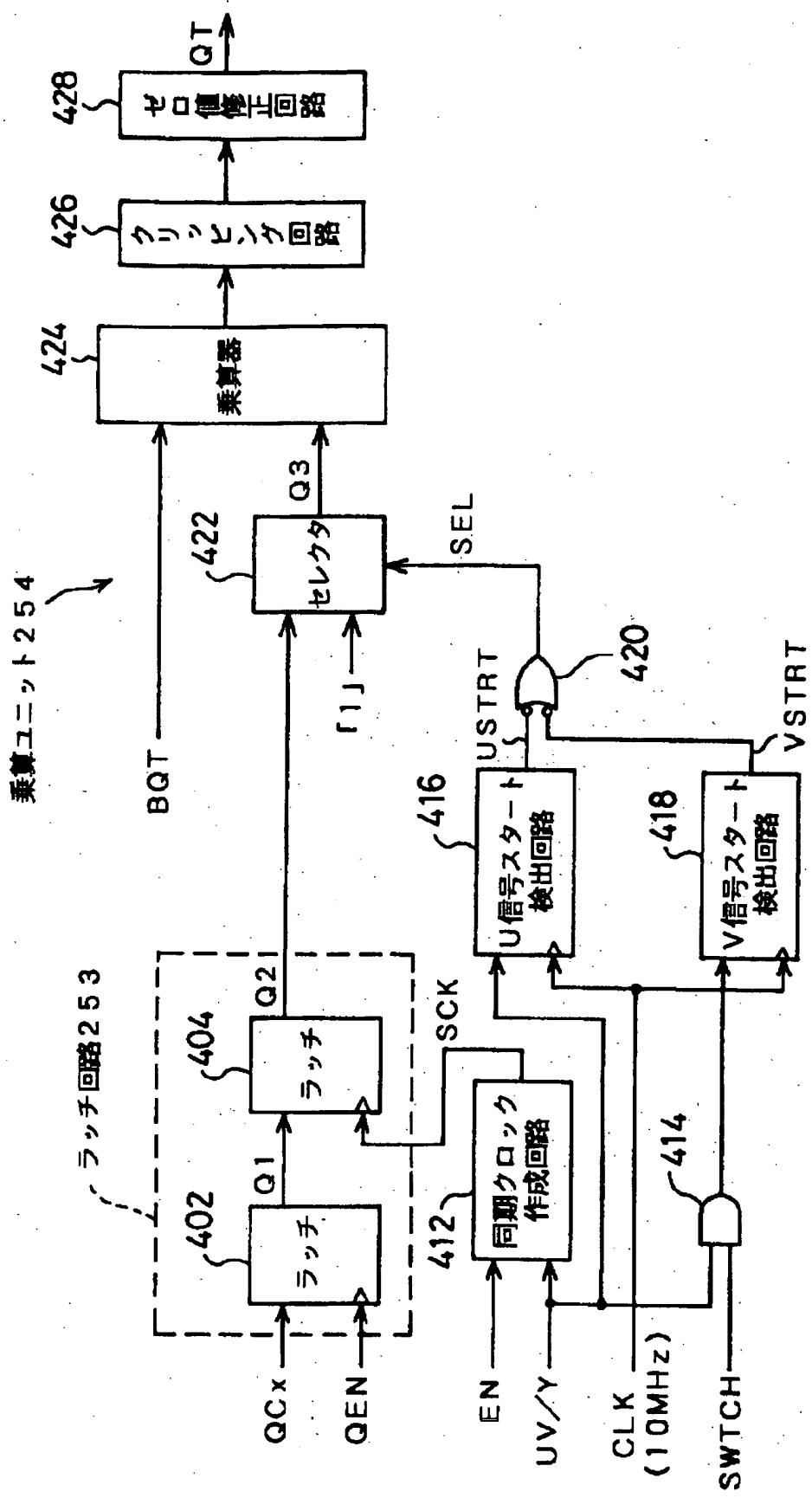
SSSS = 1

[Drawing 22]

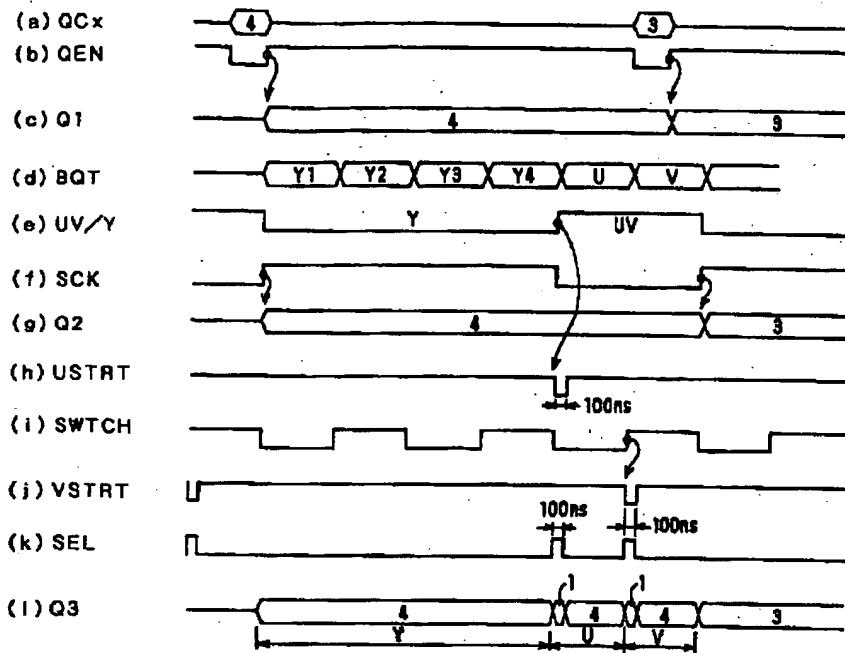
復号後の量子化  
されたDCT係数  
 $QF(u, v)$



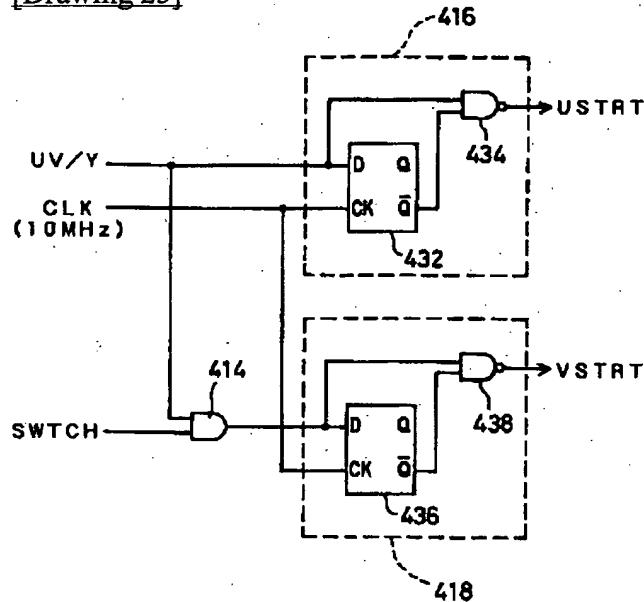
[Drawing 23]



[Drawing 24]

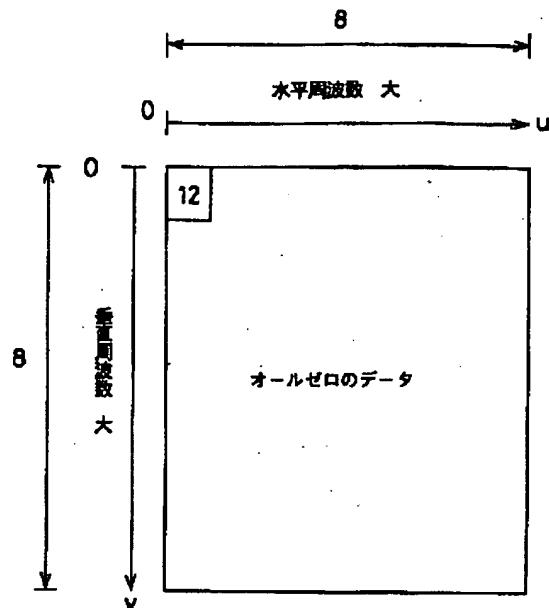


[Drawing 25]

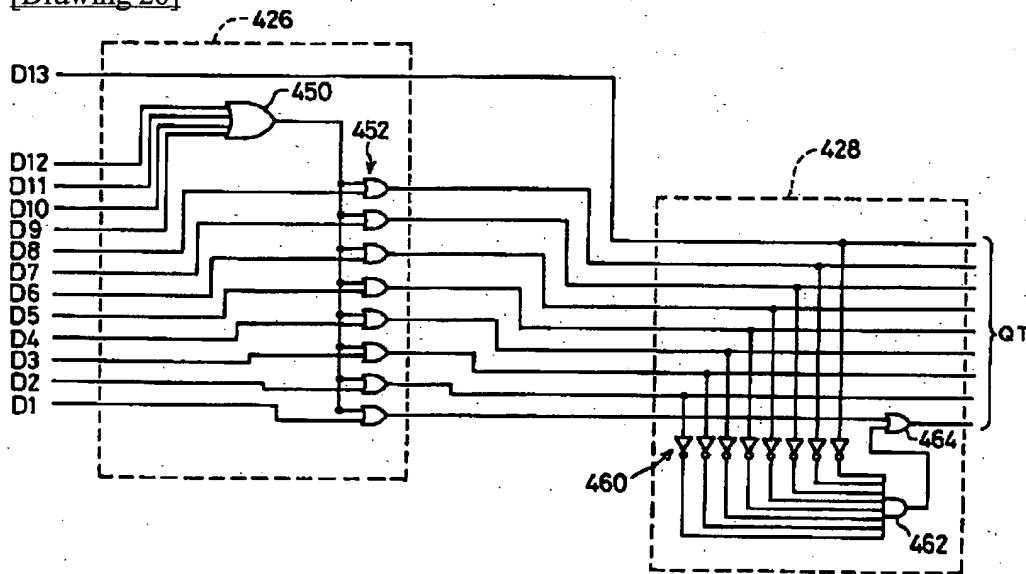


[Drawing 28]

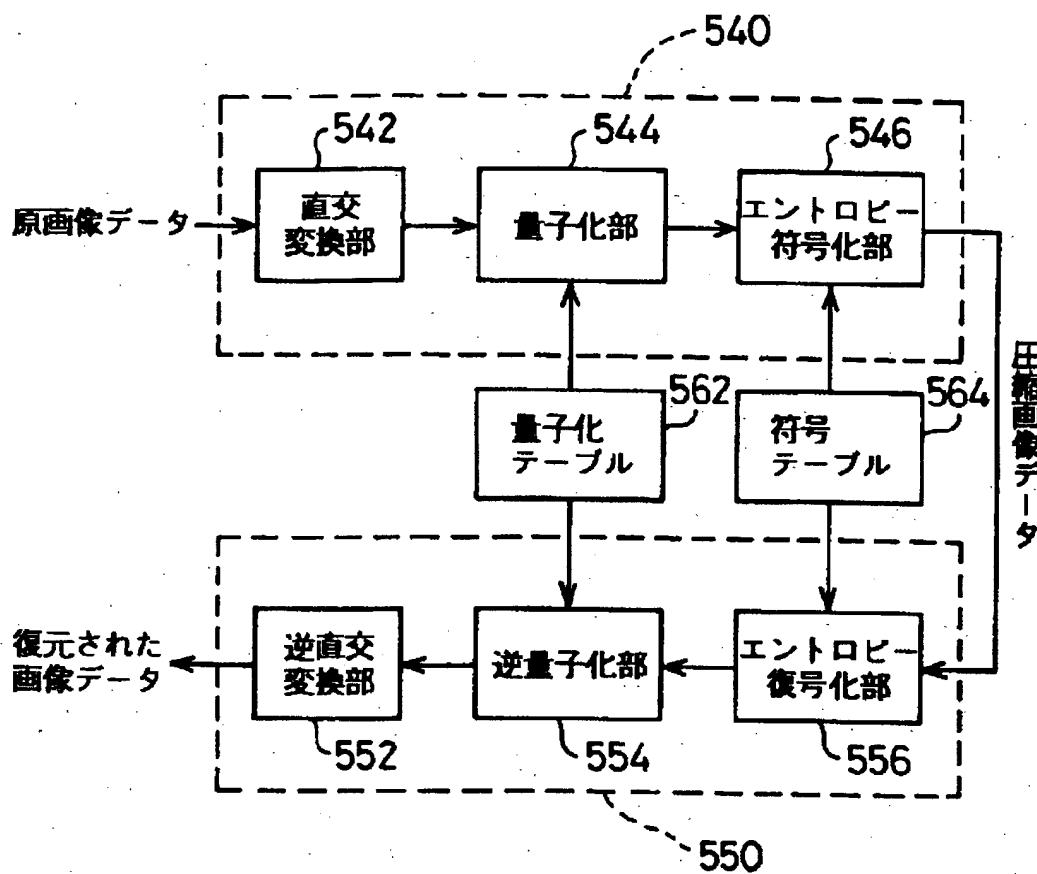
ヌルテンデータから復号されたDCT係数QF(u, v)



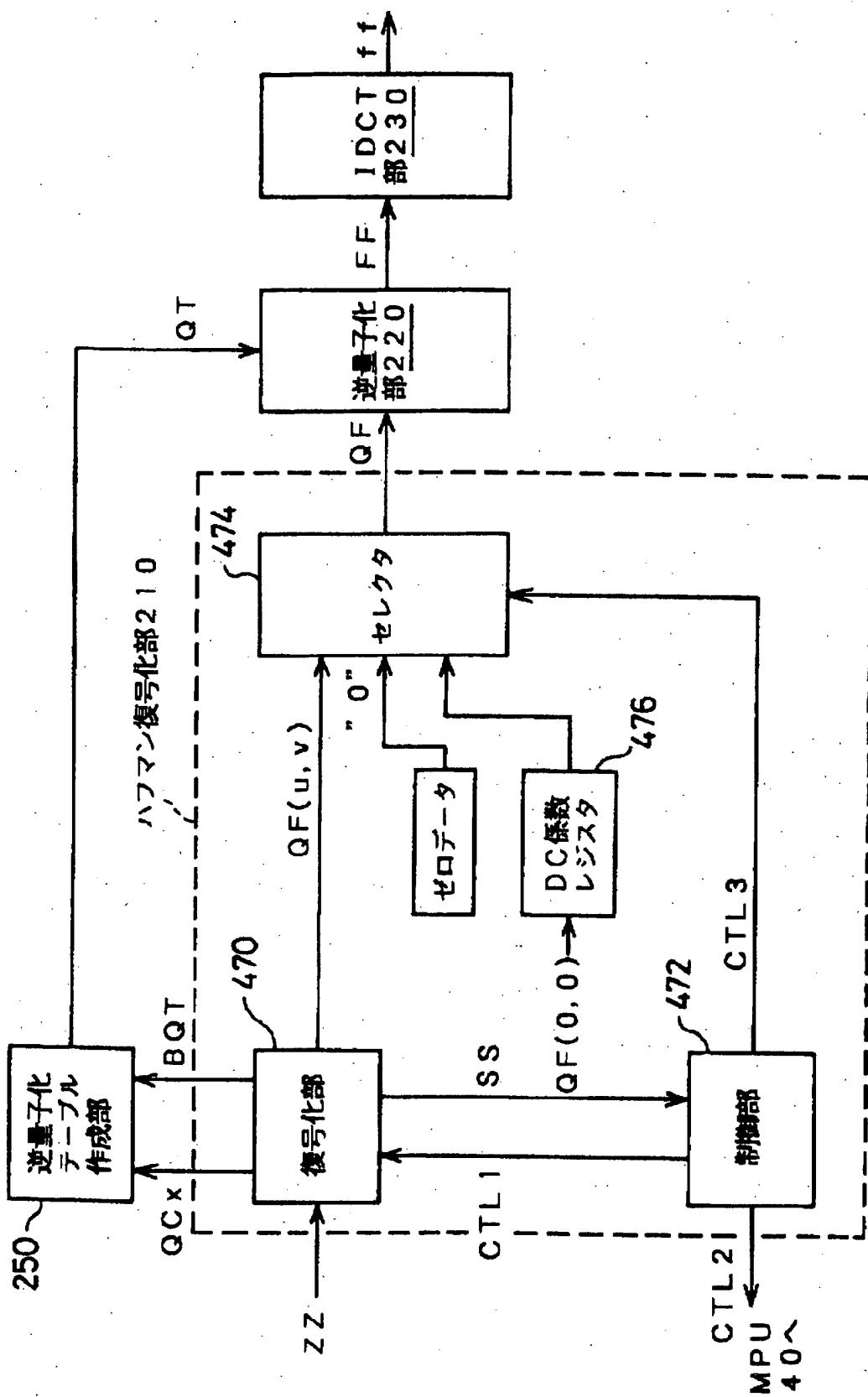
[Drawing 26]



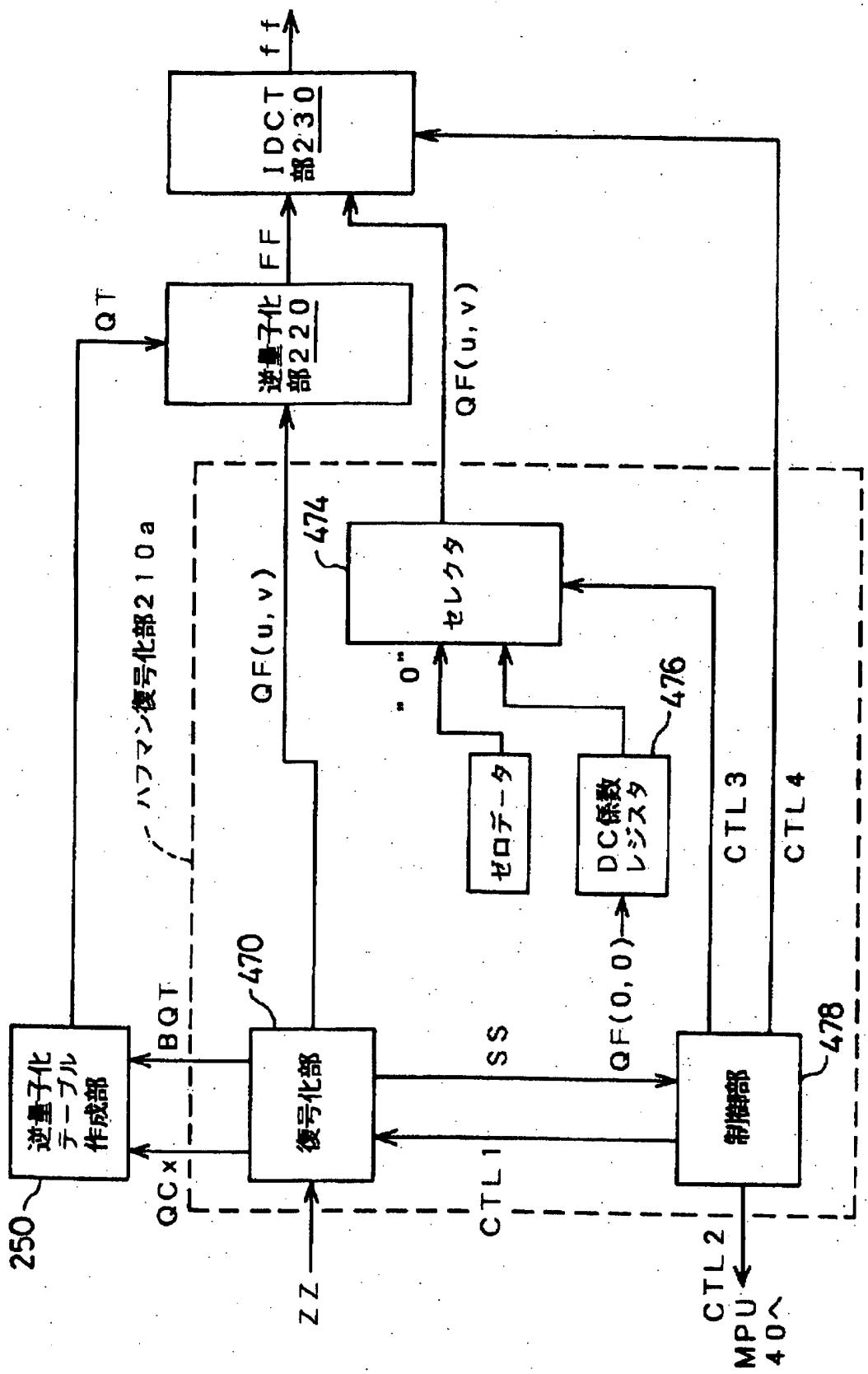
[Drawing 31]



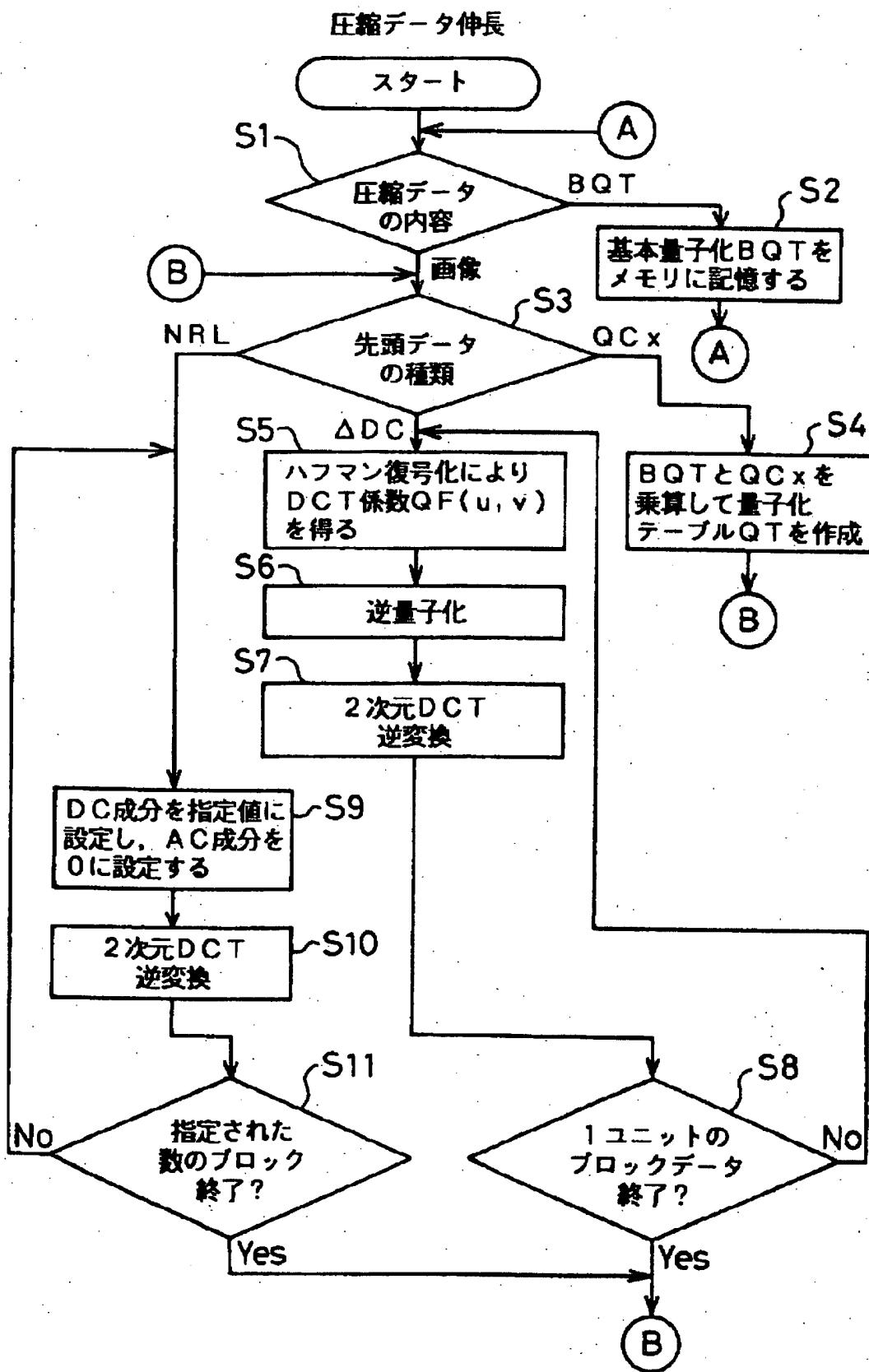
[Drawing 27]



### [Drawing 29]



[Drawing 30]



[Translation done.]

# PATENT ABSTRACTS OF JAPAN

(11)Publication number : 07-123413

(43)Date of publication of application : 12.05.1995

(51)Int.Cl.

H04N 7/30  
G06T 9/00  
H03M 7/30  
H03M 7/40  
H04N 1/41



(21)Application number : 05-265812

(71)Applicant : SEIKO EPSON CORP  
HUDSON SOFT CO LTD

(22)Date of filing : 28.09.1993

(72)Inventor : MIYANE TOSHIKI  
SEKIMOTO UICHI

(30)Priority

Priority number : 04258398

Priority date : 28.09.1992

Priority country : JP

05 66831

25.03.1993

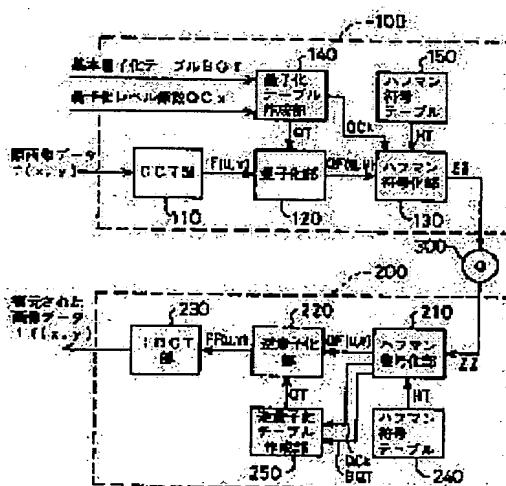
JP

## (54) METHOD AND DEVICE FOR COMPRESSING/EXTENDING IMAGE DATA

### (57)Abstract:

**PURPOSE:** To reduce the data amount of compressed image data in comparison with the conventional data amount.

**CONSTITUTION:** The compressed image data contains code data showing a quantizing level coefficient QCx among block data. An inverse quantizing table preparation part 250 prepares a quantizing table QT by multiplying a DCT coefficient QF(u,v) provided by decoding compressed image data ZZ and the quantizing level coefficient QCx. An inverse quantizing part 220 performs inverse quantizing by using the quantizing table QT. Since the quantizing level coefficient QCx is inserted between the block data of the compressed image data, the quantizing table QT is updated each time the quantizing level coefficient QCx is decoded. On the other hand, data expressing plural continuous picture element blocks in the same pattern color are inserted to the compressed image data in any special data format (null run data).



### LEGAL STATUS

[Date of request for examination] 28.04.2000

[Date of sending the examiner's decision of rejection] 13.01.2004

[Kind of final disposal of application other than the examiner's decision of rejection or

[application converted registration]  
[Date of final disposal for application]  
[Patent number] 3673529  
[Date of registration] 28.04.2005  
[Number of appeal against examiner's decision  
of rejection] 2004-02865  
[Date of requesting appeal against examiner's  
decision of rejection] 12.02.2004  
[Date of extinction of right]

(19)日本国特許庁 (JP)

## (12) 公開特許公報 (A)

(11)特許出願公開番号

特開平7-123413

(43)公開日 平成7年(1995)5月12日

(51)Int.Cl. <sup>6</sup>	識別記号	庁内整理番号	F I	技術表示箇所
H 0 4 N 7/30				
G 0 6 T 9/00				
H 0 3 M 7/30	A 8842-5J	8420-5L	H 0 4 N 7/ 133 G 0 6 F 15/ 66	Z 3 3 0 H
			審査請求 未請求 請求項の数29 FD (全34頁)	最終頁に続く

(21)出願番号 特願平5-265812  
 (22)出願日 平成5年(1993)9月28日  
 (31)優先権主張番号 特願平4-258398  
 (32)優先日 平4(1992)9月28日  
 (33)優先権主張国 日本(JP)  
 (31)優先権主張番号 特願平5-66831  
 (32)優先日 平5(1993)3月25日  
 (33)優先権主張国 日本(JP)

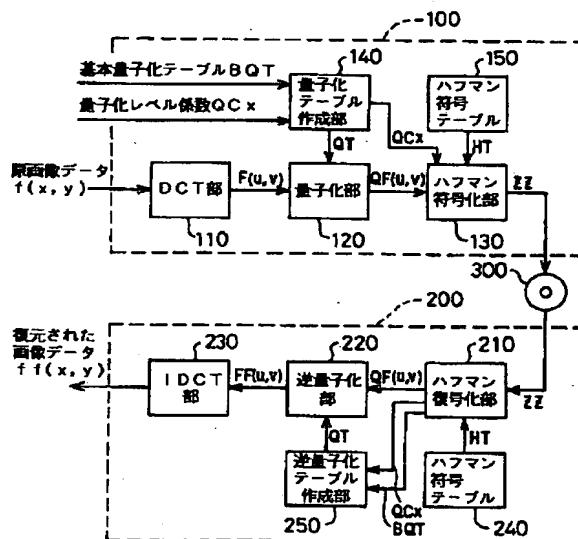
(71)出願人 000002369  
 セイコーエプソン株式会社  
 東京都新宿区西新宿2丁目4番1号  
 (71)出願人 591095856  
 株式会社ハドソン  
 北海道札幌市豊平区平岸3条7丁目26番地  
 (72)発明者 宮根 俊樹  
 長野県諏訪市大和三丁目3番5号 セイコーエプソン株式会社内  
 (72)発明者 関本 宇一  
 長野県諏訪市大和三丁目3番5号 セイコーエプソン株式会社内  
 (74)代理人 弁理士 下出 隆史 (外1名)

(54)【発明の名称】 画像データの圧縮・伸長方法およびそのための装置

## (57)【要約】

【目的】 圧縮画像データのデータ量を従来に比べて低減する。

【構成】 圧縮画像データは、ブロックデータの間に量子化レベル係数QCxを示す符号データを含んでいる。逆量子化テーブル作成部250は、圧縮画像データZZを復号して得られたDCT係数QF(u,v)と量子化レベル係数QCxを乗算することによって量子化テーブルQTを作成する。逆量子化部220は、量子化テーブルQTを用いて逆量子化を行なう。量子化レベル係数QCxは、圧縮画像データのブロックデータの間に挿入されているので、量子化レベル係数QCxが復号化されたびに量子化テーブルQTが更新される。また、同一のパターンの色を有する連続した複数の画素ブロックを表わすデータは、特殊なデータ形式(ヌルランデータ)で圧縮画像データ中に挿入されている。



1

## 【特許請求の範囲】

【請求項1】 画像データを直交変換し、量子化し、さらにエントロピー符号化して得られた符号データを含む圧縮画像データを伸長する方法であって、

(A) 複数の画素ブロックに対する符号データと、前記複数の画素ブロックの先頭に位置する第1の画素ブロックに対する第1の係数コードと、任意の第2の画素ブロックに対する第2の係数コードとを含む圧縮画像データを準備する工程と、

(B) 前記符号データをエントロピー復号化することによって、量子化された変換係数を生成するとともに、前記第1と第2の係数コードをエントロピー復号化することによって第1と第2の係数を生成する工程と、

(C) 前記第1と第2の係数のそれと、所定の基本量子化テーブルの基本量子化レベルとを乗じることによって、第1と第2の量子化テーブルを作成する工程と、

(D) 前記第1の画素ブロックに対する前記量子化された変換係数を前記第1の量子化テーブルで逆量子化するとともに、前記第2の画素ブロックに対する前記量子化された変換係数を前記第2の量子化テーブルで逆量子化することによって、逆量子化された変換係数を求める工程と、

(E) 前記逆量子化された変換係数を逆直交変換することによって、伸長された画像データを求める工程と、を備えることを特徴とする圧縮画像データの伸長方法。

【請求項2】 請求項1記載の圧縮画像データの伸長方法であって、

前記符号データは、前記複数の画素ブロックの配列順序に従って配列された複数のデータユニットを含むとともに、前記第1の画素ブロックに対する第1のデータユニットの直前に前記第1の係数コードが配置され、前記第2の画素ブロックに対する第2のデータユニットの直前に前記第2の係数コードが配置されている、圧縮画像データの伸長方法。

【請求項3】 請求項2記載の圧縮画像データの伸長方法であって、

前記第1と第2の係数コードは、前記第1と第2の係数と前記変換係数の直流成分に対する複数のエントロピー符号語を含む符号テーブル内のエントロピー符号語である圧縮画像データの伸長方法。

【請求項4】 請求項1記載の圧縮画像データの伸長方法であって、前記工程(D)は、

前記第1の画素ブロックから、前記第2の画素ブロックの直前に配列された第3の画素ブロックまでの一連の画素ブロックに対して、前記第1の量子化テーブルを用いて逆量子化する工程、

を含む圧縮画像データの伸長方法。

【請求項5】 請求項1記載の圧縮画像データの伸長方法であって、前記工程(C)は、

前記第1と第2の係数のそれぞれと前記基本量子化テー

2

ブルに含まれている量子化レベルとの乗算結果が所定の最大値を超える場合には、前記乗算結果が前記最大値を超えた量子化レベルの値を前記最大値に等しく設定する工程、

を含む圧縮画像データの伸長方法。

【請求項6】 請求項1記載の圧縮画像データの伸長方法であって、前記工程(C)は、

前記基本量子化テーブル内の所定の基本量子化レベルについて、前記乗算を実質的に行なわずに、そのまま前記第1と第2の逆量子化テーブルの量子化レベルとする工程、

を含む圧縮画像データの伸長方法。

【請求項7】 請求項6記載の圧縮画像データの伸長方法であって、

前記所定の基本量子化レベルは、前記変換係数の直流成分に関する量子化レベルである、圧縮画像データの伸長方法。

【請求項8】 請求項1記載の圧縮画像データの伸長方法であって、

前記圧縮画像データは、さらに、同一の画像パターンを有する連続した複数の画素ブロックの個数を示す同一パターンブロックデータを含み、

さらに、前記工程(B)は、

前記同一パターンブロックデータで表わされた前記複数の画素ブロックに対する前記量子化された変換係数を作成する際に、前記変換係数の所定の成分を予め指定された値に設定するとともに、前記所定の成分以外の前記変換係数の成分の値をゼロに設定することによって作成する工程、

30 を有する圧縮画像データの伸長方法。

【請求項9】 請求項8記載の圧縮画像データの伸長方法であって、

前記所定の成分は直流成分である、圧縮画像データの伸長方法。

【請求項10】 請求項8記載の圧縮画像データの伸長方法であって、

前記工程(D)は、前記同一パターンブロックデータから復号化された前記変換係数については前記逆量子化を省略する工程、

40 を含む圧縮画像データの伸長方法。

【請求項11】 画像データを直交変換し、量子化し、さらにエントロピー符号化して得られた符号データを含む圧縮画像データを伸長する方法であって、

(A) 複数の第1の画素ブロックに対する符号データと、同一の画像パターンを有する連続した複数の第2の画素ブロックの個数を示す同一パターンブロックデータとを含む圧縮画像データを準備する工程と、

(B) 前記符号データをエントロピー復号化することによって前記複数の第1の画素ブロックに対する第1の量子化された変換係数を作成するとともに、前記同一バ

ターンブロックデータを含む圧縮画像データを伸長する方法。

ーンブロックデータで表わされた前記複数の第2の画素ブロックに対する第2の量子化された変換係数を作成する際に、前記変換係数の所定の成分を予め指定された値に設定するとともに、前記所定の成分以外の前記変換係数の成分の値をゼロに設定することによって前記第2の量子化された変換係数を作成する工程と、

(C) 前記第1と第2の量子化された変換係数を量子化テーブルを用いて逆量子化することによって、第1と第2の逆量子化された変換係数を求める工程と、

(D) 前記第1と第2の逆量子化された変換係数を逆直交変換することによって、伸長された画像データを作成する工程と、

を備えることを特徴とする圧縮画像データの伸長方法。

【請求項12】 画像データを直交変換し、量子化し、さらにエントロピー符号化して得られた符号データを含む圧縮画像データを伸長する装置であって、複数の画素ブロックに対する符号データと、前記複数の画素ブロックの先頭に位置する第1の画素ブロックに対する第1の係数コードと、任意の第2の画素ブロックに対する第2の係数コードとを含む圧縮画像データを記憶する記憶手段と、

前記符号データをエントロピー復号化することによって、量子化された変換係数を生成するとともに、前記第1と第2の係数コードをエントロピー復号化することによって第1と第2の係数を生成するエントロピー復号化手段と、

前記第1と第2の係数のそれぞれと、所定の基本量子化テーブルの基本量子化レベルとを乗することによって、第1と第2の量子化テーブルを作成する逆量子化テーブル作成手段と、

前記第1の画素ブロックに対する前記量子化された変換係数を前記第1の量子化テーブルで逆量子化するとともに、前記第2の画素ブロックに対する前記量子化された変換係数を前記第2の量子化テーブルで逆量子化することによって、逆量子化された変換係数を求める逆量子化手段と、

前記逆量子化された変換係数を逆直交変換することによって、伸長された画像データを求める逆直交変換手段と、

を備えることを特徴とする圧縮画像データの伸長方法。

【請求項13】 請求項12記載の圧縮画像データの伸長装置であって、

前記符号データは、前記複数の画素ブロックの配列順序に従って配列された複数のデータユニットを含むとともに、前記第1の画素ブロックに対する第1のデータユニットの直前に前記第1の係数コードが配置され、前記第2の画素ブロックに対する第2のデータユニットの直前に前記第2の係数コードが配置されている、圧縮画像データの伸長装置。

【請求項14】 請求項13記載の圧縮画像データの伸

長装置であって、

前記第1と第2の係数コードは、前記第1と第2の係数と前記変換係数の直流成分に対する複数のエントロピー符号語を含む符号テーブル内のエントロピー符号語である圧縮画像データの伸長装置。

【請求項15】 請求項12記載の圧縮画像データの伸長装置であって、前記逆量子化手段は、

前記第1の画素ブロックから、前記第2の画素ブロックの直前に配列された第3の画素ブロックまでの一連の画素ブロックに対して、前記第1の量子化テーブルを用いて逆量子化する手段、

を含む圧縮画像データの伸長装置。

【請求項16】 請求項12記載の圧縮画像データの伸長装置であって、前記逆量子化テーブル作成手段は、前記第1と第2の係数のそれぞれと前記基本量子化テーブルに含まれている量子化レベルとの乗算結果が所定の最大値を超える場合には、前記乗算結果が前記最大値を超えた量子化レベルの値を前記最大値に等しく設定するクリッピング手段、

を含む圧縮画像データの伸長装置。

【請求項17】 請求項12記載の圧縮画像データの伸長装置であって、前記逆量子化テーブル作成手段は、前記基本量子化テーブル内の所定の基本量子化レベルについて、前記乗算を実質的に行なわずに、そのまま前記第1と第2の逆量子化テーブルの量子化レベルとする手段、

を含む圧縮画像データの伸長装置。

【請求項18】 請求項17記載の圧縮画像データの伸長装置であって、

前記所定の基本量子化レベルは、前記変換係数の直流成分に関する量子化レベルである、圧縮画像データの伸長装置。

【請求項19】 請求項12記載の圧縮画像データの伸長装置であって、

前記圧縮画像データは、同一の画像パターンを有する連続した複数の画素ブロックの個数を示す同一パターンブロックデータを含み、

さらに、前記エントロピー復号化手段は、

前記同一パターンブロックデータで表わされた前記複数の画素ブロックに対する前記量子化された変換係数を作成する際に、前記変換係数の所定の成分を予め指定された値に設定するとともに、前記所定の成分以外の前記変換係数の成分の値をゼロに設定することによって作成する手段、

を有する圧縮画像データの伸長装置。

【請求項20】 請求項19記載の圧縮画像データの伸長装置であって、

前記所定の成分は直流成分である、圧縮画像データの伸長装置。

【請求項21】 請求項19記載の圧縮画像データの伸

長装置であって、

前記エントロピー復号化手段は、前記同一パターンブロックデータから復号化された前記変換係数を、前記逆量子化手段をバイパスして前記逆直交変換手段に供給する手段、

を含む圧縮画像データの伸長装置。

【請求項22】 画像データを直交変換し、量子化し、さらにエントロピー符号化して得られた符号データを含む圧縮画像データを伸長する装置であって、

複数の第1の画素ブロックに対する符号データと、同一の画像パターンを有する連続した複数の第2の画素ブロックの個数を示す同一パターンブロックデータとを含む圧縮画像データを記憶する第1の記憶手段と、前記複数の第2の画素ブロックに関する変換係数の所定の成分に対する指定値を記憶する第2の記憶手段と、前記符号データをエントロピー復号化することによって前記複数の第1の画素ブロックに対する第1の量子化された変換係数を作成するとともに、前記同一パターンブロックデータで表わされた前記複数の第2の画素ブロックに対する第2の量子化された変換係数を作成する際に、前記所定の成分を前記第2の記憶手段に記憶された指定値に設定するとともに前記所定の成分以外の前記変換係数の値をゼロに設定することによって第2の量子化された変換係数を作成するエントロピー復号化手段と、前記第1と第2の量子化された変換係数を量子化テーブルを用いて逆量子化することによって、第1と第2の逆量子化された変換係数を求める逆量子化手段と、

前記第1と第2の逆量子化された変換係数を逆直交変換することによって、伸長された画像データを作成する逆直交変換手段と、

を備えることを特徴とする圧縮画像データの伸長装置。

【請求項23】 画像データの圧縮方法であって、

(A) 画像データを画像内の複数の画素ブロック毎に直交変換して変換係数を求める工程と、

(B) 前記複数の画素ブロックの先頭に位置する第1の画素ブロックに対する第1の係数と、任意の第2の画素ブロックに対する第2の係数とを指定する工程と、

(C) 前記第1と第2の係数のそれと所定の基本量子化テーブルの基本量子化レベルとを乗ずることによって、第1と第2の量子化テーブルを作成する工程と、

(D) 前記第1と第2の画素ブロックに対する前記変換係数を前記第1と第2の量子化テーブルでそれぞれ量子化することによって、量子化された変換係数を求める工程と、

(E) 前記量子化された変換係数をエントロピー符号化することによって符号データを作成するとともに、前記第1と第2の係数をエントロピー符号化することによって第1と第2の係数コードを作成する工程と、

(F) 前記符号データと前記第1と第2の係数コードとを含む圧縮画像データを作成する工程と、

を備えることを特徴とする画像データの圧縮方法。

【請求項24】 請求項23記載の画像データの圧縮方法であって、

前記符号データは、前記複数の画素ブロックの配列順序に従って配列された複数のデータユニットを含むとともに、前記第1の画素ブロックに対する第1のデータユニットの直前に前記第1の係数コードを含み、前記第2の画素ブロックに対する第2のデータユニットの直前に前記第2の係数コードを含む、

画像データの圧縮方法。

【請求項25】 請求項24記載の画像データの圧縮方法であって、前記工程(E)は、

前記第1と第2の係数と前記変換係数の直流成分に対する複数のエントロピー符号語を含む符号テーブルを用いて、前記第1と第2の係数をエントロピー符号化する工程、

を含む画像データの圧縮方法。

【請求項26】 請求項23記載の画像データの圧縮方法であって、前記工程(E)は、同一の画像パターンを

有する連続した複数の画素ブロックの個数を示す同一パターンブロックデータを作成する工程を含み、

前記工程(F)は、前記符号データおよび前記第1と第2の係数コードとともに、前記同一パターンブロックデータを含む圧縮画像データを作成する工程を含む画像データの圧縮方法。

【請求項27】 画像データを直交変換し、量子化し、さらにエントロピー符号化して得られた符号データを含む圧縮画像データを伸長する方法であって、

(A) 複数の画素ブロックに対する符号データと、前記複数の画素ブロックのうちの任意の第1の画素ブロックに対する第1の係数コードとを含む圧縮画像データを準備する工程と、

(B) 前記符号データをエントロピー復号化することによって、量子化された変換係数を生成するとともに、前記第1の係数コードをエントロピー復号化することによって第1の係数を生成する工程と、

(C) 前記第1の係数と、所定の基本量子化テーブルの基本量子化レベルとを乗ずることによって、第1の量子化テーブルを作成する工程と、

(D) 前記第1の画素ブロックを含む少なくとも1つの画素ブロックに対する前記量子化された変換係数を前記第1の量子化テーブルで逆量子化することによって、逆量子化された変換係数を求める工程と、

(E) 前記逆量子化された変換係数を逆直交変換することによって、伸長された画像データを求める工程と、を備えることを特徴とする圧縮画像データの伸長方法。

【請求項28】 請求項27記載の圧縮画像データの伸長方法であって、

前記圧縮画像データは、さらに、前記第1の画素ブロック以降に配列された第2の画素ブロックに対する第2の

係数コードを含み、

前記工程（B）は、前記第2の係数コードをエントロピー復号化することによって第2の係数を生成する工程を含み、

前記工程（C）は、前記第2の係数と前記基本量子化テーブルの基本量子化レベルとを乗することによって、第2の量子化テーブルを作成する工程を含み、

前記工程（D）は、前記第1の画素ブロックから前記第2の画素ブロックの直前に配列された第3の画素ブロックまでの一連の画素ブロックに対する前記量子化された変換係数を、前記第1の量子化テーブルを用いて逆量子化するとともに、前記第2の画素ブロックを含む少なくとも1つの画素ブロックに対する前記量子化された変換係数を前記第2の量子化テーブルで逆量子化する工程、を含む圧縮画像データの伸長方法。

【請求項29】 請求項28記載の圧縮画像データの伸長方法であって、

前記符号データは、前記複数の画素ブロックの配列順序に従って配列された複数のデータユニットを含むとともに、前記第1の画素ブロックに対する第1のデータユニットの直前に前記第1の係数コードが配置され、前記第2の画素ブロックに対する第2のデータユニットの直前に前記第2の係数コードが配置されている、圧縮画像データの伸長方法。

#### 【発明の詳細な説明】

##### 【0001】

【産業上の利用分野】この発明は、画像データを圧縮・伸長する方法およびそのための装置に関する。

##### 【0002】

【従来の技術】図31は、従来の画像データの圧縮／伸長装置の構成を示すブロック図である。画像データ圧縮装置540は、直交変換部542において原画像データをM×N画素のブロック毎に直交変換した後、量子化部544において量子化を行ない、さらに、エントロピー符号化部546において符号化を行なって圧縮画像データを作成する。一方、画像データ伸長装置550は、まずエントロピー復号化部556において圧縮画像データを復号化し、逆量子化部554で逆量子化した後に、逆直交変換部552によって画像データを復元する。なお、量子化部544と逆量子化部554は同一の量子化テーブル562を使用し、また、エントロピー符号化部546とエントロピー復号化部556も同一の符号テーブル564を使用する。

##### 【0003】

【発明が解決しようとする課題】ところで、1つの画像内に、高画質で復元したい第1の部分と低画質で復元してもよい第2の部分とを含む場合がある。このような場合には、第1の部分に対しては量子化レベルの小さな量子化テーブルを用いて量子化を行ない、第2の部分に対しては量子化レベルの大きな量子化テーブルを用いて量

子化を行なえば良い。量子化テーブルは、画像データの画素ブロックと同じサイズの行列、すなわちM行N列の行列である。従来の画像データ圧縮／伸長装置では、1つの画像内において複数の量子化テーブルを使用する場合には、M行N列の複数の量子化テーブルを圧縮装置540から伸長装置550に転送しなければならず、圧縮画像データのデータ量を増大させるという問題があった。

【0004】また、画像の中に一様な色を有する画素ブロックが連続するような単純な画像部分についても、それらの画素ブロックに対するM行N列の直交変換係数をすべてエントロピー符号化する。従って、このような単純な画像部分を表す圧縮画像データも、かなりのデータ量になるという問題があった。

【0005】この発明は、従来技術における上述の課題を解決するためになされたものであり、圧縮画像データのデータ量を従来に比べて低減することのできる技術を提供することを目的とする。

##### 【0006】

【課題を解決するための手段および作用】上述の課題を解決するため、この発明による圧縮画像データ伸長方法は、(A)複数の画素ブロックに対する符号データと、前記複数の画素ブロックの先頭に位置する第1の画素ブロックに対する第1の係数コードと、任意の第2の画素ブロックに対する第2の係数コードとを含む圧縮画像データを準備する工程と、(B)前記符号データをエントロピー復号化することによって、量子化された変換係数を生成するとともに、前記第1と第2の係数コードをエントロピー復号化することによって第1と第2の係数を生成する工程と、(C)前記第1と第2の係数のそれぞれと、所定の基本量子化テーブルの基本量子化レベルとを乗ることによって、第1と第2の量子化テーブルを作成する工程と、(D)前記第1の画素ブロックに対する前記量子化された変換係数を前記第1の量子化テーブルで逆量子化するとともに、前記第2の画素ブロックに対する前記量子化された変換係数を前記第2の量子化テーブルで逆量子化することによって、逆量子化された変換係数を求める工程と、(E)前記逆量子化された変換係数を逆直交変換することによって、伸長された画像データを求める工程と、を備える。

【0007】第1と第2の係数を与えるだけで量子化レベルの異なる第1と第2の量子化テーブルを作成できるので、2つの量子化テーブルを圧縮画像データに含める場合に比べて圧縮画像データのデータ量を低減することができる。

【0008】前記符号データは、前記複数の画素ブロックの配列順序に従って配列された複数のデータユニットを含むとともに、前記第1の画素ブロックに対する第1のデータユニットの直前に前記第1の係数コードが配置され、前記第2の画素ブロックに対する第2のデータユ

ニットの直前に前記第2の係数コードが配置されている  
ようにしてもよい。

【0009】こうすれば、第1と第2の係数を第1と第2の画素ブロックにそれぞれ容易に対応づけられる。

【0010】また、前記第1と第2の係数コードは、前記第1と第2の係数と前記変換係数の直流成分に対する複数のエントロピー符号語を含む符号テーブル内のエントロピー符号語であることが好ましい。

【0011】こうすれば、第1と第2の係数を一意にかつ瞬時に復号することができる。

【0012】前記工程(D)は、前記第1の画素ブロックから、前記第2の画素ブロックの直前に配列された第3の画素ブロックまでの一連の画素ブロックに対して、前記第1の量子化テーブルを用いて逆量子化する工程、を含むようとするのが好ましい。

【0013】第2の係数が現われるまでの複数の画素ブロックに対して第1の係数が適用されるので、これらの複数の画素ブロックのそれぞれの符号データの前に第1の係数を含めておく必要がない。

[0014] 前記工程(C)は、前記第1と第2の係数のそれぞれと前記基本量子化テーブルに含まれている量子化レベルとの乗算結果が所定の最大値を超える場合には、前記乗算結果が前記最大値を超えた量子化レベルの値を前記最大値に等しく設定する工程、を含むようにしてもよい。

【0015】こうすれば、量子化レベルのビット数を一定値以下に保つことができる。

【0016】前記工程(C)は、前記基本量子化テーブル内の所定の基本量子化レベルについては、前記乗算を実質的に行なわずに、そのまま前記第1と第2の逆量子化テーブルの量子化レベルとする工程、を含むようにしてもよい。

〔0017〕こうすれば、所定の位置の変換係数を逆量子化する際に発生する誤差を低減することができる。

【0018】前記所定の基本量子化レベルは、前記変換係数の直流成分に関する量子化レベルであるようにするのが好ましい。

【0019】変換係数の直流成分に対する量子化レベルを変えないようにすれば、直流成分の逆量子化にともなう誤差を低減することができる。

【0020】前記圧縮画像データは、さらに、同一の画像パターンを有する連続した複数の画素ブロックの個数を示す同一パターンブロックデータを含み、さらに、前記工程（B）は、前記同一パターンブロックデータで表わされた前記複数の画素ブロックに対する前記量子化された変換係数を作成する際に、前記変換係数の所定の成分を予め指定された値に設定するとともに、前記所定の成分以外の前記変換係数の成分の値をゼロに設定することによって作成する工程、を有するようしてもよい。

【0021】こうすれば、互いに同一の画像パターンを

有する連続した複数の画素ブロックを比較的少ないデータで表わすことができる。

〔0022〕前記所定の成分は直流成分であることが好ましい。

【0023】こうすれば、連続した複数の画素ブロックを同一の色で塗ることができる。

【0024】前記工程(D)は、前記同一パターンプロックデータから復号化された前記変換係数については前記逆量子化を省略する工程、を含むようにしてもよい。

10 【0025】こうすれば、逆量子化による誤差の発生を防止することができる。

【0026】この発明による他の圧縮画像データ伸長方法は、(A)複数の第1の画素ブロックに対する符号データと、同一の画像パターンを有する連続した複数の第2の画素ブロックの個数を示す同一パターンブロックデータとを含む圧縮画像データを準備する工程と、(B)前記符号データをエンタロピー復号化することによって前記複数の第1の画素ブロックに対する第1の量子化された変換係数を作成するとともに、前記同一パターンブ

20 ロックデータで表わされた前記複数の第2の画素ブロックに対する第2の量子化された変換係数を作成する際に、前記変換係数の所定の成分を予め指定された値に設定するとともに、前記所定の成分以外の前記変換係数の成分の値をゼロに設定することによって前記第2の量子化された変換係数を作成する工程と、(C) 前記第1と第2の量子化された変換係数を量子化テーブルを用いて逆量子化することによって、第1と第2の逆量子化された変換係数を求める工程と、(D) 前記第1と第2の逆量子化された変換係数を逆直交変換することによって、

30 伸長された画像データを作成する工程とを備える。

【0027】この方法では、互いに同一のパターンの色を有する連続した複数の画素ブロックを比較的少ないデータで表わすことができる。

【0028】この発明による圧縮画像データ伸長装置は、複数の画素ブロックに対する符号データと、前記複数の画素ブロックの先頭に位置する第1の画素ブロックに対する第1の係数コードと、任意の第2の画素ブロックに対する第2の係数コードとを含む圧縮画像データを記憶する記憶手段と、前記符号データをエンタロピー復

40 号化することによって、量子化された変換係数を生成するとともに、前記第1と第2の係数コードをエントロピー復号化することによって第1と第2の係数を生成するエントロピー復号化手段と、前記第1と第2の係数のそれと、所定の基本量子化テーブルの基本量子化レベルとを乗することによって、第1と第2の量子化テーブルを作成する逆量子化テーブル作成手段と、前記第1の画素ブロックに対する前記量子化された変換係数を前記第1の量子化テーブルで逆量子化するとともに、前記第2の画素ブロックに対する前記量子化された変換係数を前記第2の量子化テーブルで逆量子化することによつ

50

て、逆量子化された変換係数を求める逆量子化手段と、前記逆量子化された変換係数を逆直交変換することによって、伸長された画像データを求める逆直交変換手段と、を備える。

【0029】また、この発明による他の圧縮画像データ伸長装置は、複数の第1の画素ブロックに対する符号データと、同一の画像パターンを有する連続した複数の第2の画素ブロックの個数を示す同一バータンブロックデータとを含む圧縮画像データを記憶する第1の記憶手段と、前記複数の第2の画素ブロックに関する変換係数の所定の成分に対する指定値を記憶する第2の記憶手段と、前記符号データをエントロピー復号化することによって前記複数の第1の画素ブロックに対する第1の量子化された変換係数を作成するとともに、前記同一バータンブロックデータで表わされた前記複数の第2の画素ブロックに対する第2の量子化された変換係数を作成する際に、前記所定の成分を前記第2の記憶手段に記憶された指定値に設定するとともに前記所定の成分以外の前記変換係数の値をゼロに設定することによって第2の量子化された変換係数を作成するエントロピー復号化手段と、前記第1と第2の量子化された変換係数を量子化テーブルを用いて逆量子化することによって、第1と第2の逆量子化された変換係数を求める逆量子化手段と、前記第1と第2の逆量子化された変換係数を逆直交変換することによって、伸長された画像データを作成する逆直交変換手段と、を備える。

【0030】この発明による画像データの圧縮方法は、(A) 画像データを画像内の複数の画素ブロック毎に直交変換して変換係数を求める工程と、(B) 前記複数の画素ブロックの先頭に位置する第1の画素ブロックに対する第1の係数と、任意の第2の画素ブロックに対する第2の係数とを指定する工程と、(C) 前記第1と第2の係数のそれぞれと所定の基本量子化テーブルの基本量子化レベルとを乗することによって、第1と第2の量子化テーブルを作成する工程と、(D) 前記第1と第2の画素ブロックに対する前記変換係数を前記第1と第2の量子化テーブルでそれぞれ量子化することによって、量子化された変換係数を求める工程と、(E) 前記量子化された変換係数をエントロピー符号化することによって符号データを作成するとともに、前記第1と第2の係数をエントロピー符号化することによって第1と第2の係数コードを作成する工程と、(F) 前記符号データと前記第1と第2の係数コードとを含む圧縮画像データを作成する工程とを備える。

【0031】第1と第2の係数を与えるだけで量子化レベルの異なる第1と第2の量子化テーブルを作成できるので、2つの量子化テーブルを圧縮画像データに含める場合に比べて圧縮画像データのデータ量を低減することができる。

【0032】この発明によるさらに他の圧縮画像データ

を伸長する方法は、(A) 複数の画素ブロックに対する符号データと、前記複数の画素ブロックのうちの任意の第1の画素ブロックに対する第1の係数コードとを含む圧縮画像データを準備する工程と、(B) 前記符号データをエントロピー復号化することによって、量子化された変換係数を生成するとともに、前記第1の係数コードをエントロピー復号化することによって第1の係数を生成する工程と、(C) 前記第1の係数と、所定の基本量子化テーブルの基本量子化レベルとを乗ることによって、第1の量子化テーブルを作成する工程と、(D) 前記第1の画素ブロックを含む少なくとも1つの画素ブロックに対する前記量子化された変換係数を前記第1の量子化テーブルで逆量子化することによって、逆量子化された変換係数を求める工程と、(E) 前記逆量子化された変換係数を逆直交変換することによって、伸長された画像データを求める工程と、を備える。

【0033】こうすれば、第1と第2の係数を与えるだけで量子化レベルの異なる第1と第2の量子化テーブルを作成できるので、2つの量子化テーブルを圧縮画像データに含める場合に比べて圧縮画像データのデータ量を低減することができる

【0034】前記圧縮画像データは、さらに、前記第1の画素ブロック以降に配列された第2の画素ブロックに対する第2の係数コードを含み、前記工程(B)は、前記第2の係数コードをエントロピー復号化することによって第2の係数を生成する工程を含み、前記工程(C)は、前記第2の係数と前記基本量子化テーブルの基本量子化レベルとを乗ることによって、第2の量子化テーブルを作成する工程を含み、前記工程(D)は、前記第1の画素ブロックから前記第2の画素ブロックの直前に配列された第3の画素ブロックまでの一連の画素ブロックに対する前記量子化された変換係数を、前記第1の量子化テーブルを用いて逆量子化するとともに、前記第2の画素ブロックを含む少なくとも1つの画素ブロックに対する前記量子化された変換係数を前記第2の量子化テーブルで逆量子化する工程、を含むようになるのが好ましい。

【0035】こうすれば、圧縮画像データに第1と第2の係数コードを含めておくだけで、圧縮画像データの伸長時に、2つの量子化テーブルを用いて量子化を行なうことができる。

【0036】

【実施例】以下では、次の各項目について順次説明を行なう。

- A. 圧縮／伸長装置の全体構成と基本動作；
- B. 量子化レベル係数QCxによる量子化テーブルQTの調整；
- C. ハフマン符号化と圧縮データの構成；
- D. 逆量子化テーブル作成部の詳細構成；
- E. ヌルランデータの復号化

【0037】A. 圧縮／伸長装置の構成と動作：図1は、この発明の一実施例を適用した画像データの圧縮装置100と伸長装置200の機能を示すブロック図である。

【0038】画像データ圧縮装置100は、原画像データ $f(x, y)$ に対してディスクリートコサイン変換を行なうDCT部110と、DCT変換で得られた変換係数 $F(u, v)$ を量子化する量子化部120と、量子化された変換係数 $QF(u, v)$ をハフマン符号化して圧縮画像データ $ZZ$ を作成するハフマン符号化部130と、量子化テーブル作成部140と、ハフマン符号テーブルメモリ150とを備えている。量子化テーブル作成部140は、後述するように、基本量子化テーブルBQTと量子化レベル係数QCxに基づいて量子化テーブルQ $T$ を作成する。圧縮画像データ $ZZ$ は、CD-ROMなどの記憶媒体に記憶されて、画像データ圧縮装置100から画像データ伸長装置200に供給される。

【0039】画像データ伸長装置200は、圧縮画像データ $ZZ$ をハフマン復号化するハフマン復号化部210と、復号された量子化後の変換係数 $QF(u, v)$ を逆量子化する逆量子化部220と、逆量子化された変換係数 $FF(u, v)$ にディスクリートコサイン逆変換を行なって画像データ $f'f'(x, y)$ を得るIDCT部230と、ハフマン符号テーブルメモリ240と、逆量子化テーブル作成部250とを備えている。逆量子化テーブル作成部250は、圧縮画像データ $ZZ$ から復号された基本量子化テーブルBQTと量子化レベル係数QCxとをハフマン復号化部210から受取り、これらに基づいて量子化テーブルQ $T$ を作成する。この量子化テーブルQ $T$ は、圧縮装置100で用いられた量子化テーブルQ $T$ と同じである。また、ハフマン符号テーブルメモリ240に記憶されているハフマン符号テーブルHTも圧縮装置100のハフマン符号テーブルメモリ150に記憶されているものと同じである。

【0040】図2は、画像データ圧縮装置100の具体的な構成を示すブロック図である。この画像データ圧縮装置100は、CPU101と、メインメモリ102と、キーボード103と、マウス104と、磁気ディスク装置105と、光磁気ディスク装置106とを備えている。図1に示す画像データ圧縮装置100の各処理部110～140は、メインメモリ102に記憶されたソフトウェアプログラムによって実現されている。また、基本量子化テーブルBQTとハフマン符号テーブルHTは、磁気ディスク装置105に記憶されている。この画像データ圧縮装置100は、ビデオゲームを作成するためのワークステーションであり、画像データ圧縮のためのソフトウェアプログラムの他に、ビデオゲームを作成するための各種のプログラムがCPU101により実行される。完成したゲームプログラムは、圧縮画像データ

とともに光磁気ディスク装置106に格納される。そして、この光磁気ディスクを用いて、ビデオゲームのプログラムと圧縮画像データを含むCD-ROMが製作される。

【0041】図3は、画像データ伸長装置200を含むビデオゲーム装置20の構成を示すブロック図である。このビデオゲーム装置20は、SCSIバス36を介して接続されたCD-ROMドライブ32の他、画像処理とこれに関する総ての処理を統括的に司るマイクロプロセッサ（以下、MPUという）40と、このMPU40に直接接続されたメインメモリ（以下、M-RAMと呼ぶ）41と、同じくBIOSプログラムを記憶したROM42と、MPU40のバス（M-BUS）43に接続された各種ユニット、即ち画像信号コントロールユニット45、画像データ伸長ユニット200、特定の画像信号を出力するVDPユニット49、ビデオ信号の合成と出力を行なうビデオエンコーダユニット50、音声データを扱う音声データ出力ユニット52とを備える。

【0042】また、このビデオゲーム装置20内には、画像信号コントロールユニット45のローカルバス（K-BUS）54に接続されたメモリ（以下、K-RAMと呼ぶ）55、画像データ伸長ユニット200のローカルバスに接続されたメモリ（以下、R-RAMと呼ぶ）251、VDPユニット49のローカルバスに接続されたビデオメモリ（以下、V-RAMと呼ぶ）59、ビデオエンコーダユニット50からの出力信号を通常の映像信号（NTSC）に変換し、カラーテレビ28に出力するNTSCコンバータ60が備えられている。

【0043】画像信号コントロールユニット45、画像データ伸長ユニット200、ビデオエンコーダユニット50および音声データ出力ユニット52は、それぞれロジック回路により構成されている。

【0044】図4（A）は、ゲームの背景画像となる原画像の一例を示す平面図である。この原画像は、一様な色で塗られた背景BGに火山の自然画がはめ込まれた画像である。図4（B）は、1つの画素ブロックPBを含む原画像の一部を拡大して示している。一般に、画素ブロックPBは、 $M \times N$ 個の画素PXを含むように設定できる。2つの整数M、Nの値としては8または16が好ましく、この実施例では、 $M=N=8$ である。なお、整数MとNを異なる値に設定しても良い。後述するように、圧縮画像データZZの中で背景BGを表すデータ部分は、一様色の画素ブロックPBが連続することを示す特別なデータ形式（ヌルランデータ）を有している。

【0045】画像データ圧縮装置100のDCT部110は、次の式1に従って、各画素ブロックPB毎に2次元DCT変換を行なう。

【数1】

15

$$F(u, v) = \frac{1}{4} C(u)C(v) \sum_{x=0}^7 \sum_{y=0}^7 f(x, y) \cos \frac{\pi(2x+1)u}{16} \cos \frac{\pi(2y+1)v}{16}$$

$$C(u), C(v) = \frac{1}{\sqrt{2}} \quad (u, v = 0 \text{ の時})$$

$$= 1 \quad (u, v \neq 0 \text{ の時})$$

ここで、 $f(x, y)$  は 1 つの画素ブロック PB に含まれる  $8 \times 8$  個の画像データの配列、 $x, y$  は各画素ブロック PB 内の各画素の位置を示す座標、 $F(u, v)$  は変換係数の配列、 $u, v$  は周波数空間の座標である。

【0046】図5は、変換係数 $F(u, v)$  の配列を示す説明図である。変換係数 $F(u, v)$  は画素ブロック PB と同じ  $8 \times 8$  の配列である。左上端の変換係数 $F(0, 0)$  は DC 成分（または DC 係数）と呼ばれており、その他の変換係数は AC 成分（または AC 係数）と呼ばれている。DC 成分は、画素ブロック PB における画像データの平均値を示している。また、AC 成分は、画素ブロック PB における画像データの変化を示している。隣接する画素の画像データにはある程度の相関があるので、AC 係数の中で低周波成分の値は比較的大きく、高周波成分の値は比較的小さい。また、高周波成分が画質に与える影響は比較的小さい。

【0047】図6は、画像データ圧縮装置 100 と画像データ伸長装置 200 の基本動作を示す説明図である。DCT 部 110 は、図6(a) に示す DCT 係数 $F(u, v)$  を生成する。

【0048】量子化テーブル作成部 140 は、次の式 2 に示すように、基本量子化テーブル BQT (図6(c)) と量子化レベル係数 QCx とを乗することによって量子化テーブル QT (図6(d)) を作成する。

【数2】

$$QT(u, v) = QCx \times BQT(u, v)$$

【0049】図6の例では QCx = 1 なので、量子化テーブル QT は基本量子化テーブル BQT と同一である。

【0050】量子化部 120 は、DCT 係数 $F(u, v)$  を量子化テーブル QT で線形量子化することによって、図6(b) に示す量子化された DCT 係数 $FF(u, v)$  を求める。線形量子化とは、除算を行なつ \*

$$ff(x, y) = \frac{1}{4} \sum_{u=0}^7 \sum_{v=0}^7 C(u)C(v) FF(u, v) \cos \frac{\pi(2x+1)u}{16} \cos \frac{\pi(2y+1)v}{16}$$

$$C(u), C(v) = \frac{1}{\sqrt{2}} \quad (u, v = 0 \text{ の時})$$

$$= 1 \quad (u, v \neq 0 \text{ の時})$$

【0055】B. 量子化レベル係数 QCx による量子化テーブル QT の調整：量子化テーブル QT は、前記式 2 に従って基本量子化テーブル BQT と量子化レベル係数 QCx とを乗算することによって作成されるので、量子化レベル係数 QCx の値を大きくすれば量子化テーブル QT 内の各量子化レベルを大きくすることができます。

16

\* て、その除算結果を整数に丸める処理である。  
【0051】ハフマン符号化部 130 は、この DCT 係数 $QF(u, v)$  をハフマン符号化することによって圧縮画像データ ZZ (図6(e)) を作成する。なお、ハフマン符号化の方法については更に後述する。圧縮画像データ ZZ は、後述するように、基本量子化テーブル BQT を表わす第 1 のデータと、量子化レベル係数 QCx と変換係数 $QF(u, v)$  を表わす第 2 のデータとを含んでいる。

【0052】圧縮画像データ ZZ が画像データ伸長装置 200 に与えられると、ハフマン復号化部 210 が圧縮画像データ ZZ を復号化して DCT 係数 $QF(u, v)$  (図6(f)) を求める。ハフマン符号化は可逆符号化なので、この DCT 係数 $QF(u, v)$  は、画像データ圧縮装置 100 の量子化部 120 によって求められた量子化後の DCT 係数 $QF(u, v)$  (図6(b)) と同一である。なお、ハフマン復号化部 210 は、DCT 係数 $QF(u, v)$  の他に、圧縮画像データ ZZ に含まれている基本量子化テーブル BQT (図6(c)) と量子化レベル係数 QCx も復号化して逆量子化テーブル作成部 250 に与える。

【0053】逆量子化テーブル作成部 250 は、基本量子化テーブル BQT と量子化レベル係数 QCx とを乗算することによって量子化テーブル QT (図6(d)) を作成する。逆量子化部 220 は、この量子化テーブル QT と DCT 係数 $QF(u, v)$  とを乗算し、図6(g) に示す復号された DCT 係数 $FF(u, v)$  を求める。

【0054】IDCT 部 230 は、この DCT 係数 $FF(u, v)$  に対して次の式 3 に示す 2 次元 DCT 逆変換を行ない、復元された画像データ $f(x, y)$  を生成する。

【数3】

量子化レベル係数 QCx の値は、画像データ圧縮装置 100 において画像データを圧縮する際に、予め定められた複数の値 (0 ~ 15) の中からオペレータが選択する。

【0056】図7は、量子化レベル係数 QCx を 4 に指定した場合の圧縮／伸長動作を示す説明図である。量子

化テーブル作成部140と逆量子化テーブル作成部250は、上記の数式2に従って、図7(d)に示す量子化テーブルQTを作成する。但し、この実施例では量子化レベルの最大値が15に制限されており、乗算の結果が15以上になる量子化レベルの値はすべて強制的に15に設定される。

【0057】図7(d)に示す量子化テーブルQTを用いて図7(a)のDCT係数F(u, v)を線形量子化すると、図7(b)に示すように、DC成分が1であり、AC成分がすべて0であるDCT係数QF(u, v)が得られる。このように、量子化レベル係数QCxの値を大きくすると、量子化されたDCT係数QF(u, v)における0の数が増加するので、データ圧縮率を高めることができる。但し、図7(g)に示す復号されたDCT係数FF(u, v)は、図7(a)に示す元のDCT係数F(u, v)とかなり異なる値を示しているので、量子化レベル係数QCxが1に等しい場合(図6)よりも画質の劣化は大きい。

【0058】ところで、DCT係数のDC成分は画素ブロックPB内における画像データの平均値を示しているので、画質に対する影響がかなり大きい。従って、量子化レベル係数QCxの値に係わらずに、DC成分用の量子化レベルを基本量子化テーブルBQTにおける値と同じに保つようにするのが好ましい。図8は、このような場合の処理を示す説明図であり、DC成分用の量子化レベルQT(0, 0)が強制的に1に保たれている。QT(0, 0)=1の場合には、図8(g)に示す復号されたDCT係数FF(u, v)のDC成分が、図8(a)に示す元のDCT係数F(u, v)のDC成分と同じ値に保たれるので、図7の場合と同程度の圧縮率で画質の劣化を比較的小さく抑えることができる。なお、DC成分用の量子化レベルQT(0, 0)は必ずしも1にする必要はない、他の任意の値を設定しておいてもよい。

【0059】量子化レベル係数QCxとしては0を選択することも可能である。QCx=0の場合には、図9に示すように、量子化テーブルQT内のすべての量子化レベルが1に設定される。量子化されたDCT係数QF(u, v)は元のDCT係数F(u, v)と同じであるので、圧縮率は小さいが高画質で圧縮/伸長を行なうことができる。

【0060】以上をまとめると、量子化テーブル作成部140と逆量子化テーブル作成部250は、次のような特徴を有している。

(1) 基本量子化テーブルBQTと量子化レベル係数QCxとを、数式2に従って乗算することによって量子化テーブルQTを作成する。

(2) 乗算の結果が量子化レベルの最大値(=15)以上のものは、最大値に等しく設定する(図6~図8)。

(3) DC成分用の量子化レベルQT(0, 0)は、量子化レベル係数QCxの値に係わらず基本量子化テーブ

ルBQTにおける値と一緒に保つ(図8)。

(4) 量子化レベル係数QCxが0の場合には、すべての量子化レベルを1に設定する(図9)。

【0061】量子化テーブルQTを求める上述の演算は、画像データ圧縮装置100においてはソフトウェアプログラムによって行なわれ、画像データ伸長装置200においては専用のハードウェアで構成された逆量子化テーブル作成部250によって行なわれる。逆量子化テーブル作成部250の具体的な回路構成については更に後述する。

【0062】なお、量子化レベル係数QCxは、画像データ圧縮装置100で画像データを圧縮する際に、キーボード103やマウス104を用いて画素ブロックPBごとに異なる値を指定することができる。

【0063】C. ハフマン符号化と圧縮データの構成: 画像データ圧縮装置100のハフマン符号化部130

(図1)は、DC係数符号化部とAC係数符号化部とで構成されている。図10(A)は、DC係数符号化部の機能を示すブロック図である。ブロック遅延部131と加算器132は、図10(B)に示すように、各画素ブロックPBのDC係数DC<sub>i</sub>と1つ前の画素ブロックPBのDC係数DC<sub>i-1</sub>との差分△DCを算出する。

【0064】カテゴリ化処理部133は、図11に示すカテゴリ化テーブルに従って、DC係数の差分△DCに対応するカテゴリSSSSと識別データIDとを求める。カテゴリSSSSは、DC係数の差分△DCの範囲を示す番号である。識別データIDは、カテゴリSSSSで指定される複数の差分△DCの中の小さい方から何番目の値であるかを示すデータである。

【0065】カテゴリSSSSは、さらに1次元ハフマン符号化部134(図10)においてDC係数用のハフマン符号語HFDに変換される。図12は、1次元ハフマン符号化部134によって使用されるハフマン符号テーブルHTDCの一例を示す説明図である。この実施例では、原画像データf(x, y)がYUV信号(輝度信号Yと2つの色差信号U, V)で表現されているものとする。U信号/V信号共用のDC係数用ハフマン符号テーブルは、0~9のカテゴリSSSSの符号語を含むだけである。一方、Y信号用のDC係数用ハフマン符号テーブルは、0~9のカテゴリSSSSの符号語の他に、15~31のカテゴリSSSSの符号語を含んでいる。SSSS=15のハフマン符号語は、後述するヌルランデータであることを示している。ヌルランデータは、一様色の画素ブロックPBが連続することを示すデータである。また、SSSS=16~31のハフマン符号語は、量子化レベル係数QCxの値を示す符号である。例えば、SSSS=16に対するハフマン符号語「11111000」はQCx=0を示しており、SSSS=31に対するハフマン符号語「11111111」はQCx=15を示している。なお、図12のハフマン符号語は、カテゴリS

$SSS = 1 \sim 9$ 、および $15 \sim 31$ のすべてに関して一意復号可能で、かつ、瞬時復号可能である。

【0066】図13は、ハフマン符号化部130内のAC係数符号化部の機能を示すブロック図である。AC係数の配列F(u, v) ( $u = v = 0$ を除く)は、まずシグザグスキャン部135によって1次元に並び直される。図14は、シグザグスキャンの順路を示す説明図である。

【0067】判定部136は、1次元に並び直されたAC係数の値が0か否かを判定する。AC係数の値が0であれば、ランレンジスカウンタ137が、連続する0のAC係数をゼロラン長NNNNに変換する。AC係数が0でなければ、そのAC係数の値がカテゴリ化部138によってカテゴリSSSSと識別データIDに変換される。この際、図11に示すカテゴリ化テーブルが参照される。

【0068】ゼロラン長NNNNとカテゴリSSSSとは、2次元ハフマン符号化部139においてAC係数用のハフマン符号語HFACに変換される。図15は、AC係数用の2次元ハフマン符号テーブルHTACを示す説明図である。また、図16は、ハフマン符号テーブルHTACの中で、 $NNNN = 0$ と $NNNN = 1$ の部分(図15における最上部2行)のハフマン符号語の一例を示している。なお、 $NNNN/SSSS = 0/0$ のハフマン符号語「1111」は、1つの画素ブロックに対する符号データの終了を示している。

【0069】図17は、ハフマン符号化の一例を示す説明図である。図17(B)は、DC係数の符号化を示している。1つ前の画素ブロックにおけるDC係数の値を0と仮定すると、 $\Delta DC = F(0, 0) = 12$ である。図11のカテゴリ化テーブルによれば $\Delta DC = 12$ のカテゴリSSSSは4であり、識別データIDは「1100」である。また、図12のDC係数用ハフマン符号テーブルによれば、カテゴリSSSS=4のハフマン符号語HFDCは「011」である。なお、ここではY信号用のハフマン符号テーブルを使用する。DC係数に対するハフマン符号(HF+ID)は、図17(B)に示すように「0111100」となる。

【0070】図17(C)はAC係数の符号化を示している。まず、シグザグスキャンによって、AC係数が一次元の配列に並べられる。この配列は、ゼロラン長NNNNと、ゼロでない値のカテゴリSSSS(図11参照)とに変換される。ゼロラン長NNNNとカテゴリSSSSの組み合わせは、図15および図16に示すAC係数用ハフマン符号テーブルによってハフマン符号語HFACに変換され、ゼロでないAC係数の識別データIDと組み合わされて、図17(C)に示すようにハフマン符号(HFAC+ID)が作成される。

【0071】図18は、圧縮データの構成を示す説明図である。圧縮データの全体は、図18(A)に示すよう

に、ヘッダ部と圧縮データ部とダミー部とで構成されている。ヘッダ部は、それぞれ1バイトの4つのデータDFH, DFL, DLH, DLLを有している。最初の2つのデータDFH, DFLは、圧縮データ部に含まれるデータの種類を示している。圧縮データ部のデータには、基本量子化テーブルBQTのデータ、フルカラー自然画像圧縮データ、ランレンジス画像圧縮データなどの種類がある。ヘッダの後部16ビットのデータ(DLH+DLL)は、圧縮データ部とダミー部の合計のデータ長を示している。圧縮データ部はハフマン符号を含む可変長のデータなので、ダミー部との合計のデータ長が、ワード(=2バイト)の整数倍の長さになるように調整されている。

【0072】図18(B)は、基本量子化テーブルBQTを表わす圧縮データの構成を示している。この1セットの圧縮データは、Y信号用の基本量子化テーブルBQTを表わすデータと、U信号/V信号共用の基本量子化テーブルBQTを表わすデータとを含んでいる。なお、基本量子化テーブルBQTを表わすデータはハフマン符号化しておかなくてもよい。

【0073】図18(C)は、フルカラー自然画像の圧縮データの構成を示している。圧縮データ部には、量子化レベル係数QCxを表わす符号データ(図12におけるカテゴリSSSS=16~31の符号語)と、各画素ブロックの符号データであるブロックデータと、一様色の複数の画素ブロックを示すヌルランデータとを含んでいる。

【0074】図18(D)に示すように、1ユニットのブロックデータは4組のY信号用データと、1組のU信号データと、1組のV信号用データとで構成されている。図19は、YUVの各信号のブロックの関係を示す説明図である。図19(A)に示すように、この実施例における1画面は、256画素×240走査線の大きさを有している。Y信号に関しては、間引きをせずに、8×8画素の画素ブロック毎にDCT変換が行なわれる。一方、U信号とV信号に関しては、図19(B)に示すように、横方向と縦方向に1/2に間引き(サブサンプリング)されて、間引き後の8×8画素のブロックに対してDCT変換が行なわれる。従って、図19(C)に示すように、Y信号の4つの画素ブロックY1~Y4の領域がU信号とV信号の1つの画素ブロックの領域に対応している。なお、Y信号を間引きせずにU信号とV信号を間引きするのは、人間の目が輝度の変化(Y信号の変化)には比較的敏感であるが、色の変化(U信号とV信号の変化)には比較的鈍感だからである。U信号とV信号のみを間引くことによって、画質を過度に劣化させずに圧縮率を高めることができる。なお、図18(D)に示す1ユニットのブロックデータは、図19(C)に示す各領域のハフマン符号データを順に並べたものである。

21

【0075】 ブロックデータ内の1つの画素ブロックに対する符号データは、図18(F)に示すように、DC係数の1つのハフマン符号データと、AC係数の複数のハフマン符号データとで構成されている。DC係数のハフマン符号データは、前述したように、カテゴリSSSのハフマン符号語HFDと識別データIDとで構成される(図18(G))。また、AC係数のハフマン符号データは、ゼロラン長NNNNとカテゴリSSSSとの組み合わせに対するハフマン符号語HFACと、識別データIDとで構成される(図18(H))。

【0076】量子化レベル係数QCxの符号データは、圧縮データ部の先頭と、量子化レベル係数QCxの値を変更したい画素ブロックのブロックデータの直前に挿入されている。2番目の量子化レベル係数QCxが挿入される前の複数の画素ブロックに対しては、先頭の量子化レベル係数QCxが共通に使用される。また、3番目の量子化レベル係数QCx(図示せず)が挿入される前の複数のブロックに対しては、2番目の量子化レベル係数QCxが共通に使用される。

【0077】なお、圧縮データ部の先頭に量子化レベル係数QCxの符号語が含まれていない場合には、QCx=1であると見なされる。従って、圧縮データ部の先頭に量子化レベル係数QCxが挿入されておらず、途中に量子化レベル係数QCxが1回だけ挿入されている場合にも、量子化レベル係数QCxが2つ指定されていることと等価である。

【0078】量子化レベル係数QCxを表わすハフマン符号は、ブロックデータの間に挿入されているので、新たな量子化レベル係数QCxが復号化された時点の次のブロックデータに対してこの新たな量子化レベル係数QCxを容易に適用することができる。また、図12に示すように、量子化レベル係数QCxの符号データはDC係数用のハフマン符号語で表わされているので、これがブロックデータの間に挿入されても、この符号データがブロックY1用のDC係数の符号データであるか、量子化レベル係数QCxの符号データであるかを直ちに判断することができる。

【0079】圧縮データ部に含まれているヌルランデータは、図18(E)に示すように、ヌルランデータであることを示すDC係数用符号語「NRL」と、ブロック数と、識別データIDとで構成されている。

【0080】図20は、ヌルランデータによって表わされる画像を示す説明図である。図20(A)の原画像の背景BGは一様色で塗られている。図20(A)の楕円の部分は、図20(B)に示すようにすべての画素が同じ画像データ値( $f(x, y) = 12$ )を有する画素ブロックが18個連続しているものと仮定する。図20(C)は、これらの画素ブロックを表わすヌルランデータを示している。このヌルランデータは、16画素ブロック分の第1のヌルランデータNRD1と、2画素ブロ

10

ック分の第2のヌルランデータNRD2を含んでいる。【0081】各ヌルランデータNRD1、NRD2の先頭には、ヌルランデータであることを示すDC係数用符号語「NRL」(図12のカテゴリSSSS=15の符号語「1111011」)を有している。図18(F)に示すように、通常のブロックデータの先頭にはDC係数のハフマン符号が配置されているので、先頭にあるDC係数用符号語を復号化することによって、ヌルランデータと、ブロックデータと、量子化レベル係数QCxの符号データとを一意にかつ瞬時に識別することができる。

20

【0082】図20(C)に示すように、ブロック数は、AC係数用ハフマン符号語で表わされている。図21は、AC係数用ハフマン符号テーブル(図15)のうちでヌルランデータに使用される部分を示す図である。ヌルランデータに使用される場合には、ゼロラン長NNNNは([ブロック数]-1)に等しいと設定される。また、AC係数の値は1であるとして、カテゴリSSSS=1のハフマン符号語が使用される。図20(C)に示す第1のヌルランデータNRD1におけるブロック数のデータ(NNNN/SSSS=15/1)は一様色の画素ブロックが16個連続していることを示している。また、第2のヌルランデータNRD2におけるブロック数のデータ(NNNN/SSSS=1/1)は一様色の画素ブロックが2個連続していることを示している。

【0083】各ヌルランデータNRD1、NRD2の後端には、識別データIDが付加されている。この実施例では、ID=1に固定されている。

30

【0084】ヌルランデータは、このように、20ビット程度のデータによって連続した複数の画素ブロックが一様色であることを表わすことが可能である。一方、通常のブロックデータによって一様色の1セットのブロック(図19に示すY信号を4画素ブロック、U信号、V信号を各1画素ブロック含む)を表わすには、約300~約400ビット必要である。しかも、複数セットの画素ブロックが一様色であることを示す場合にも、各セットについて約300~約400ビット必要である。従って、ヌルランデータを使用すれば、連続する一様色の多数の画素ブロックを表わす圧縮データのデータ量をかなり低減することが可能である。

40

【0085】なお、ヌルランデータで表わされる一様色の画素ブロックの輝度信号Yや色差信号U、Vの値は、圧縮データには含まれておらず、ビデオゲームを記述するソフトウェアプログラムの中において指定されている。オペレータは、ビデオゲーム用のソフトウェアプログラマを作成する際に、一様色の画素ブロックの領域(図20(A)では背景BG)の範囲をマウス104で指定するとともに、これらのブロックの輝度や色調をキーボード103やマウス104を用いて指定する。こうすれば、例えばビデオゲーム装置20(図3)を用いてゲームを実行している途中に特定のイベントが発生した

50

22

場合に、背景BGの色を時間的に変化させるなどの特殊な視覚的効果を生じさせることができる。なお、スルランデータを復号化する具体的な回路構成については更に後述する。

【0086】D. 逆量子化テーブル作成部の詳細構成：図22は、図1に示す逆量子化テーブル作成部250の内部構成を示すブロック図である。逆量子化テーブル作成部250は、基本量子化テーブルBQTを記憶するRAM251と、RAM251のアドレスを生成するアドレス生成回路252と、量子化レベル係数QCxを保持するラッチ回路253と、量子化レベル係数QCxと基本量子化テーブルBQTとを乗算して量子化テーブルQTを生成する乗算ユニット254とを備えている。乗算ユニット254によって作成された量子化テーブルQTは、逆量子化部220に供給される。

【0087】まず、CD-ROMに収納されている圧縮画像データZZがハフマン復号化部210に与えられると、基本量子化テーブルBQTの符号データが最初に復号化されてRAM251に供給される。この基本量子化テーブルBQTは、アドレス生成回路252から与えられるライトアドレスに従ってRAM251に記憶される。RAM251に記憶された基本量子化テーブルBQTは、すべての画素ブロックに対して使用される。

【0088】アドレス生成回路252は、ハフマン復号化部210から出力されるDCT係数データQF(u,v)と同期してリードアドレスを発生し、このリードアドレスに応じてRAM251から基本量子化テーブルBQTが読み出される。一方、ハフマン復号化部210で復号化された量子化レベル係数QCxは、ラッチ回路253でラッチされて、次の量子化レベル係数QCxが与えられるまでラッチ回路253に保存される。従って、量子化レベル係数QCxが新たに供給されるまでは、複数の画素ブロックに対して同じ量子化レベル係数QCxが共通に使用される。

【0089】図23は、逆量子化テーブル作成部250(図22)に含まれているラッチ回路253と乗算ユニット254の内部構成を示すブロック図である。ラッチ回路253は、2つのラッチ402, 404で構成されている。乗算ユニット254は、同期クロック作成回路412と、AND回路414と、U信号スタート検出回路416と、V信号スタート検出回路418と、NAND回路420と、セレクタ422と、乗算器424と、クリッピング回路426と、ゼロ値修正回路428とを有している。

【0090】図24は、図23に示す回路の動作を示すタイミングチャートである。ハフマン復号化部210(図1)で量子化レベル係数QCxが復号化されると、量子化レベル係数QCxのデータがイネーブル信号QE Nとともにラッチ回路253に与えられる(図24(a), (b))。第1のラッチ402はイネーブル信

号QE Nの立ち上がりエッジで量子化レベル係数QCxをラッチして出力Q1を第2のラッチ404に供給する(図24(c))。

【0091】図24(d)に示すように、RAM251(図22)からはY信号用の基本量子化テーブルBQTが4回読出された後に、U信号/V信号共用の基本量子化テーブルBQTが2回読出される。

【0092】同期クロック作成回路412(図23)には、RAM251から読出される基本量子化テーブルBQTに同期して、Y信号の期間でLレベルとなり、U信号とV信号の期間ではHレベルとなるブロック識別信号UV/Yがアドレス生成回路252から与えられる(図24(e))。また、この際、イネーブル信号ENも同期クロック作成回路412に与えられる。同期クロック作成回路412は、ブロック識別信号UV/Yを反転して同期クロック信号SCK(図24(f))を生成し、これを第2のラッチ404のクロック入力端子に供給する。第2のラッチ404は、この同期クロック信号SCKの立ち上がりエッジで第1のラッチ402の出力Q1をラッチして、その出力Q2(図24(g))をセレクタ422のデータ入力端子に供給する。なお、セレクタ422の他のデータ入力端子には、固定値「1」が与えられている。

【0093】U信号スタート検出回路416は、ブロック識別信号UV/Yと10MHzの基本クロック信号CLKとに基づいて、U信号用の基本量子化テーブルBQTのスタート時刻を示すUスタート信号USTRT(図24(h))を生成する。このUスタート信号USTRTは、ブロック識別信号UV/Yの立ち上がりエッジから100ナノ秒だけLレベルになる信号である。

【0094】AND回路414には、6つのブロックY1~Y4, U, Vの各期間でレベルが交互に切り替わるブロック切替信号SWTCH(図24(i))と、ブロック識別信号UV/Yとが与えられている。V信号スタート検出回路418は、AND回路414の出力と10MHzの基本クロック信号CLKとに基づいて、V信号用の基本量子化テーブルBQTのスタート時刻を示すVスタート信号VSTART(図24(j))を生成する。このVスタート信号VSTARTは、ブロック識別信号UV/YがHレベルの期間において、ブロック切替信号SWTCHの立ち上がりエッジから100ナノ秒だけLレベルになる信号である。

【0095】図25は、U信号スタート検出回路416とV信号スタート検出回路418の内部構成を示すブロック図である。これらの回路416, 418は、どちらもDフリップフロップとNAND回路とで構成されている。U信号スタート検出回路416のDフリップフロップ432のD入力端子にはブロック識別信号UV/Yが供給されており、そのクロック入力端子には10MHzの基本クロック信号CLKが入力されている。NAND

回路434の入力端子には、ブロック識別信号UV/YとDフリップフロップ432の反転出力とが与えられている。なお、ブロック識別信号UV/Yは基本クロック信号CLKに同期している。

【0096】ブロック識別信号UV/YがLレベルの間は、Dフリップフロップ432の反転出力がHレベルなので、NAND回路434の出力USTRTはHレベルに保たれる(図24(h)参照)。ブロック識別信号UV/YがLレベルからHレベルに切り替わった直後は、NAND回路434の出力(Uスタート信号USTRT)がLレベルになり、100ナノ秒後に基本クロック信号CLKのエッジでDフリップフロップ432が入力をラッチすると、NAND回路434の出力USTRTは再びHレベルに戻る。

【0097】V信号スタート検出回路418の動作も、U信号スタート検出回路416の動作と同じである。ただし、Dフリップフロップ436のD入力端子には、ブロック識別信号UV/Yとブロック切替信号SWTCHの論理積が与えられているので、ブロック識別信号UV/YがHレベルの期間において、ブロック切替信号SWTCHの立ち上がりエッジから100ナノ秒だけVスタート信号VSTARTがLレベルになる。

【0098】NAND回路420(図23)には、Uスタート信号USTRTとVスタート信号VSTARTとが入力されており、その出力(選択信号SEL)はセレクタ422のセレクト入力端子に供給されている。選択信号SEL(図24(k))は、U信号用の期間の始めとV信号用の期間の始めに100ナノ秒だけHレベルになる。セレクタ422は、選択信号SELがLレベルの場合にはラッチ回路253から与えられた出力Q2をそのまま出力し、一方、選択信号SELがHレベルの場合には固定値「1」を出力する。セレクタ422の出力Q3は、乗算器424によって基本量子化テーブルBQTと乗算される。

【0099】図24(1)に示すように、U信号用の期間の始めとV信号用の期間の始めには、セレクタ422の出力Q3は量子化レベル係数QCxの値に係わらずに必ず「1」になる。U信号用とV信号用の期間の始めの100ナノ秒の期間は、DC係数用の量子化レベルを算出するための期間である。従って、図23の回路を使用すれば、U信号とV信号に使用するDC係数用の量子化レベルと、指定された量子化レベル係数QCxとの乗算を実質的に行なわないようにすることができる。換言すれば、図23に示す回路は、図8(c), (d)に示す演算を実現する回路である。

【0100】図23に示すように、乗算器424の出力はクリッピング回路426とゼロ値修正回路428によって修正されて最終的な量子化テーブルQQTとなる。図26は、これらの2つの回路426, 428の内部構成を示すブロック図である。

【0101】クリッピング回路426は、4入力OR回路450と、8個の2入力OR回路452とで構成されている。これらの回路は、量子化レベルを9ビット(最上位ビットは符号ビット)で表現する場合の回路である。4入力OR回路450には、乗算器424(図23)から出力された14ビットのデータのうちで符号ビットD13を除く上位4ビットD9~D12が入力されている。8個の2入力OR回路452の一方の入力端子には、4入力OR回路450の出力が与えられており、他方の入力端子には乗算器424の出力の下位8ビットD1~D8が与えられている。上位4ビットD9~D12の少なくとも1つの値が「1」の場合には、8つの2入力OR回路452の出力がすべて「1」になる。従って、乗算器424の出力が十進数で255以上の場合には、クリッピング回路426の出力が255に設定される。

【0102】ゼロ値修正回路428では、クリッピング回路426内の7つのビットD2~D8用の7つの2入力OR回路452の出力がそのまま出力される。また、これらの7つの2入力OR回路452の出力と符号ビットD13とが、8つのインバータ460にそれぞれ与えられている。8つのインバータ460の出力は8入力AND回路462に与えられており、このAND回路462の出力は2入力OR回路464に供給されている。この2入力OR回路464には、最下位ビットD1用の2入力OR回路452の出力が与えられている。この結果、乗算器424の出力の13ビットD1~D13の値がすべて「0」の場合には、ゼロ値修正回路428は、最下位ビットの値のみが「1」で他の8ビットの値が「0」の量子化レベルQQTを出力する。換言すれば、ゼロ値修正回路420は、図9(c), (d)に示す演算を実現している。

【0103】E. ヌルランデータの復号化: 図27は、画像データ伸長装置200におけるハフマン復号化部210(図1)の内部構成を示すブロック図である。ハフマン復号化部210は、圧縮画像データEZをハフマン復号化する復号化部470と、制御部472と、セレクタ474と、DC係数レジスタ476とを備えている。

【0104】復号化部470は、与えられた圧縮データの種類が基本量子化テーブルBQT、量子化レベル係数QCx、ブロックデータ、ヌルランデータのいずれであるかを判断し、圧縮データの種類を示す状態信号SSを制御部472に供給する。制御部472は、この状態信号SSに応じて制御信号CTL1, CTL2, CTL3を復号化部470と、セレクタ474と、ビデオゲーム装置のMPU40(図3)とにそれぞれ供給する。復号化された基本量子化テーブルBQTと量子化レベル係数QCxは、復号化部470から逆量子化テーブル作成部250に供給される。復号後の量子化されたDC係数QF(u, v)は、復号化部470からセレクタ474

に供給される。

【0105】セレクタ474のデータ入力端子には、復号化部470から与えられたDCT係数QF(u, v)の他に、ゼロデータと、DC係数レジスタ476に登録されたDC係数QF(0, 0)とが与えられている。DC係数レジスタ476には、ゲームのソフトウェアプログラムにおいて記述されている一様色の画素ブロックのDC係数QF(0, 0)の値が、ビデオゲーム装置のMPU40によって書き込まれる。なお、DC係数QF(0, 0)は、YUV信号に対してそれぞれ異なる値が登録される。

【0106】ここで、図20(C)に示す2つのヌルランデータNRD1, NRD2を復号化する場合を考える。第1のヌルランデータNRD1の先頭のデータNRLが復号化部470によって検出されると、ヌルランデータであることを知らせる状態信号SSが復号化部470から制御部472に出力される。制御部472は、状態信号SSに応じて直ちに復号化部470とMPU40に制御信号CTL1, CTL2をそれぞれ出力して復号化動作を止めないように制御する。また、制御部472は、セレクタ474に制御信号CTL3を供給し、最初のブロックのDC係数として、DC係数レジスタ476に登録されているDC係数QF(0, 0)を選択させる。図20(B)に示すように、原画像データf(x, y)の値が12の場合には、QF(0, 0) = 12がDC係数レジスタ476に登録されている。制御部472は、さらに、63個のAC係数として、すべてゼロデータを選択するようにセレクタ474を制御する。図28は、こうして作成されたDCT係数QF(u, v)を示している。

【0107】第1のヌルランデータNRD1は、一様色のブロックが16個連続していることを示しているので、16個の画素ブロックのそれぞれについて、図28に示すDCT係数QF(u, v)が作成される。第2のヌルランデータNRD2についても同様に、2つのブロックのそれぞれについて、図28に示すDCT係数QF(u, v)が作成される。

【0108】ヌルランデータの処理が終了すると、制御部472が復号化部470とMPU40に対して制御信号CTL1, CTL2を出力し、復号化動作を再開するよう制御する。ヌルランデータによって一様色であることが指定されている画素ブロックの輝度信号Yや色差信号U, Vの値は、MPU40からDC係数レジスタ476に書き込むDC係数の値を変更することによって容易に変更することができる。言い換えれば、ヌルランデータを利用すれば、一様色の画像領域の色を圧縮画像データ以外のデータに応じて所望の色に変更することができる。この実施例では、一様色のブロックの輝度信号Yや色差信号U, Vの値は、ビデオゲームを記述するソフトウェアプログラムの中で指定されている。

【0109】図29は、ハフマン復号化部の他の構成を示すブロック図である。このハフマン復号化部210aでは、復号化部470から出力されたDCT係数QF(u, v)がセレクタ474をバイパスして逆量子化部220に直接与えられている。セレクタ474は、DC係数レジスタ476から与えられるDC係数QF(0, 0)とゼロデータの一方を選択し、逆量子化部220をバイパスしてIDCT部230に直接供給している。制御部478は、図27と同様の3つの制御信号CTL1～CTL3の他に、第4の制御信号CTL4をIDCT部230に出力している。

【0110】制御部478は、圧縮データの種類を示す状態信号SSに応じて復号化部470と、MPU40と、IDCT部230とに、制御信号CTRL1, CTL2, CTL4をそれぞれ出力する。復号化部470で通常のブロックデータが復号されると、復号化されたDCT係数QF(u, v)が逆量子化部220に供給される。

【0111】図20(C)に示す第1のヌルランデータNRD1の先頭のデータNRLが復号化部470によって検出されると、ヌルランデータであることを知らせる状態信号SSが復号化部470から制御部478に出力される。制御部478は、状態信号SSに応じて直ちに復号化部470とMPU40に制御信号CTL1, CTL2をそれぞれ出力して復号化動作を止めないように制御する。また制御部478は、セレクタ474に与える制御信号CTL3のレベルを切換えて、DC係数レジスタ476に登録されているDC係数QF(0, 0)を選択させる。制御部478は、さらに、63個のAC係数として、すべてゼロデータを選択するようにセレクタ474を制御する。制御部478は、同時にIDCT部230に対して制御信号CTL4を出力し、セレクタ474の出力を選択して逆変換するよう制御する。

【0112】ヌルランデータの処理が終了すると、制御部478から復号化部470とMPU40に制御信号CTL1, CTL2が出力され、復号動作を再開するよう制御する。

【0113】このように、図29に示す回路では、ヌルランデータの処理の際には、セレクタ474から出力されるDCT係数QF(u, v)が逆量子化部220をバイパスしてIDCT部230に直接供給されるので、逆量子化による演算誤差が生じないという利点がある。例えば、特定の色の部分を検出し透明色にするクロマキー処理をビデオエンコーダユニット50(図3)によって行う時に、量子化による演算誤差が最小に押さえられるので、所望の色の画素ブロックを確実に透明色にすることが可能である。

【0114】上記の実施例においては、ヌルランデータによってDC係数のみを任意に設定できるようにしたが、AC係数の所定の部分(例えばQF(1, 0))とQ

$F(0, 1)$ ）に対して所望の値を設定できるようにすることも可能である。この場合には、ヌルランデータは同一の画像パターンを有する一連の画素ブロックの個数を表わす。

【0115】F. 変形例：なお、この発明は上記実施例に限られるものではなく、その要旨を逸脱しない範囲において種々の態様において実施することが可能であり、例えば次のような変形も可能である。

【0116】（1）上記実施例では直交変換として2次元DCT変換を取り上げたが、本発明には任意の直交変換（例えばK-L変換やアダマール変換）が利用可能である。また、エントロピー符号化としては、ハフマン符号化以外の任意の符号化（例えば算術符号化やMEL符号化）を利用することができる。

【0117】（2）画像データ圧縮装置100はハードウェアで実現してもよく、また、画像データ伸長装置200をソフトウェアで実現してもよい。図30は、ソフトウェアによって図18に示す圧縮データを伸長する処理の手順を示すフローチャートである。ステップS1では、ヘッダ部の値から圧縮データの内容を判断する。圧縮データが基本量子化テーブルBQTを表わしている場合には、ステップS2において基本量子化テーブルBQTをメモリに記憶し、ステップS1に戻る。

【0118】圧縮データが画像を表わしている場合には、ステップS3において圧縮データ部内に含まれている各データユニットの先頭データの種類が判断される。ここで、データユニットとは、図18（C）に示す量子化レベル係数QCxの符号データ、ブロックデータ、ヌルランデータのそれぞれを意味している。図18-

（C）、（D）、（E）、（F）に示すように、データユニットの先頭データとしては、量子化レベル係数QCxを表わすハフマン符号語と、ブロックデータのDC係数△DCのハフマン符号語と、ヌルランを表わすハフマン符号語NRLの3種類がある。

【0119】先頭データが量子化レベル係数QCxのハフマン符号語である場合には、ステップS4において、この量子化レベル係数QCxと基本量子化テーブルBQTを乗することによって量子化テーブルQTTを作成する。なお、この乗算の際に、前述した特徴、すなわち、（1）乗算結果が最大値以上の量子化レベルを最大値に設定すること、（2）DC係数は量子化しないこと、（3）量子化レベル係数QCxの値が0の時にはすべての量子化レベルを1に設定すること、が実現される。ステップS4において量子化テーブルQTTが作成されると、ステップS3に戻る。

【0120】ステップS3において先頭データがDC係数△DCのハフマン符号語であると判断された場合には、ステップS5において1画素ブロック分のデータが復号化されてDCT係数QF(u, v)が得られる。ステップS6では逆量子化が行なわれ、ステップS7で

は、2次元DCT逆変換が行なわれる。1セットのブロックデータ（図18（D））に含まれるすべての画素ブロックについてステップS5～S7の処理が繰り返されると、ステップS8からステップS3に戻る。

【0121】ステップS3において、先頭データがヌルランデータであることを示すハフマン符号語NRLであると判断された場合には、ステップS9において、1画素ブロックのDC成分が予め指定された値に設定されるとともに、63個のAC成分がすべて0に設定される。

10ステップS10では、こうして作成されたDCT係数に対して2次元DCT逆変換が行なわれる。ヌルランデータで指定されたブロック数だけステップS9, S10の処理が繰り返されると、ステップS11からステップS3に戻る。こうして、ステップS3～S11の処理を圧縮データ部のすべてに亘って実行することによって画像データff(x, y)が復元される。

【0122】（3）上記実施例では、本発明をビデオゲーム装置に適用した例を示したが、本発明はあらゆる種類の画像処理装置に適用することが可能である。

20【0123】（4）上記実施例では、量子化レベル係数QCxや量子化テーブルQTT内の各量子化レベルを整数としたが、これらは小数を含む数でもよい。

#### 【0124】

【発明の効果】以上説明したように、請求項1, 12および23に記載した発明によれば、第1と第2の係数を与えるだけで量子化レベルの異なる第1と第2の量子化テーブルを作成できるので、2つの量子化テーブルを圧縮画像データに含める場合に比べて圧縮画像データのデータ量を低減することができるという効果がある。

30【0125】請求項2, 13および24およびに記載した発明によれば、第1と第2の係数を第1と第2の画素ブロックにそれぞれ容易に対応づけられるという効果がある。

【0126】請求項3, 14および25に記載した発明によれば、第1と第2の係数を一意にかつ瞬時に復号することができるという効果がある。

【0127】請求項4および15に記載した発明によれば、第2の係数が現われるまでの複数の画素ブロックに対して第1の係数が適用されるので、これらの複数の画素ブロックのそれぞれの符号データの前に第1の係数を含めておく必要がなく、従って、圧縮画像データのデータ量を低減することができるという効果がある。

40【0128】請求項5および16に記載した発明によれば、量子化レベルのビット数を一定値以下に保つことができるという効果がある。

【0129】請求項6および17に記載した発明によれば、所定の位置の変換係数を逆量子化するさいに発生する誤差を低減することができるという効果がある。

50【0130】請求項7および18に記載した発明によれば、変換係数の直流成分に対する量子化レベルを変えな

いようにすれば、直流成分の逆量子化にともなう誤差を低減することができるという効果がある。

【0131】請求項8および19に記載した発明によれば、互いに同一の画像パターンを有する連続した複数の画素ブロックを比較的少ないデータで表わすことができるという効果がある。

【0132】請求項9および20に記載した発明によれば、連続した複数の画素ブロックを同一の色で塗ることができるという効果がある。

【0133】請求項10および21に記載した発明によれば、逆量子化による誤差の発生を防止することができるという効果がある。

【0134】請求項11に記載した発明によれば、互いに同一のパターンの色を有する連続した複数の画素ブロックを比較的少ないデータで表わすことができるという効果がある。

【0135】請求項26に記載した発明によれば、互いに同一の画像パターンを有する連続した複数の画素ブロックを比較的少ないデータで表わすことができるという効果がある。

【0136】請求項27に記載した発明によれば、第1の画素ブロックを含む少なくとも1つの画素ブロックに対して新たな量子化テーブルを使用する場合に第1の係数コードを圧縮画像データに含めておけばよいで、圧縮画像データのデータ量を低減することができるという効果がある。

【0137】請求項28に記載した発明によれば、第1と第2の係数を与えるだけで量子化レベルの異なる第1と第2の量子化テーブルを作成できるので、2つの量子化テーブルを圧縮画像データに含める場合に比べて圧縮画像データのデータ量を低減することができるという効果がある。

【0138】また、請求項29に記載した発明によれば、第1と第2の係数を第1と第2の画素ブロックにそれぞれ容易に対応づけられるという効果がある。

#### 【図面の簡単な説明】

【図1】この発明の一実施例を適用した画像データの圧縮装置と伸長装置の機能を示すブロック図。

【図2】画像データ圧縮装置100の具体的な構成を示すブロック図。

【図3】画像データ伸長装置200の具体的な構成を示すブロック図。

#### 【図4】原画像を示す平面図。

【図5】DCT係数F(u, v)の配列を示す説明図。

【図6】圧縮／伸長の基本動作を示す説明図。

【図7】量子化レベル係数QCxが4の場合の圧縮／伸長動作を示す説明図。

【図8】量子化においてDC成分を変更しない場合の動作を示す説明図。

#### 【図9】量子化レベル係数QCxが0の場合の圧縮／伸

長動作を示す説明図。

【図10】DC係数符号化部の機能を示すブロック図。

【図11】ハフマン符号化におけるカテゴリ化テーブルを示す図。

【図12】DC係数用のハフマン符号テーブルHTDCの一例を示す説明図。

【図13】AC係数符号化部の機能を示すブロック図。

【図14】AC係数のジグザグスキャンの順路を示す説明図。

10 【図15】AC係数用の2次元ハフマン符号テーブルを示す説明図。

【図16】ハフマン符号テーブルの内容を示す図。

【図17】ハフマン符号化の一例を示す説明図。

【図18】圧縮データの構成を示す説明図。

【図19】YUVの各信号のブロックの関係を示す説明図。

【図20】ヌルランデータによって表わされる画像を示す説明図。

20 【図21】AC係数用のハフマン符号テーブルの他の部分を示す図。

【図22】逆量子化テーブル作成部250の内部構成を示すブロック図。

【図23】ラッチ回路253と乗算ユニット254の内部構成を示すブロック図。

【図24】図23に示す回路の動作を示すタイミングチャート。

【図25】U信号スタート検出回路416とV信号スタート検出回路418の内部構成を示すブロック図。

【図26】クリッピング回路426とゼロ値修正回路428の内部構成を示すブロック図。

30 【図27】ハフマン復号化部210の内部構成を示すブロック図。

【図28】ヌルランデータを復号化して得られたDCT係数を示す図。

【図29】ハフマン復号化部の他の構成を示すブロック図。

【図30】ソフトウェアによって圧縮データを復号する処理の手順を示すフローチャート。

40 【図31】従来の画像データ圧縮装置と伸長装置を示すブロック図。

#### 【符号の説明】

20…ビデオゲーム装置

21…ROM

24, 26…ゲームパッド

28…カラーテレビ

30…ビデオ信号ケーブル

32…CD-ROMドライブ

34…スピーカ

36…SCSIバス

38…スピーカ

4 0 … MPU  
 4 0 a … 演算部  
 4 0 b … コントローラ  
 4 1 … メインメモリ  
 4 2 … ROM  
 4 3 … M-BUS  
 4 5 … 画像信号コントロールユニット  
 4 5 a … MPU I/F  
 4 5 b … SCS I コントローラ  
 4 5 c … AFFIN コンバータ  
 4 5 d … グラフィックコントローラ  
 4 5 e … サウンドコントローラ  
 4 9 … VDP ユニット  
 5 0 … ビデオエンコーダユニット  
 5 0 a … ルックアップテーブル  
 5 0 b … 優先順位設定部  
 5 0 c … ミキサー  
 5 0 c … 画像データ合成部  
 5 0 d … スーパーインポーズ部  
 5 0 e … DAC  
 5 2 … 音声出力ユニット  
 5 2 a … ADPCM 部  
 5 5 … RAM  
 5 9 … RAM  
 6 0 … NTSC コンバータ  
 1 0 0 … 画像データ圧縮装置  
 1 0 1 … CPU  
 1 0 2 … メインメモリ  
 1 0 3 … キーボード  
 1 0 4 … マウス  
 1 0 5 … 磁気ディスク装置  
 1 0 6 … 光磁気ディスク装置  
 1 1 0 … DCT 部  
 1 2 0 … 量子化部  
 1 3 0 … ハフマン符号化部  
 1 3 1 … ブロック遅延部  
 1 3 2 … 加算器  
 1 3 3 … カテゴリ化処理部  
 1 3 4 … 1 次元ハフマン符号化部  
 1 3 5 … ジグザグスキヤン部  
 1 3 6 … 判定部  
 1 3 7 … ランレングスカウンタ  
 1 3 8 … カテゴリ化部  
 1 4 0 … 量子化テーブル作成部  
 1 5 0 … ハフマン符号テーブルメモリ  
 2 0 0 … 画像データ伸長装置  
 2 1 0 … ハフマン復号化部  
 2 2 0 … 逆量子化部  
 2 3 0 … IDCT 部  
 2 4 0 … ハフマン符号テーブルメモリ

特開平7-123413

34  
 2 5 0 … 逆量子化テーブル作成部  
 2 5 1 … RAM  
 2 5 2 … アドレス生成回路  
 2 5 3 … ラッチ回路  
 2 5 4 … 乗算ユニット  
 4 0 2, 4 0 4 … ラッチ  
 4 1 2 … 同期クロック作成回路  
 4 1 4 … AND 回路  
 4 1 6 … U 信号スタート検出回路  
 10 4 1 8 … V 信号スタート検出回路  
 4 2 0 … NAND 回路  
 4 2 2 … セレクタ  
 4 2 4 … 乗算器  
 4 2 6 … クリッピング回路  
 4 2 8 … ゼロ値修正回路  
 4 3 2, 4 3 6 … D フリップフロップ  
 4 3 4, 4 3 8 … NAND 回路  
 4 5 0 … 4 入力 OR 回路  
 4 5 2 … 2 入力 OR 回路  
 20 4 6 0 … インバータ  
 4 6 2 … AND 回路  
 4 6 4 … OR 回路  
 4 7 0 … 復号化部  
 4 7 2 … 制御部  
 4 7 4 … セレクタ  
 4 7 6 … DC 係数レジスタ  
 4 7 8 … 制御部  
 5 4 0 … 画像データ圧縮装置  
 5 4 2 … 直交変換部  
 30 5 4 4 … 量子化部  
 5 4 6 … エントロピー符号化部  
 5 5 0 … 画像データ伸長装置  
 5 5 0 … 伸長装置  
 5 5 2 … 逆直交変換部  
 5 5 4 … 逆量子化部  
 5 5 6 … エントロピー復号化部  
 5 6 2 … 量子化テーブル  
 5 6 4 … 符号テーブル  
 BG … 背景  
 40 BQT … 基本量子化テーブル  
 CLK … 基本クロック信号  
 D 1 ~ D 1 3 … 乗算器 4 2 4 の出力  
 D 1 3 … 符号ビット  
 EN … イネーブル信号  
 f (x, y) … 原画像データ  
 f f (x, y) … 復号された画像データ  
 F (u, v) … 元の DCT 係数  
 FF (u, v) … 復号された DCT 係数  
 HFAC … AC 係数用ハフマン符号語  
 HFDCT … DC 係数用ハフマン符号語

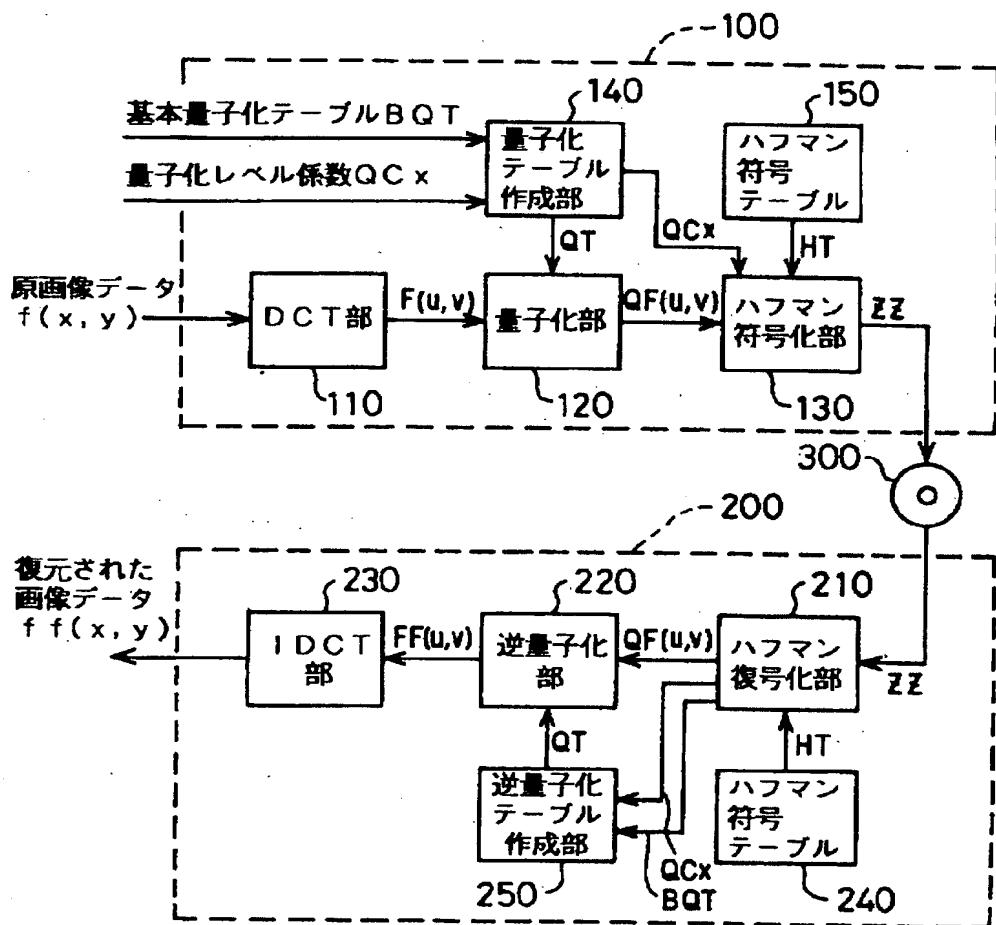
35

HTAC…A C係数用ハフマン符号テーブル  
 HTDC…D C係数用ハフマン符号テーブル  
 HT…ハフマン符号テーブル  
 ID…識別データ  
 NNNN…ゼロラン長  
 NR\_L…ヌルランデータ  
 PB…画素ブロック  
 PX…画素  
 QCx…量子化レベル係数  
 QEN…イネーブル信号  
 QF…量子化されたDCT係数

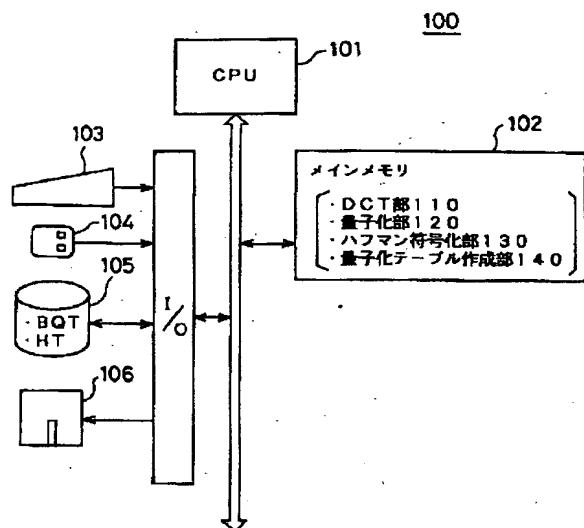
36

\* QT…量子化テーブル  
 SCK…同期クロック信号  
 SEL…選択信号  
 SS…状態信号  
 SSSS…カテゴリ  
 SWITCH…ブロック切替信号  
 U, V…色差信号  
 USTRT…Uスタート信号  
 VSTRT…Vスタート信号  
 Y…輝度信号  
 \* ZZ…圧縮画像データ

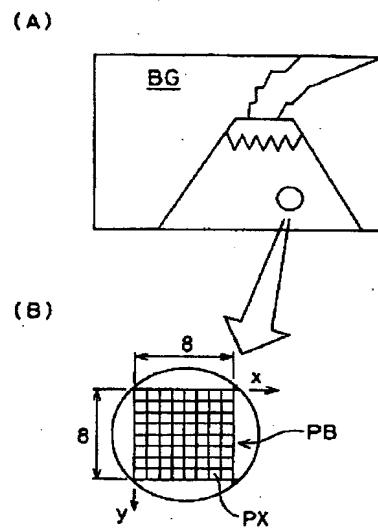
【図1】



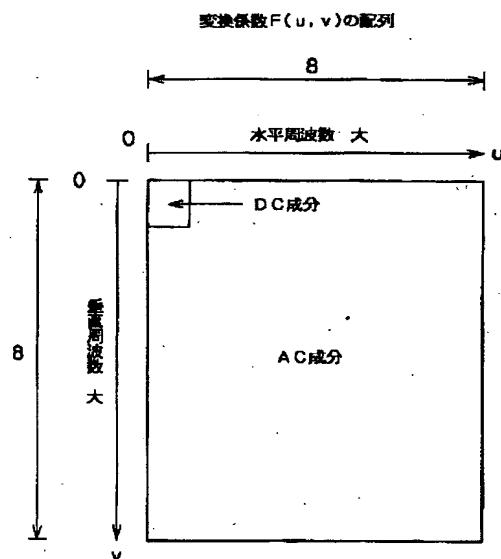
【図2】



【図4】



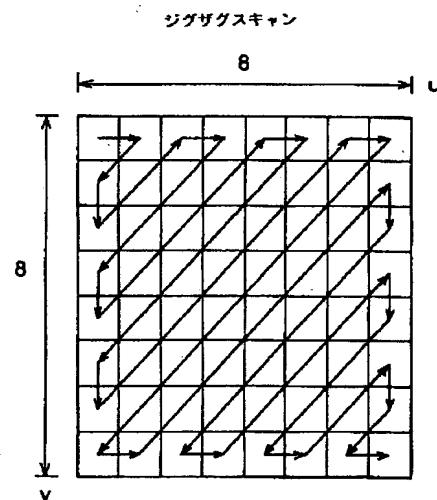
【図5】



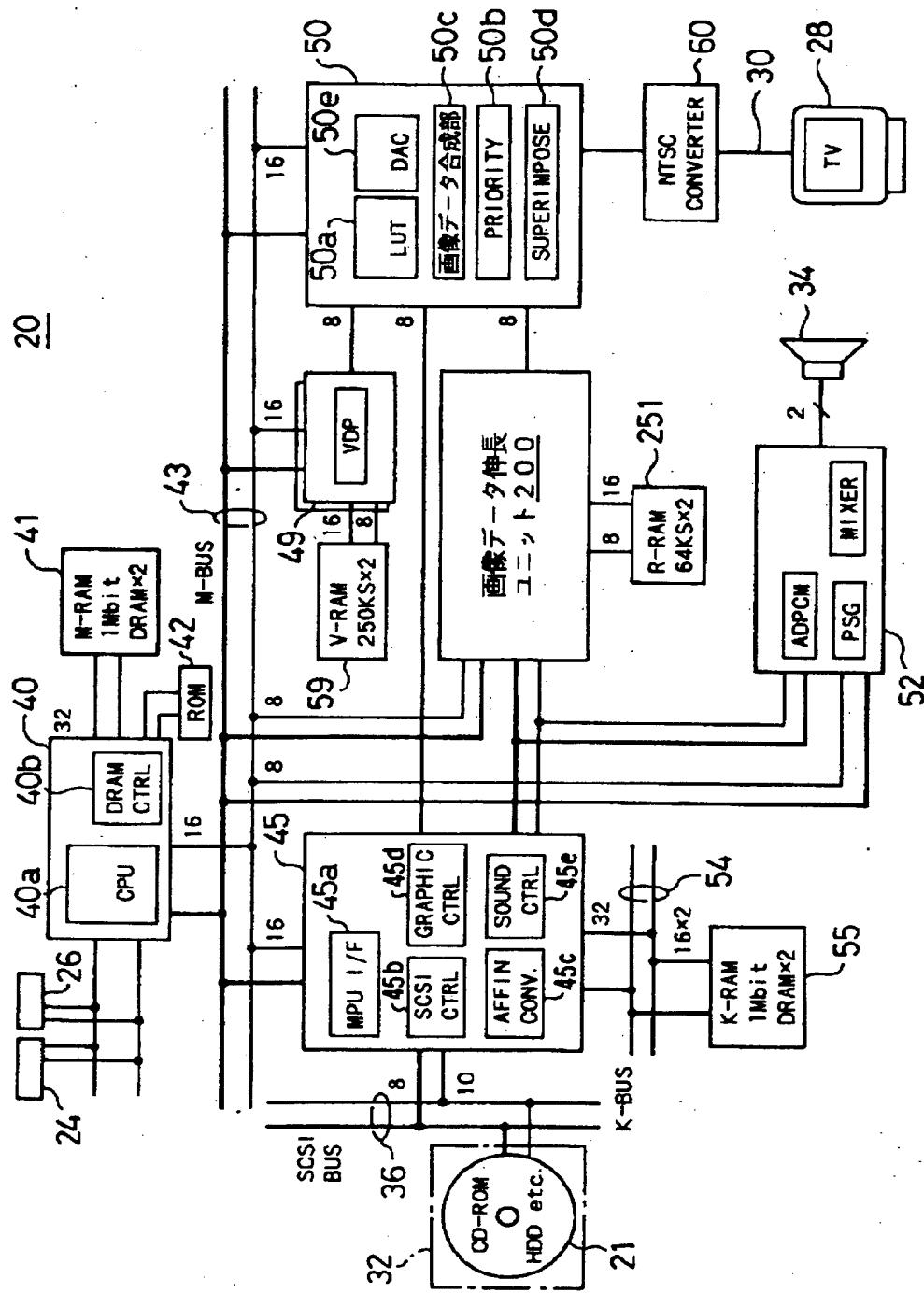
【図11】

カテゴリ SSSS	ADC/AC係数値	識別データ ビット数
0	0	0
1	-1, 1	1
2	-3, -2, 2, 3	2
3	-7, -4, 4, 7	3
4	-15, -8, 8, 15	4
5	-31, -16, 16, 31	5
6	-63, -32, 32, 63	6
7	-127, -64, 64, 127	7
8	-255, -128, 128, 255	8
9	-511, -256, 256, 511	9

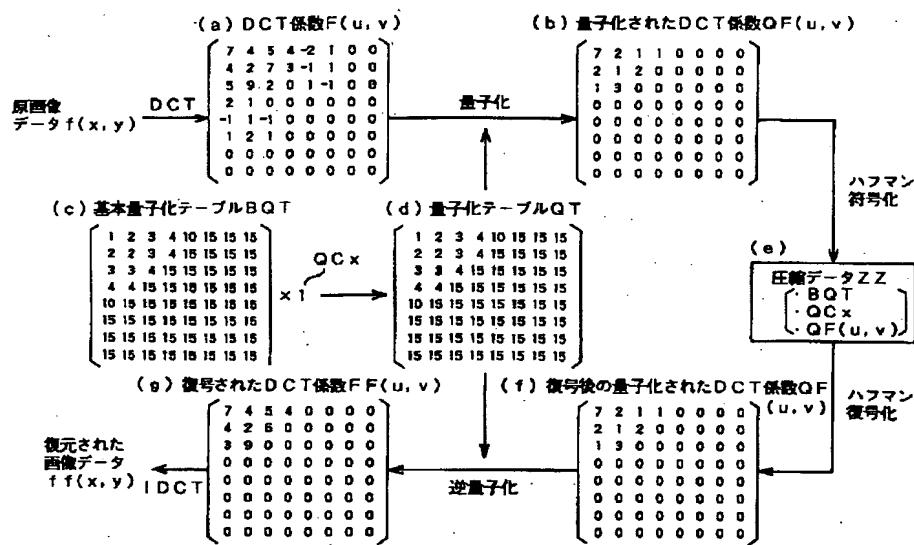
【図14】



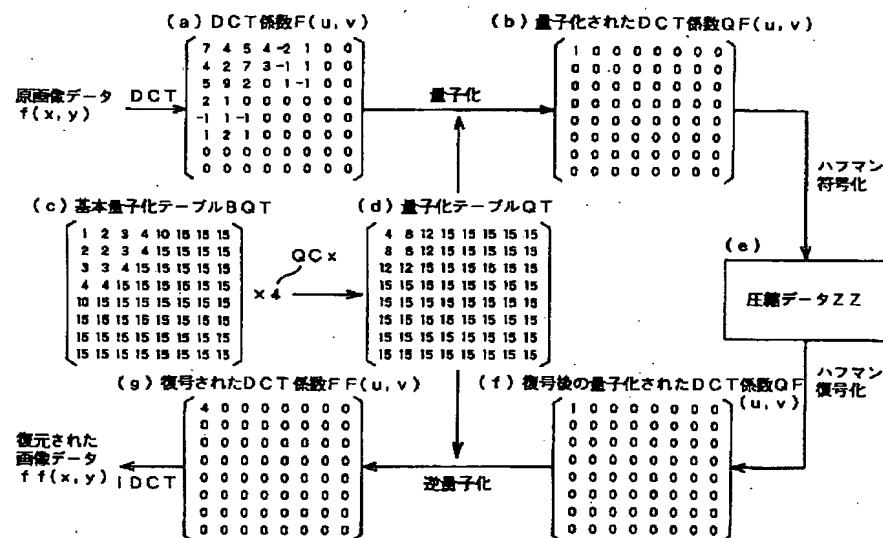
[図3]



【図6】



【図7】

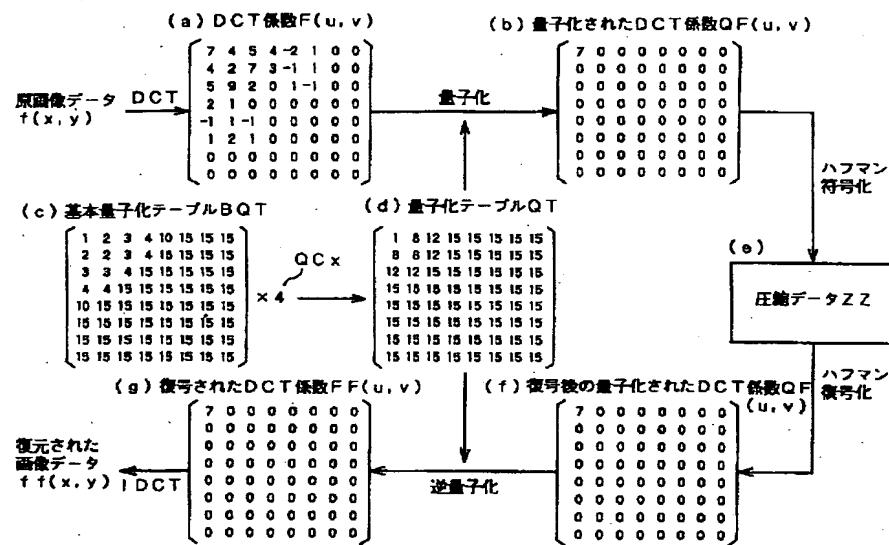


【図15】

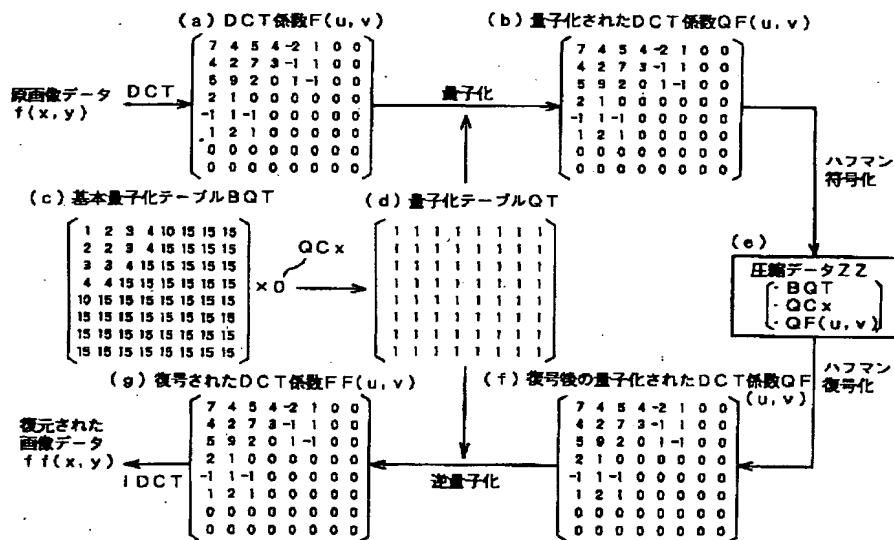
		カテゴリ SSSS			
		0	1	-----	9
ゼ リ 長 NNNN	0	EOB			
	1	ZRL			
2	-				
15	-				

AC係数  
ハフマン符号語  
HFAC

[図8]

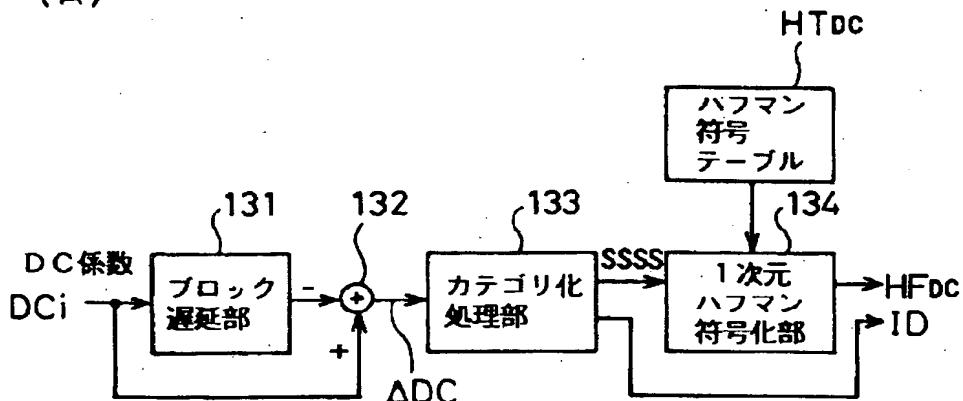


[図9]

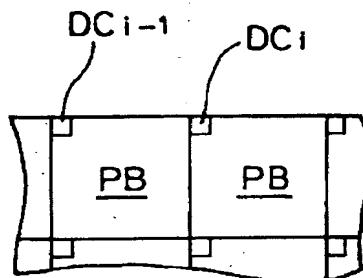


【図10】

(A)

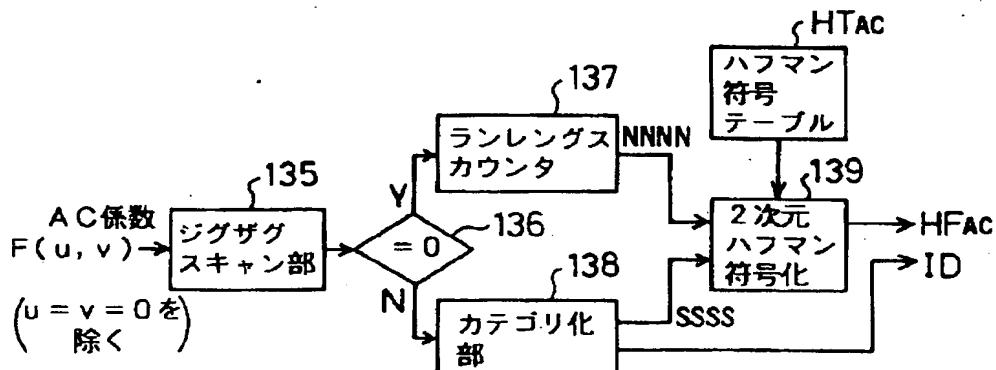


(B)



$$\Delta DC = DC_i - DC_{i-1}$$

【図13】



[図12]

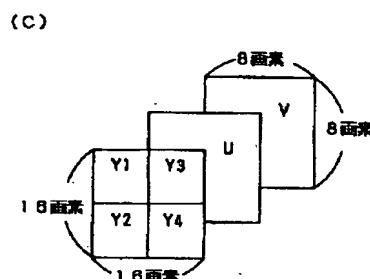
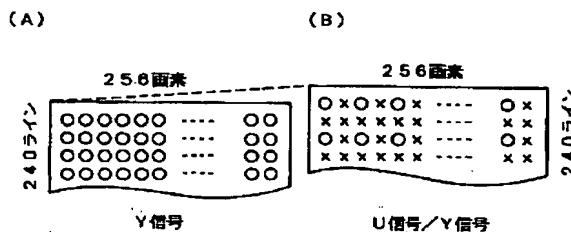
カテゴリ SSSS	△DCT係数 ハフマン符号語 HF <sub>DCT</sub>	
	Y信号用	U信号、V信号用
0	000	00
1	1110	01
2	001	10
3	010	110
4	011	1110
5	100	11110
6	101	111110
7	110	1111110
8	111100	11111110
9	1111010	11111111
10		
11		
12		
13		
14		
15(QR1)	1111011	
16(QC0)	11110000	
17(QC1)	11110001	
18(QC2)	11110010	
19(QC3)	11110011	
20(QC4)	11110100	
21(QC5)	11110101	
22(QC6)	11110110	
23(QC7)	11110111	
24(QC8)	11111000	
25(QC9)	11111001	
26(QC10)	11111010	
27(QC11)	11111011	
28(QC12)	11111100	
29(QC13)	11111101	
30(QC14)	11111110	
31(QC15)	11111111	

[図17]

[図16]

ゼロラン長 NNNN	AC係数ハフマン符号語 HF <sub>AC</sub>	
	Y信号用	U信号、V信号用
0/0 (EOB)	11111	11111
0/1	00	00
0/2	01	01
0/3	100	1010
0/4	1010	11000
0/5	11000	1110010
0/6	110110	111100010000
0/7	111011100	111100010001
0/8	111100010000	111100010010
0/9	11101111000	11101111000
1/0 (ZRL)	111101111111	111101111111
1/1	1011	100
1/2	11001	11001
1/3	1110100	1110011
1/4	111100010001	111100010011
1/5	111100010010	111100010100
1/6	111100010011	111100010101
1/7	111100010100	111100010110
1/8	111100010101	111100010111
1/9	11101111001	11101111001

[図19]

(B) DC係数符号化(DC<sub>i</sub>-1=0)

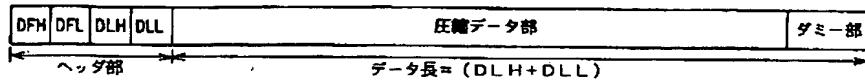
$$\Delta \text{DC} = 12 \rightarrow \text{SSSS} = 4 \rightarrow \text{HF}_{\text{DCT}} + \text{ID} = \underline{\text{0111100}} \quad \underline{\text{HF}_{\text{DCT}}} \quad \underline{\text{ID}}$$

(C) AC係数符号化

- ・ジグザグスキャン : 0, -4, 0, 3, 4, 0, 0, 5
- ・NNNN/SSSS = 1/3 1/2 0/3 2/3
- ・HF<sub>AC</sub> + ID = 11101000111100111

【図18】

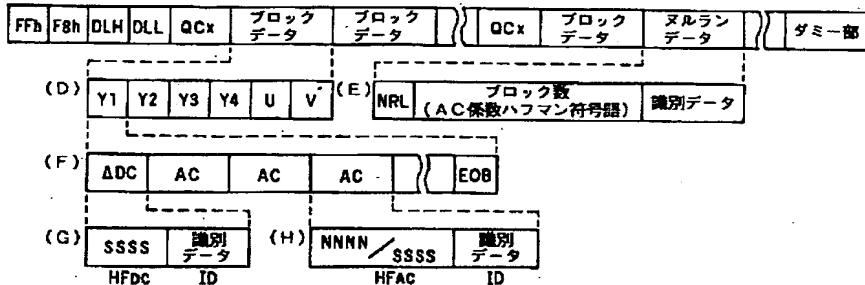
(A) 圧縮データ基本構造



(B) 基本量子化テーブルデータ

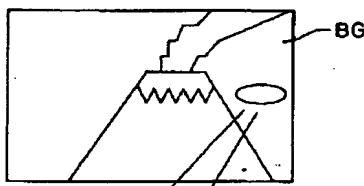
FFh	FFh	00h	82h	基本量子化テーブルデータ ( Y信号用 + U/V信号用 )		ダミー部
-----	-----	-----	-----	--------------------------------	--	------

(C) フルカラー自然画像圧縮データ

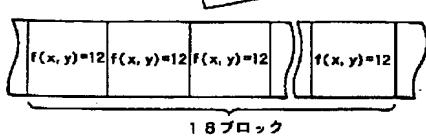


【図20】

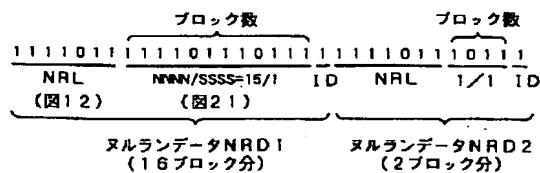
(A)



(B) 原画像データ



(C) ヌルランデータ



【図21】

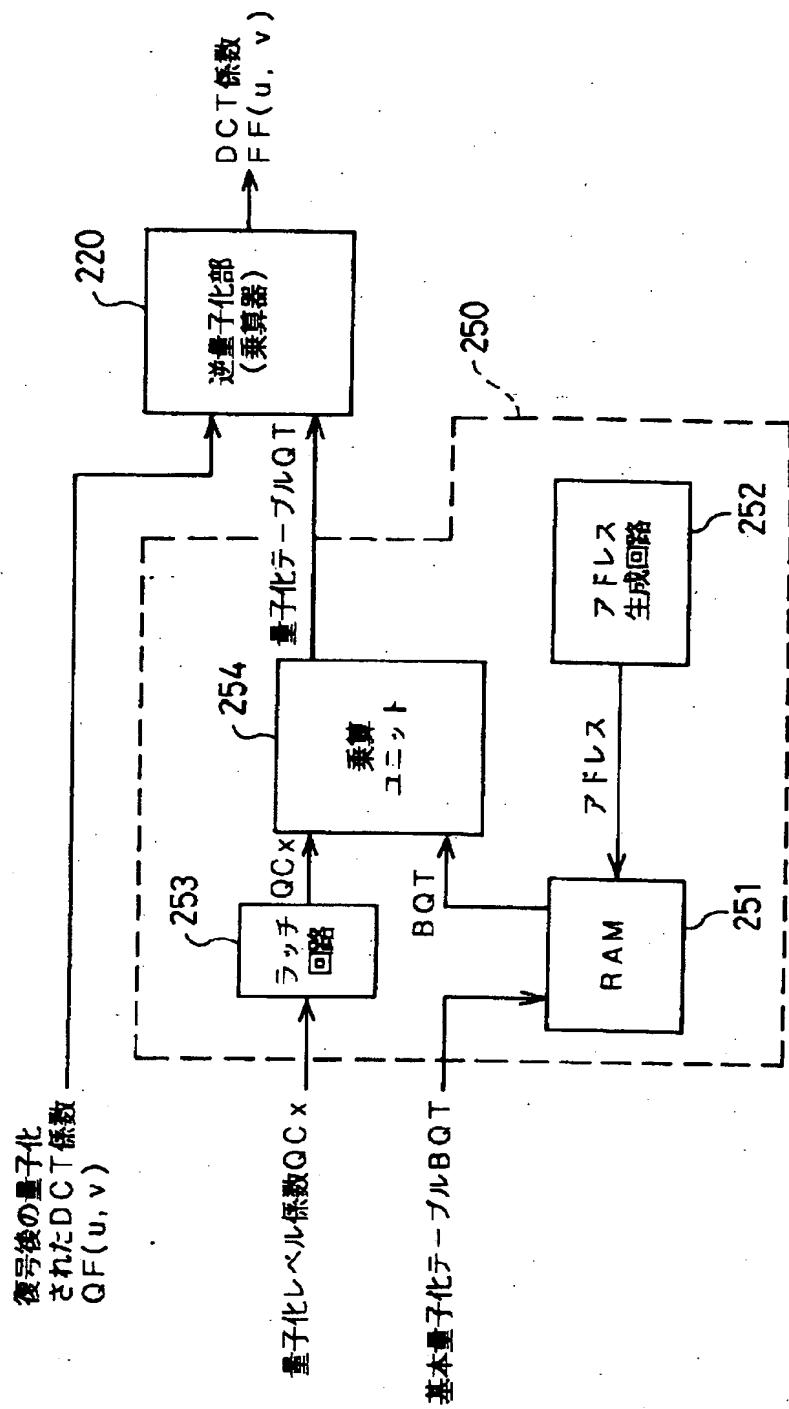
ゼロラン長 NNNN カテゴリ SSSS	AC係数ハフマン符号語HFACT	
	Y信号用	U信号、V信号用
0/1	00	00
1/1	1011	100
2/1	11010	1011
3/1	110111	11010
4/1	111000	11011
5/1	111001	111000
6/1	1110110	1110101
7/1	111011101	1110110
8/1	1110011111	111011101
9/1	11101000111	11101000111
10/1	11101001111	11101001111
11/1	11101010111	11101010111
12/1	11101011111	11101011111
13/1	11101100111	11101100111
14/1	11101101111	11101101111
15/1	11101110111	11101110111

ヌルランデータに使用する場合：

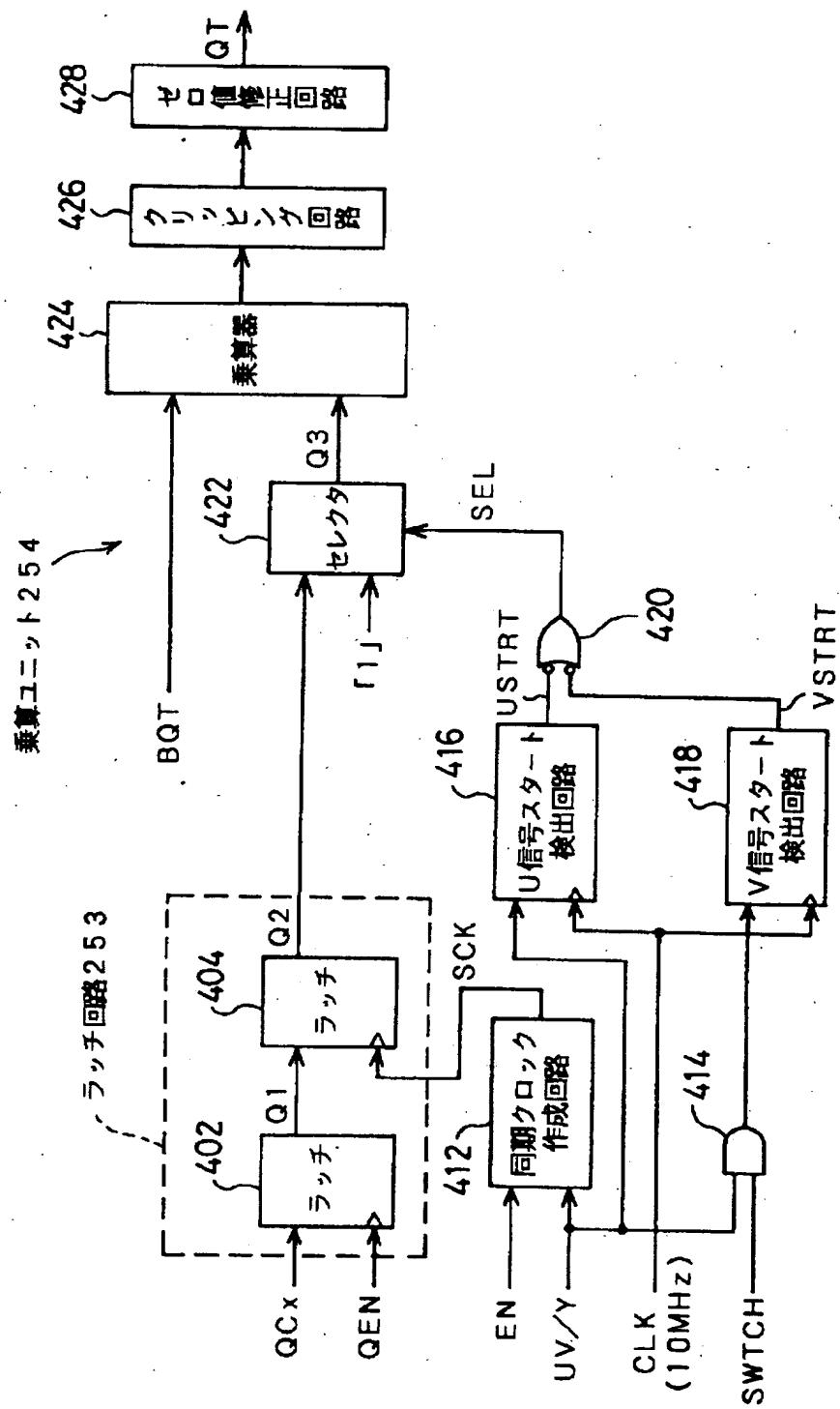
$$\text{NNNN} = [\text{ブロック数}] - 1$$

$$\text{SSSS} = 1$$

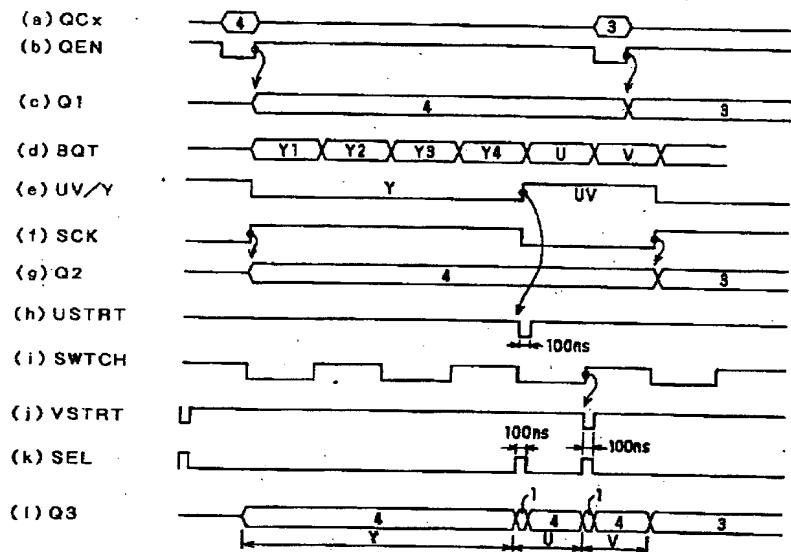
【図22】



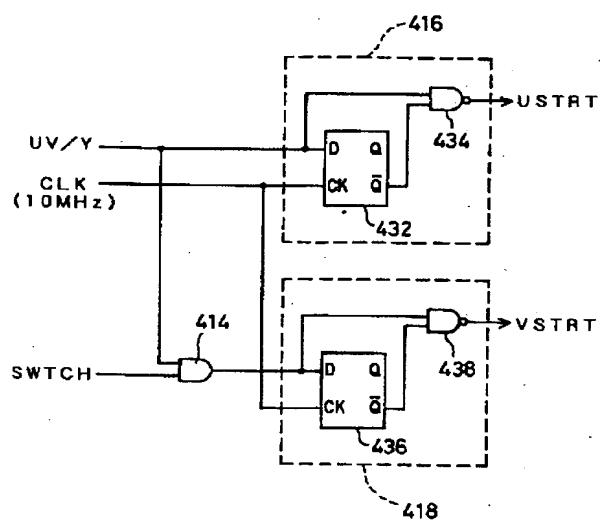
【図23】



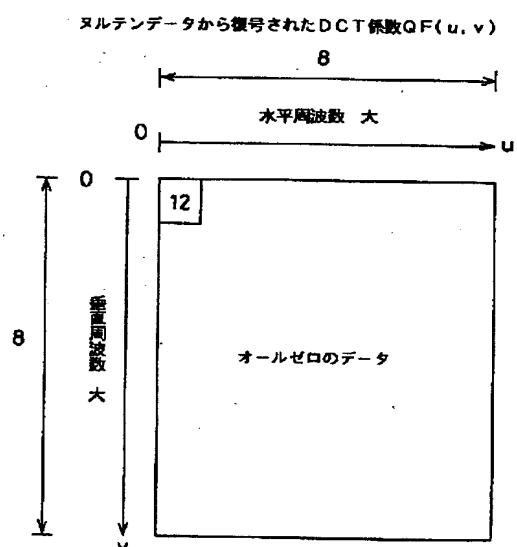
【図24】



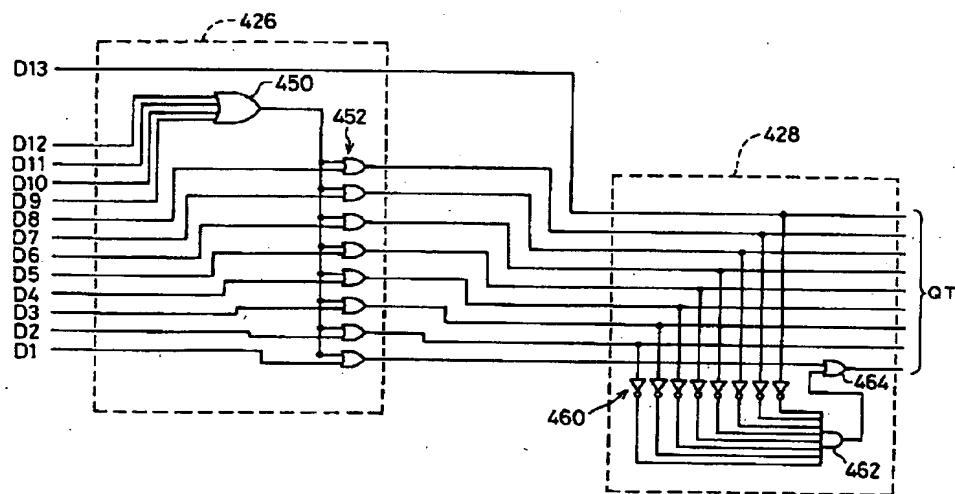
【図25】



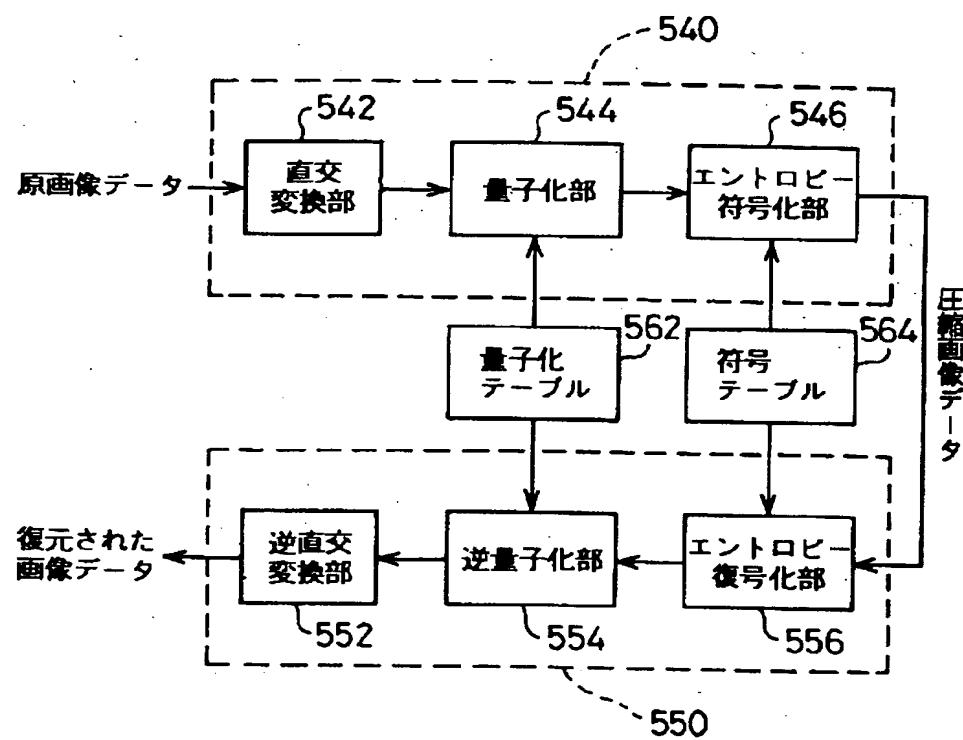
【図28】



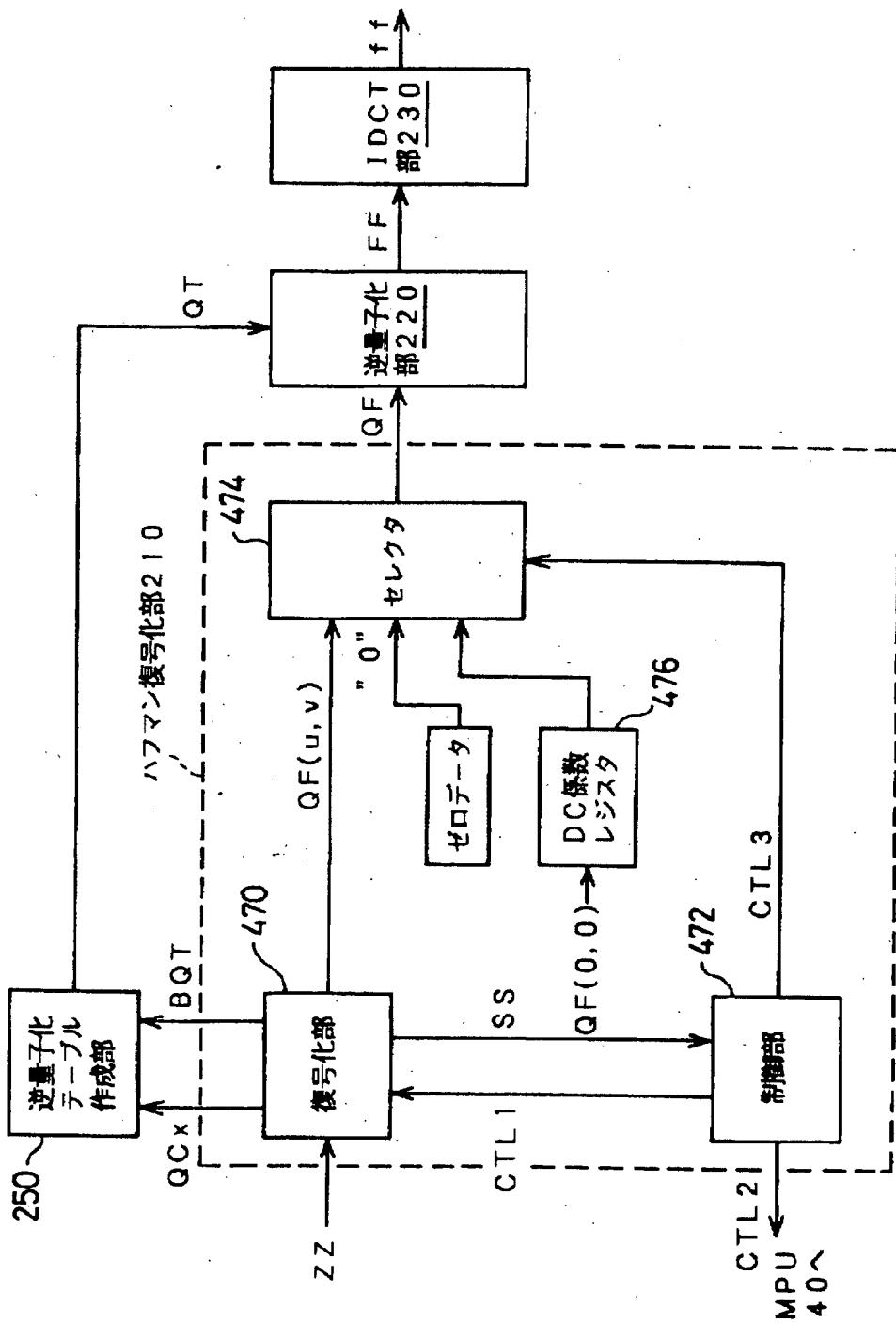
【図26】



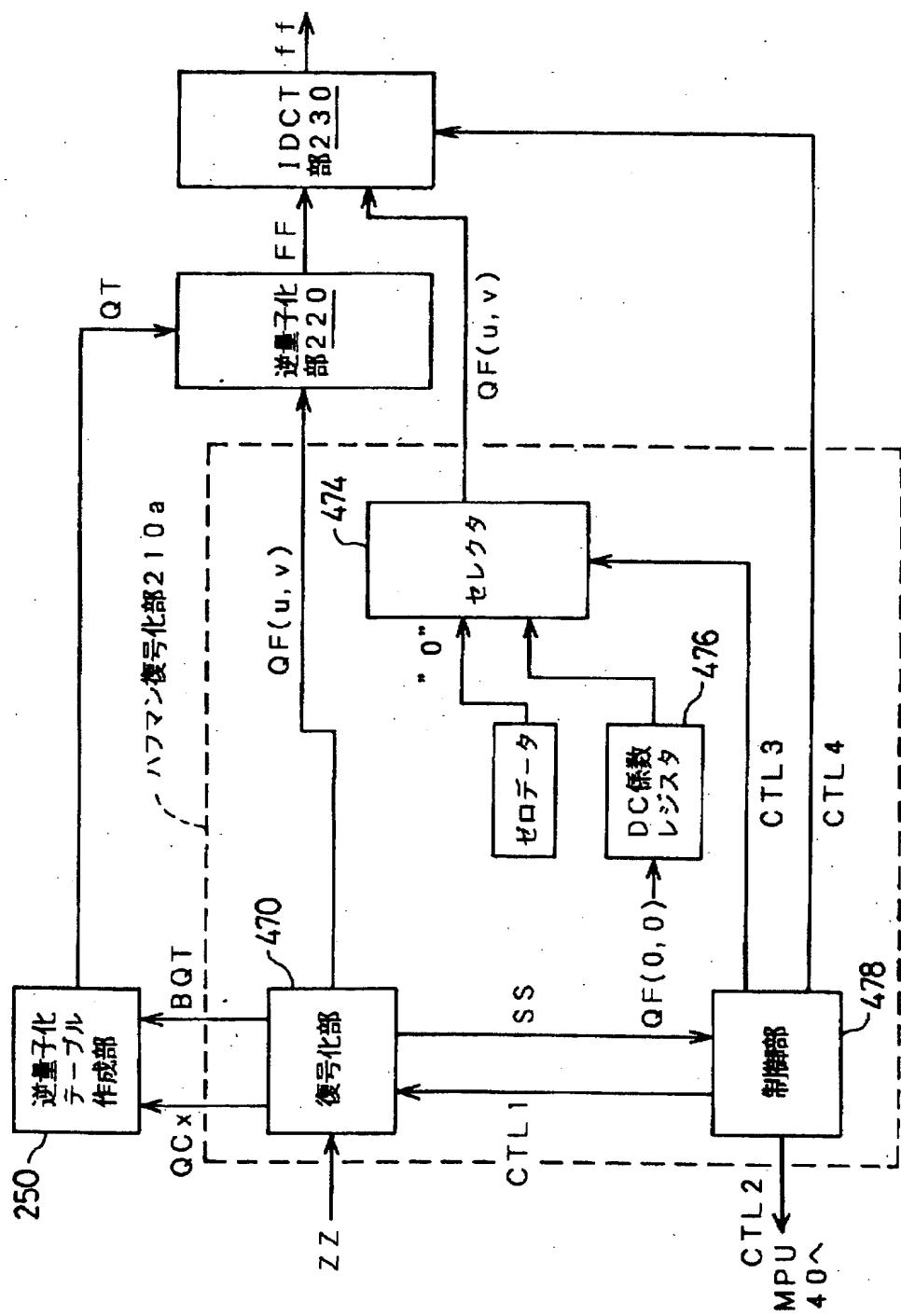
【図31】



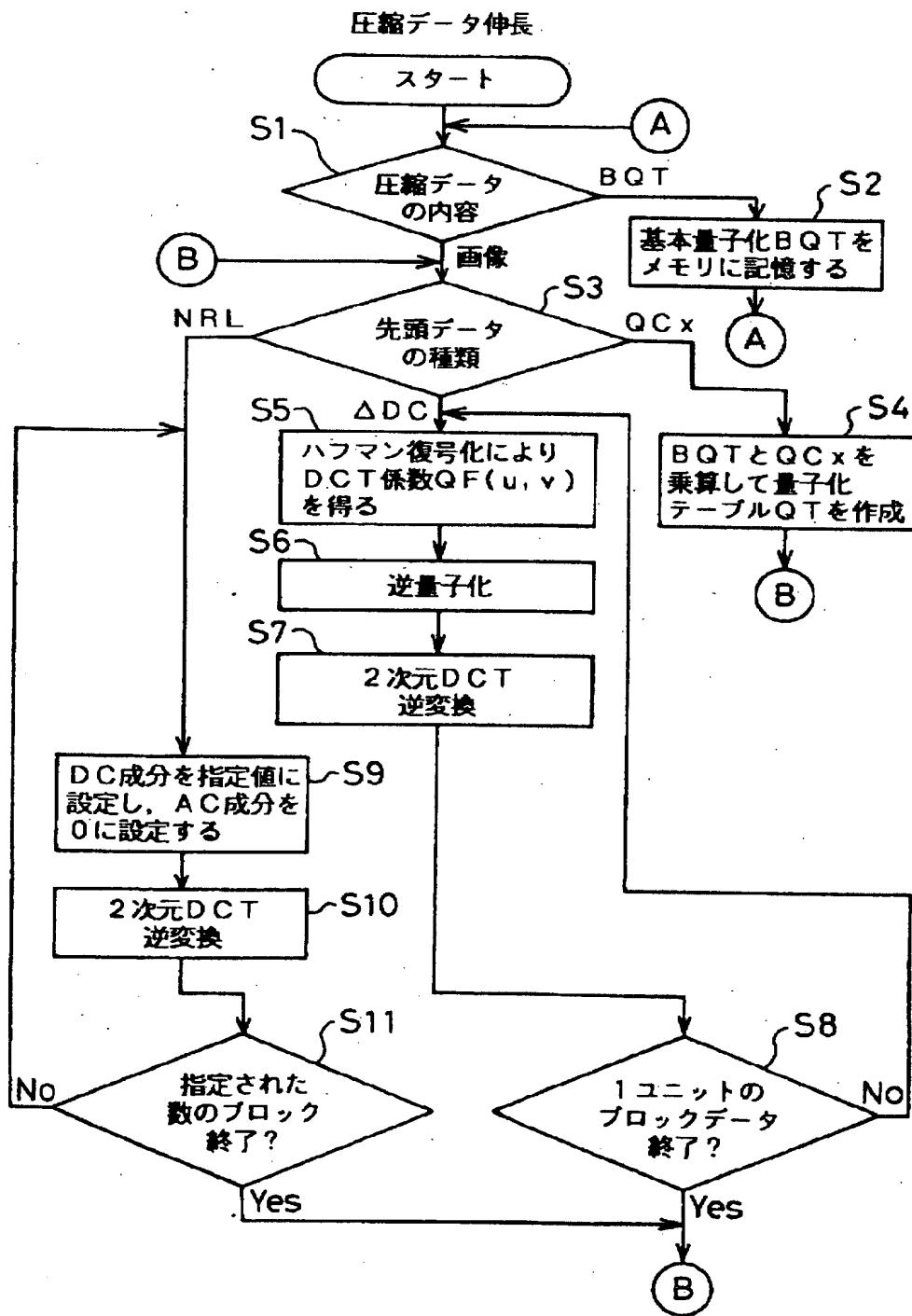
[図27]



[図29]



【図30】



フロントページの続き

(51)Int.Cl. <sup>b</sup>	識別記号	序内整理番号	F I	技術表示箇所
H 03M 7/40		8842-5J		
H 04N 1/41	B			

**This Page is Inserted by IFW Indexing and Scanning  
Operations and is not part of the Official Record**

**BEST AVAILABLE IMAGES**

Defective images within this document are accurate representations of the original documents submitted by the applicant.

Defects in the images include but are not limited to the items checked:

- BLACK BORDERS**
- IMAGE CUT OFF AT TOP, BOTTOM OR SIDES**
- FADED TEXT OR DRAWING**
- BLURRED OR ILLEGIBLE TEXT OR DRAWING**
- SKEWED/SLANTED IMAGES**
- COLOR OR BLACK AND WHITE PHOTOGRAPHS**
- GRAY SCALE DOCUMENTS**
- LINES OR MARKS ON ORIGINAL DOCUMENT**
- REFERENCE(S) OR EXHIBIT(S) SUBMITTED ARE POOR QUALITY**
- OTHER:** \_\_\_\_\_

**IMAGES ARE BEST AVAILABLE COPY.**

**As rescanning these documents will not correct the image problems checked, please do not report these problems to the IFW Image Problem Mailbox.**